Outline
 Introduction
 Two-stage Stackelberg Game
 Equilibrium and Algorithm
 Trace-based Simulations
 Conclusions
 Q&A

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Demand Response of Data Centers: A Real-time Pricing Game between Utilities in Smart Grid

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- Motivation
- Challenges
- Approaches
- 2 Two-stage Stackelberg Game
 - Stage II: DCs' Cost Minimization
 - Stage I: Non-Cooperative Pricing Game
- Equilibrium and Algorithm
 - Backward Induction: Optimal Solutions at Stage II
 - Backward Induction: Nash equilibrium at Stage I
 - Distributed Algorithm
 - Trace-based Simulations
 - Conclusions





Motivation

- DCs consumed 1.5% of the worldwide electricity supply in 2011 and this fraction is expected to grow to 8% by 2020
- 2 DC operators paid more than \$10M (Qureshi 2009)
- DC operators can save 5% 45% cost by leveraging time and location diversities of prices
- The electricity price applying on DC does not change with demand
- Demand Response of Data Centers: receiving consideration
- Pricing for DR: a *right price* not only at the *right time* but also on the *right amount* of demand



Challenges





Approaches



Stage II: DCs' Cost Minimization

Optimization Problem

DC :	minimize	$\sum_{t=1}^{T}\sum_{i=1}^{I}oldsymbol{e}_{i}(t)oldsymbol{p}_{i}(t)+\omega oldsymbol{d}_{i}\lambda_{i}^{2}(t)$	(1)
	subject to	$\frac{1}{s_i(t)\mu_i - \lambda_i(t)} + d_i \leq D_i, \forall i$	(2)
		$\sum_{i=1}^{l} \lambda_i(t) = \Lambda(t), \ \forall t,$	(3)
		$0 \leq s_i(t) \leq S_i, \ \forall i, t,$	(4)
		$0 \leq \lambda_i(t) \leq \mathbf{s}_i(t)\mu_i, \ \forall i, t,$	(5)
	variables	$\boldsymbol{s}_i(t), \lambda_i(t), \ \forall i, t.$	(6)

Stage I: Utility Revenue and Cost

Revenue

$$\mathcal{R}_i(\boldsymbol{p}(t)) = (\boldsymbol{e}_i(\boldsymbol{p}(t)) + \boldsymbol{B}_i(\boldsymbol{p}_i(t)))\boldsymbol{p}_i(t)$$

Cost
$$C_i(p(t)) = \gamma ELI = \gamma \left(\frac{e_i(p(t)) + B_i(p_i(t))}{C_i(t)} \right)^2 C_i(t),$$

Stage-I: A Non-Cooperative Pricing Game Formulation

- *Players*: the utilities in the set \mathcal{I} ;
- Strategy: $p_i^l \le p_i(t) \le p_i^u$, $\forall i \in \mathcal{I}, t \in \mathcal{T}$; Payoff function: $\sum_{t=1}^{T} u_i(p_i(t), p_{-i}(t)), \forall i \in \mathcal{I}$.

$$u_i(p_i(t), p_{-i}(t)) = \mathcal{R}_i(p(t)) - \mathcal{C}_i(p(t)),$$

Backward Induction: Optimal Solutions at Stage II

Observe that the QoS constraint must be active

$$\boldsymbol{s}_i(\lambda_i) = \left[rac{1}{\mu_i} \left(\lambda_i + ilde{\boldsymbol{D}}_i^{-1}
ight)
ight]_{0}^{\boldsymbol{S}_i},$$

Optimization Problem

 $f_i($

$$\mathsf{DC}': \min_{\lambda} \qquad \sum_{i=1}^{l} f_{i}(\lambda_{i}) \qquad (7)$$
s.t.
$$\sum_{i=1}^{l} \lambda_{i} = \Lambda, \qquad (8)$$

$$\lambda_{i} \ge 0, \ \forall i, \qquad (9)$$

$$\lambda_{i}) := \omega d_{i}\lambda_{i}^{2} + p_{i}\left(a_{i} + \frac{b_{i}}{\mu_{i}}\right)\lambda_{i} + p_{i}\left(e_{b} + \frac{b_{i}\tilde{D}_{i}^{-1}}{\mu_{i}}\right)$$

Backward Induction: Optimal Solutions at Stage II



$$e_i^*(p) = rac{A_i^2 p_i}{2\omega d_i} (rac{1}{\hat{d}d_i} - 1) + rac{A_i}{2\omega \hat{d}d_i} \sum_{j
eq i} rac{A_j p_j}{d_j} + rac{A_i \Lambda}{\hat{d}d_i} + rac{b_i}{\mu_i ilde{D}_i} + e_i^b.$$

Backward Induction: Nash equilibrium at Stage I

Existence: concave game (Rosen 1965)

Best response updates
$$p_{i}^{(k+1)} = \mathcal{BR}_{i}\left(p_{-i}^{(k)}\right) = \left[\frac{1/2 - \gamma N_{i}/C_{i}}{1 - \gamma N_{i}/C_{i}}\frac{h\left(p_{-i}^{(k)}\right)}{(-N_{i})}\right]_{\mathcal{P}_{i}}, \forall i$$

Uniqueness:
$$p_i^e = \mathcal{BR}_i(p_{-i}^e), \forall i$$

Condition
 $\omega \ge \max_i \left\{ \frac{A_i \sum_{j \ne i} A_j/d_j - A_i^2 \hat{d}(1 - 1/(d_i \hat{d}))}{2\beta_i \hat{d} d_i} \right\},$

Distributed Algorithm



 Outline
 Introduction
 Two-stage Stackelberg Game
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 Conclusions
 Q&A

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Trace-based Simulations



 Outline
 Introduction
 Two-stage Stackelberg Game
 Equilibrium and Algorithm
 Trace-based Simulations
 Conclusions
 Q&A

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Dynamic Prices



DC's cost and utilities' profit



(a) FIU trace

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Summary and Future Work

Summary

- DR of DCs: interactions between DCs and utilities via pricing
- Two-stage Stackelberg game: utilities are leaders, DCs are follows
- Flatten the demand over time and space

Future Work

- Deadline constraint
- Workload estimation errors
- Risk consideration

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