

NETWORK CONFIGURATION IS HARD!



- High-level tasks are realized through low-level commands and scripts: **hard to understand**
- Distributed configuration: **hard to manage**
- Variety of network-wide tasks cause changes to the network: **lots of dynamics**
- No changes are checked for correctness:
error-prone

SURVEY WITH NETWORK OPERATORS



20% make changes more than once a day

89% are *never* completely certain that changes will not introduce a new bug

82% are concerned changes might break existing functionality unrelated to the changes

SURVEY WITH NETWORK OPERATORS

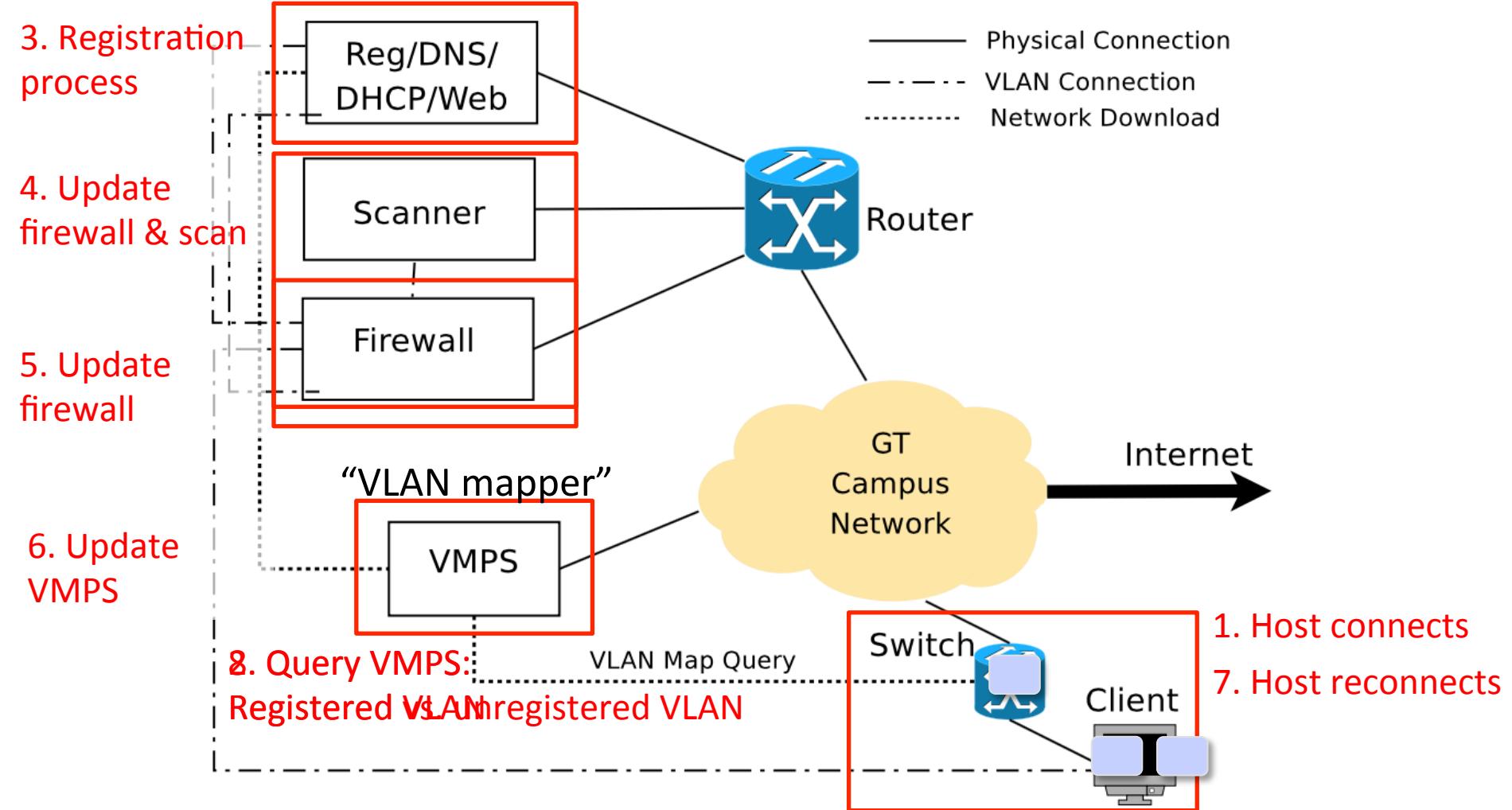


20% make changes more than once a day

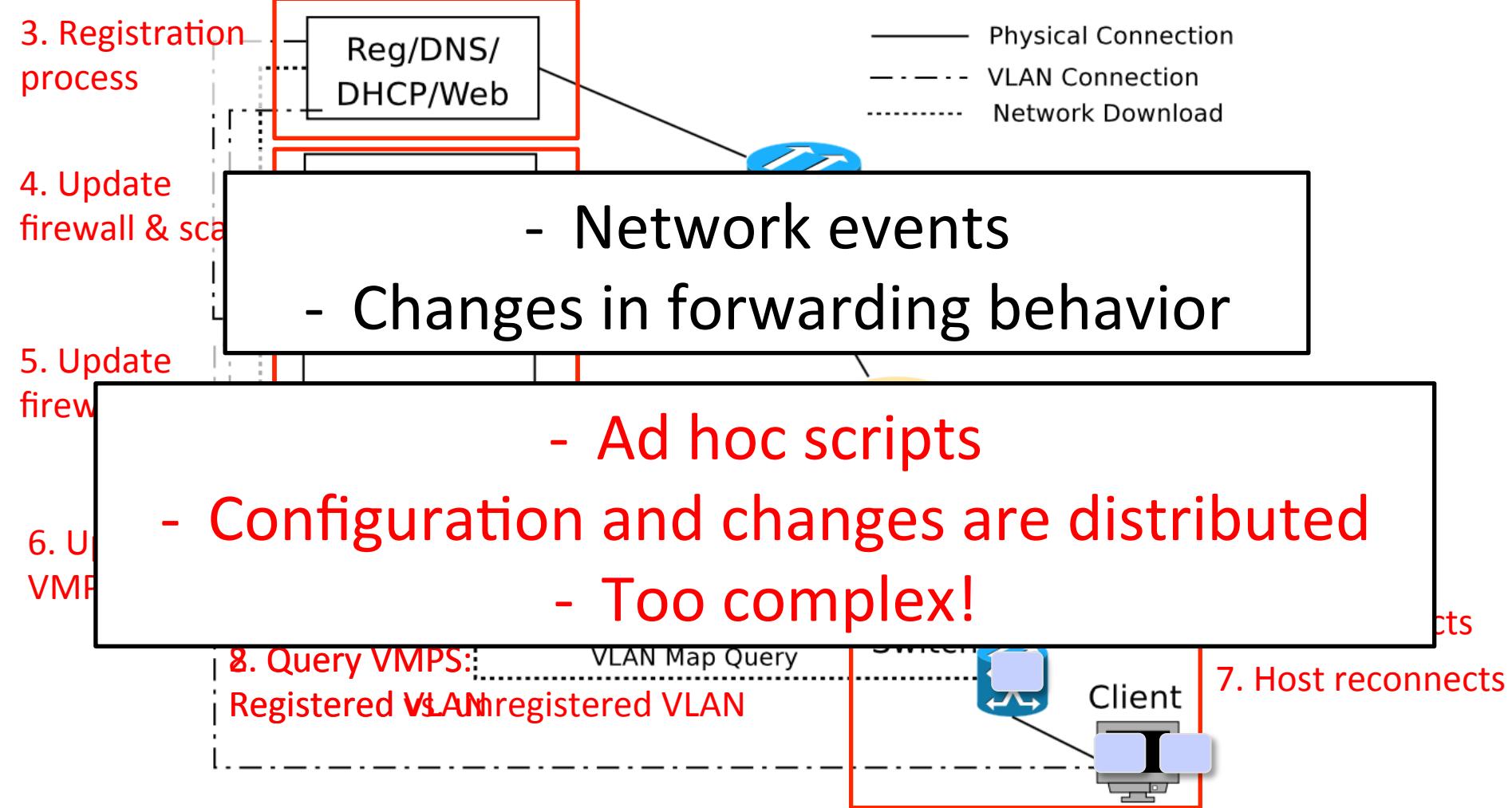
89% are *never* completely certain that changes will not introduce a new bug

“You should track down those 10-20% of operators who say they are always certain. They are LYING.”

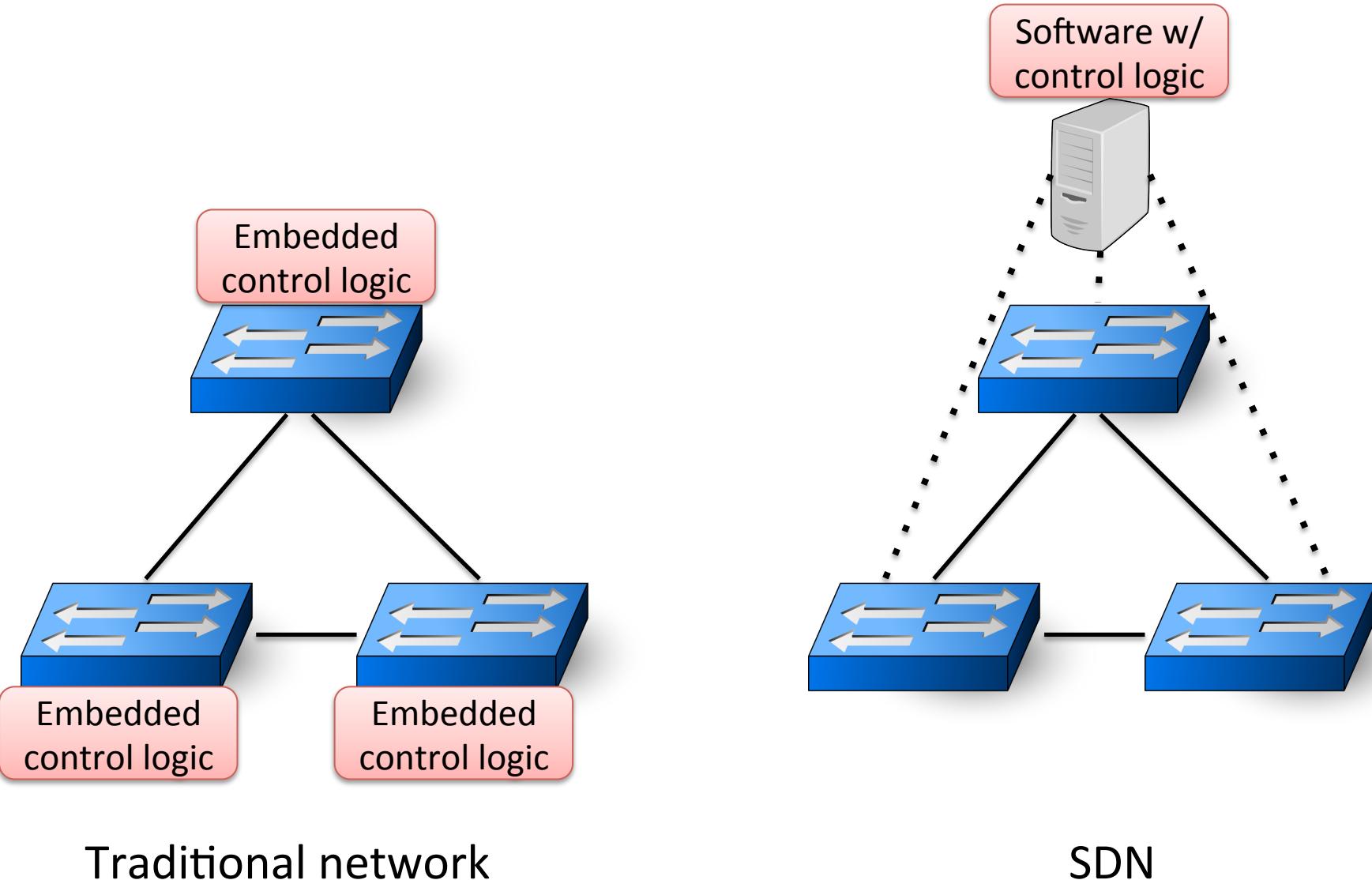
MOTIVATING EXAMPLE: THE START SYSTEM



MOTIVATING EXAMPLE: THE START SYSTEM



SOFTWARE-DEFINED NETWORKING



Traditional network

SDN

SDN IS NOT A SILVER BULLET

- Low-level commands & scripts:
hard to understand
 - Distributed configuration:
hard to manage
 - Many network-wide tasks, lots of changes: **lots of dynamics**
 - No correctness guarantee:
error-prone
- 
- Programs: e.g., C++, Java, Python, Pyretic
- Central control
- Unsolved
- Unsolved

WHAT SDN PLATFORM NEEDS

- Guidance on how to implement a network control program
 - How to provide ***dynamic control*** that handles arbitrary network events
 - E.g, Intrusion detection, traffic load shift, etc
- Verification and guarantees of program's correctness
- Huge missed opportunities in software

DIFFERENT TYPES OF NETWORK EVENTS



- Network traffic
 - Traffic load increase/decrease, security incidents
- User-specific
 - User authentication, excessive data usage
- Data-plane events
 - Topology change, switch/link failures
- ...

DIFFERENT REACTIONS TO AN EVENT

Event

Operators

Reaction



“Only block that infected host”

Host is infected!



“Block all communications in the network!”

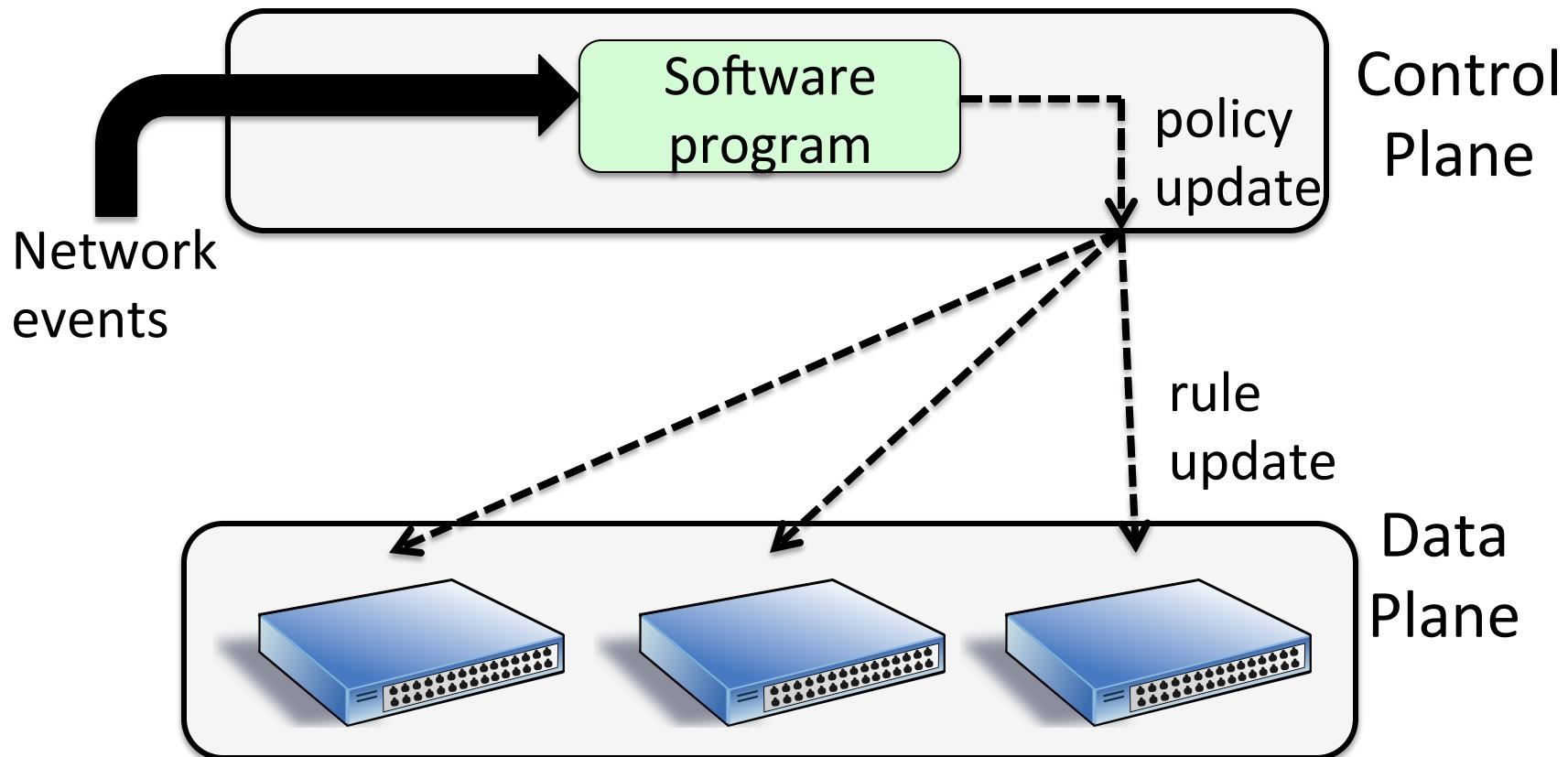


“Direct communication to our internal honeypot”

**Network events and dynamic reactions to them
should be programmatically encoded in the
network control program by operators**

DYNAMIC NETWORK CONTROL PROGRAM

- **Software program** that embeds event – reaction relationships



UNANSWERED QUESTIONS



How to **embed event-reaction logic** in software?

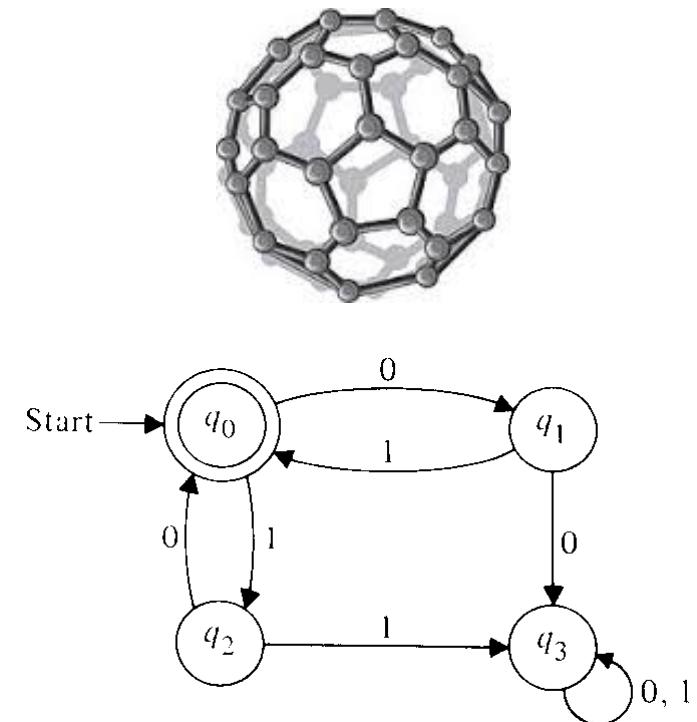
How to **verify** that the program will make
changes correctly?

Kinetic tackles these questions

- **Domain specific** language and control platform
- Helps create SDN control programs that **embed custom event-reaction** relationships
- **Verifies** program's **correctness**

OUR APPROACH

- Domain specific language
 - Constrained, but structured
- Express changing behavior as a *finite state machine*
- Verify program's correctness with a model checker (NuSMV)

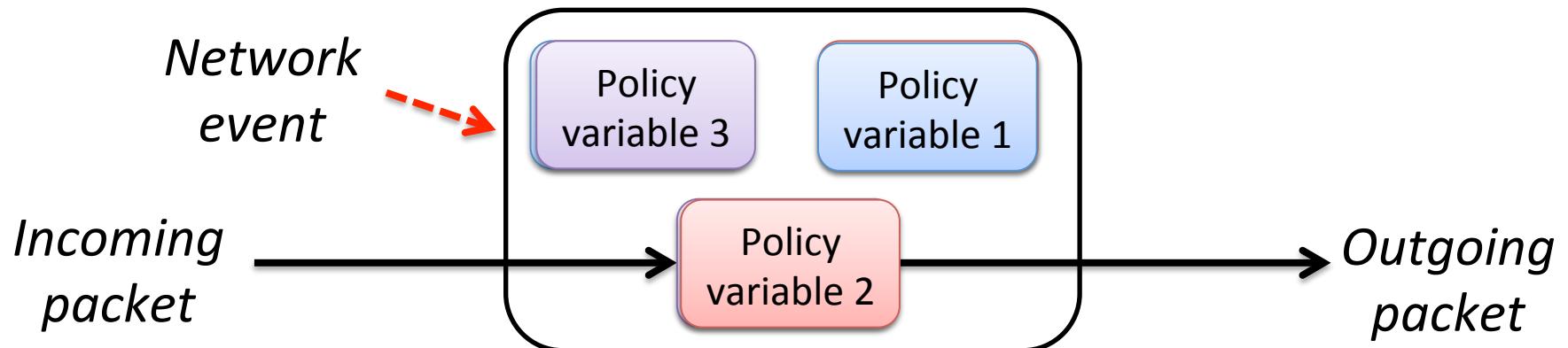


KINETIC'S DOMAIN SPECIFIC LANGUAGE

- Embedded in Python
- Borrows some abstractions from Pyretic
 - Encodes forwarding behavior in a *policy variable*

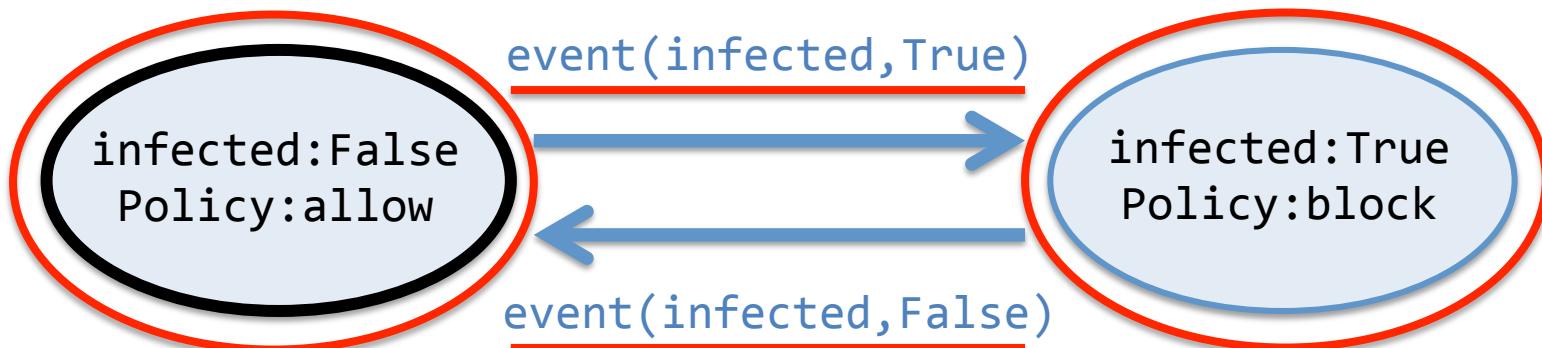


- New constructs and functions to express policies that respond to *changing conditions*



IDS EXAMPLE IN KINETIC

- **Event:** *infected*
- **State:** policy variable's value
 - *allow* or *block* packet

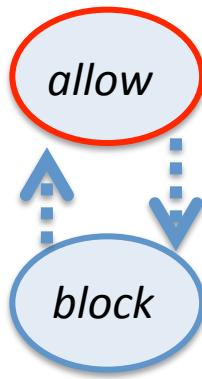


There are many different flows
Each flow can have its own independent FSM

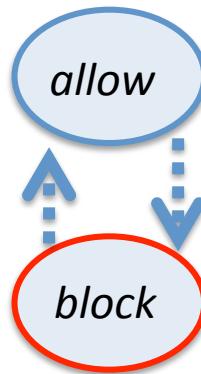
DECOMPOSING TO MULTIPLE FSMS

- FSM instance is instantiated per flow

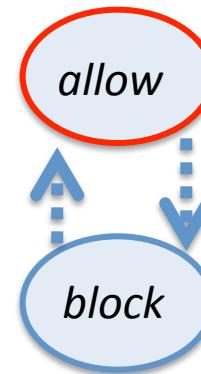
Host 1



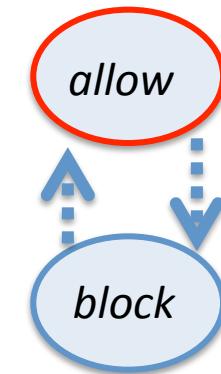
Host 2



Host 3



Host N



of hosts: N

Total # of states: $2N$

Total # of transitions: $2N$

State representation is *Linear* in N
(instead of geometric)

- In IDS example, flow is defined by source IP address (host)
- Other policies may require more flexibility (e.g., need to group packets by location)
- **Located Packet Equivalence Class (LPEC)**
 - Programmer abstraction to define *flow*

```
def lpec(pkt):
    return match(dstip=pkt['dstip'])
```

KINETIC VERIFICATION PROCESS

- Kinetic verifies correctness of the program
 - User-specified temporal properties
 - Verifies ***current and future*** forwarding behavior based on network events
- Verification process is ***automated***
 - Constrained but structured language allows automatic parsing and translation of program
- Verification runs ***before program's deployment***

VERIFICATION PROCESS

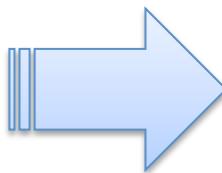
Kinetic program

```
@transition
def infected(self):
    self.case(occured(self.event),self.event)

@transition
def policy(self):
    self.case(is_true(V('infected')),C(drop))
    self.default(C(identity))

self.fsm_def = FSMDef(
    infected=FSMVar(type=BoolType(),
                      init=False,
                      trans=infected),
    policy=FSMVar(type=Type(Policy,{ drop ,
                                identity }),,
                  init=identity,
                  trans=policy))
```

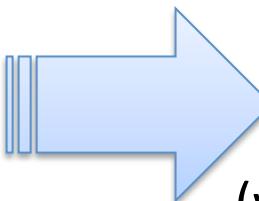
Automatically
generates



NuSMV FSM model

```
MODULE main
VAR
    policy   : {identity ,drop };
    infected : boolean;
ASSIGN
    init(policy)  := identity ;
    init(infected) := FALSE;
next(policy) :=
    case
        infected : drop ;
        TRUE      : identity ;
    esac ;
next(infected) :=
    case
        TRUE      : {FALSE,TRUE} ;
    esac ;
```

User-specified
temporal properties



True or False
(w/ counter-example)

EXAMPLES OF TEMPORAL PROPERTIES

- If a host is infected, drop packets from that host

$\boxed{A}\boxed{G}$ (infected \rightarrow $\boxed{A}\boxed{X}$ policy=drop)

For all possible transitions from
current state,

For all current and future
states,

For all possible transitions
from current state,

For the next state,

- If host is authenticated either by Web or 802.1X, and is not infected, packets should never be dropped.

$\text{AG} ((\text{authenticated_web} \mid \text{authenticated_1x}) \& \neg \text{infected} \rightarrow \text{AX policy} \neq \text{drop})$

EVALUATION

- Usability evaluation
 - User study against over 870 participants
 - Lines of code comparison with other SDN solutions
- Performance and scalability
 - Event handling and policy recompilation

KINETIC: USER STUDY

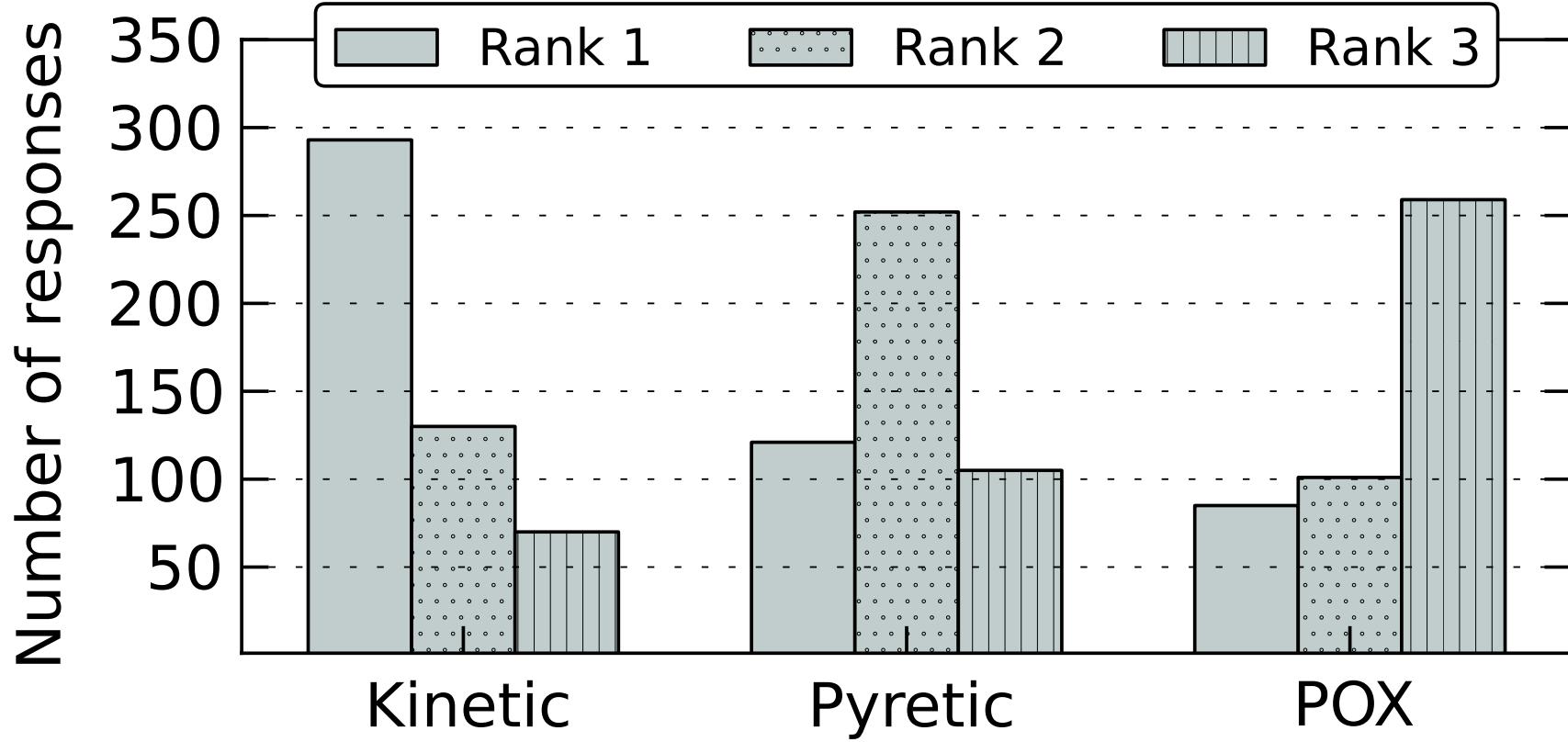
- Demographic

Profession	Experience (years)		
Operator	216	1	32
Developer	251	1-5	310
Student	123	5-10	187
Vendor	80	10-15	150
Manager	69	15-20	122
Other	138	> 20	73
Total	877		874

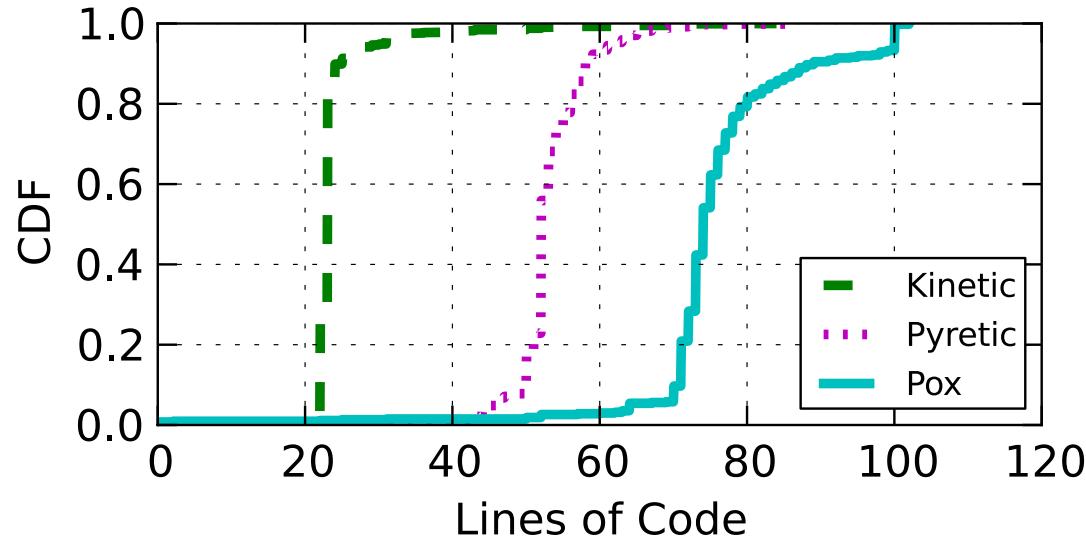
- Task

- Implement an enhanced IDS program with Kinetic, Pyretic, and POX.

RANK PLATFORMS BY PREFERENCE



LINES OF CODE COMPARISON



Programs	FL	POX	Pyretic	Kinetic
IDS/firewall	416	22	46	17
Mac learner	314	73	17	33
Server load balance	951	145	34	37
Stateful firewall	<i>None found</i>	<i>None found</i>		41
Usage-based rate limiter	<i>None found</i>	<i>None found</i>	<i>None found</i>	30

NOTABLE QUOTES

- Why did you like Kinetic?
 - **FSM-based structure and support for intuition**

“Kinetic is more intuitive: the only things I need to do is to define the FSM variable”

“intuitive and easy to understand”

“Programming state transitions in FSMs makes much more sense”

- **More concise**

“reduces the number of lines of code”

“the logic is more concise”

NOTABLE QUOTES

- Why *didn't* you like Kinetic?

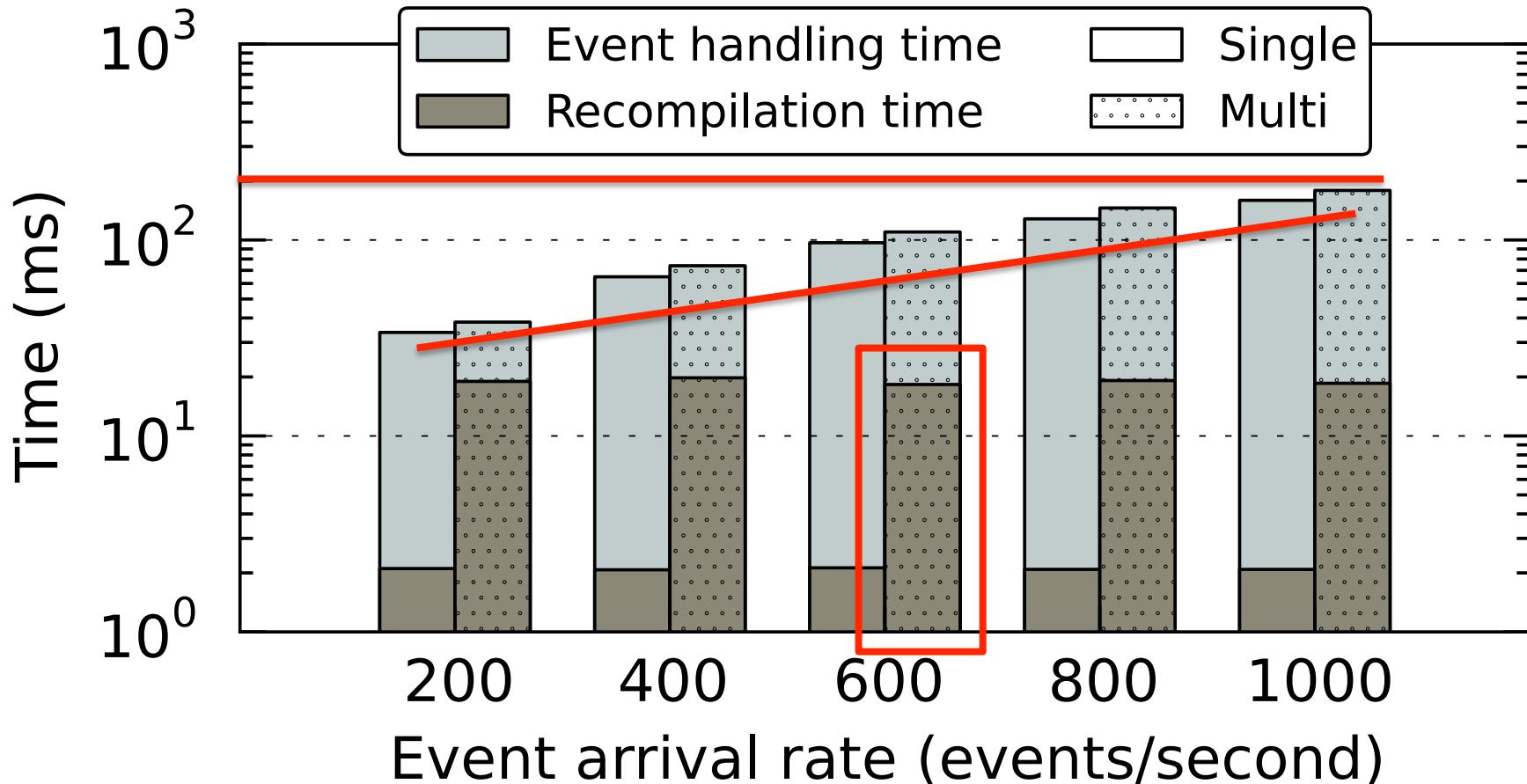
- Steep learning curve

“Kinetic took less time and was actually more understandable ...[but] the structure was very cryptic”

- Not friendly when finding why program is wrong

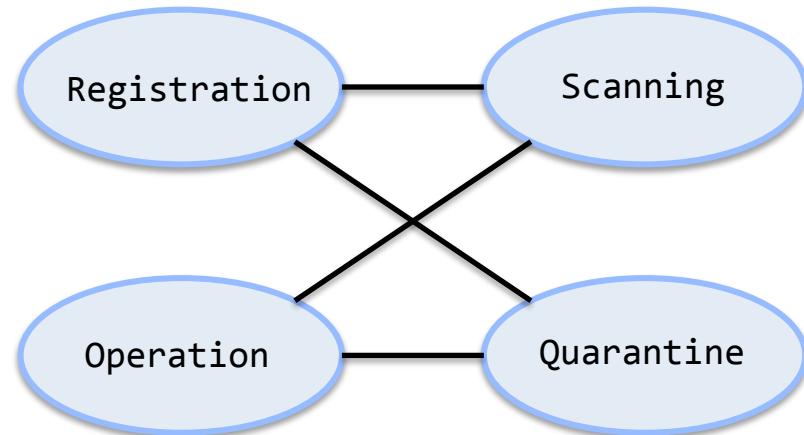
“I spent a lot more time chasing down weird bugs I had because of things I left out or perhaps didn’t understand”

Event handling and policy recompilation

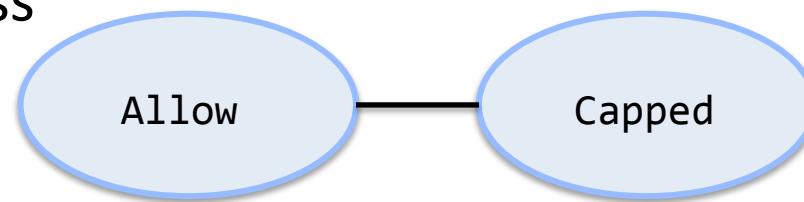


KINETIC: REAL DEPLOYMENTS

- Campus network
 - Functional access control system
 - Deployed SDN-enabled switches over 3 buildings



- Home network
 - Usage-based access control
 - Deployed 21 SDN-enabled wireless routers over 3 continents
 - Jul., 2012 – Feb., 2014
 - Presented in ACM CHI 2015



KINETIC TAKEAWAYS

- **Domain specific** language and control platform
 - Program **encodes event-reaction logic**
- **Extensive user study** shows that
 - Much **easier to express dynamics** in the network
 - Helps to **reduce lines of code**
- **Scales well** to large networks and lots of events
- **Verification process reduces bugs** in programs

DISCUSSION & FUTURE WORK



- Combining with verifications in other stacks
 - Consistent updates to data plane
 - Verification of data-plane state
- More dynamic network policies
 - Should collect more real network policies
 - Need public repository

More about Kinetic:

<http://kinetic.noise.gatech.edu>

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Questions?