

Stroboscope: Declarative Network Monitoring on a Budget



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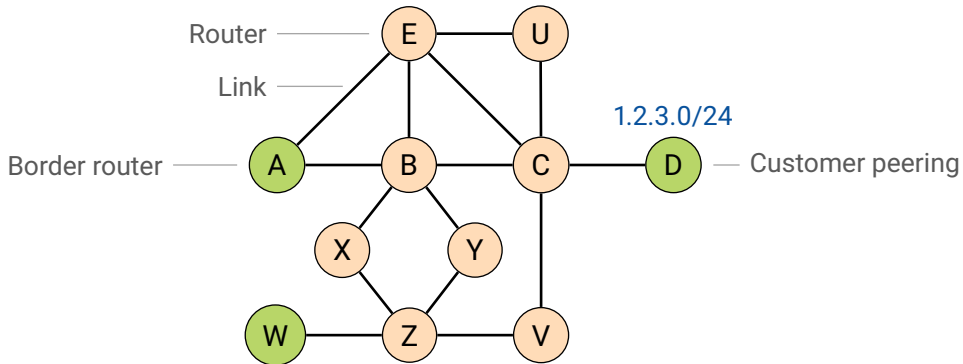
April 11, 2018

Joint work with

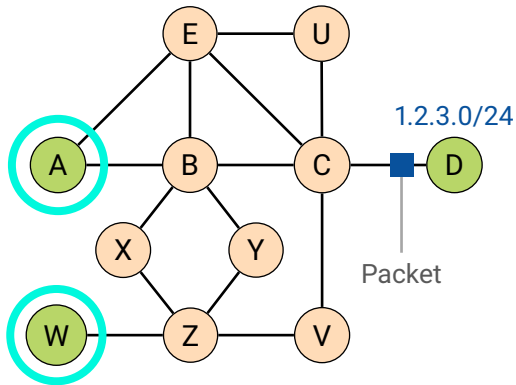
T. Bühler (ETH Zürich), I. Poese (BENOCS), S. Vissicchio (UCL) and L. Vanbever (ETH Zürich)

Adapted from original picture © Michael Magg, 2007, CC-BY-SA 3.0

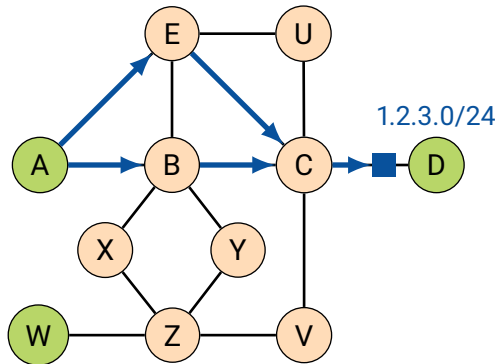
Consider this example ISP network topology



What is the ingress router for this packet arriving at router D?

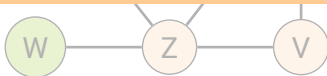


Which paths does the traffic follow?



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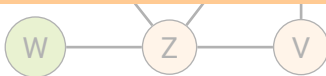
Tracking flows network-wide requires to **match measurements** across multiple vantage points



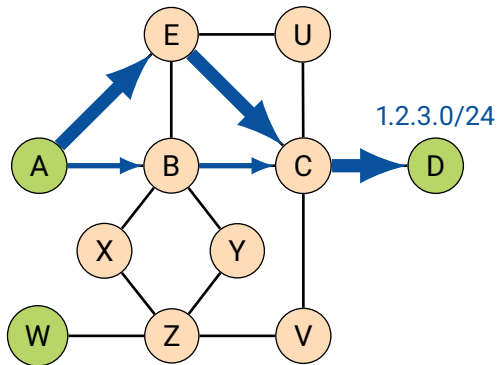
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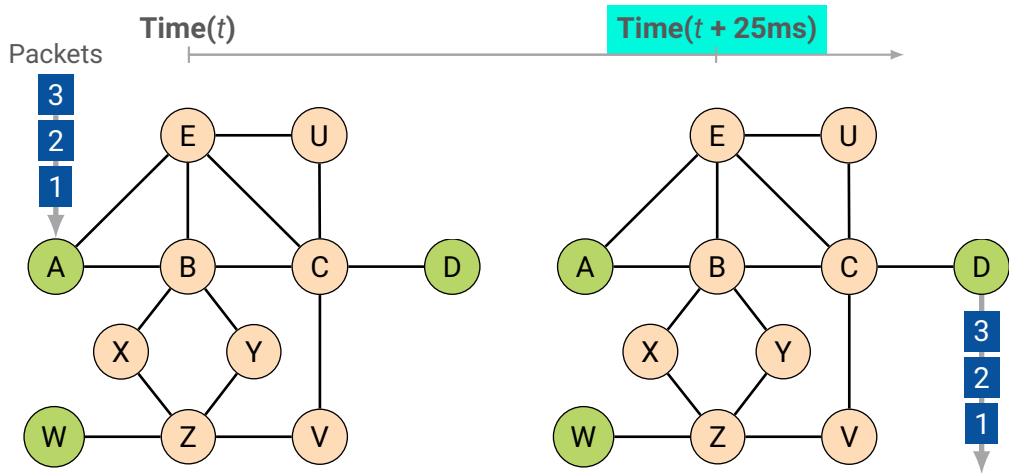
~~NetFlow, ProgME [ToN'11], FlowRadar [NSDI'16]~~



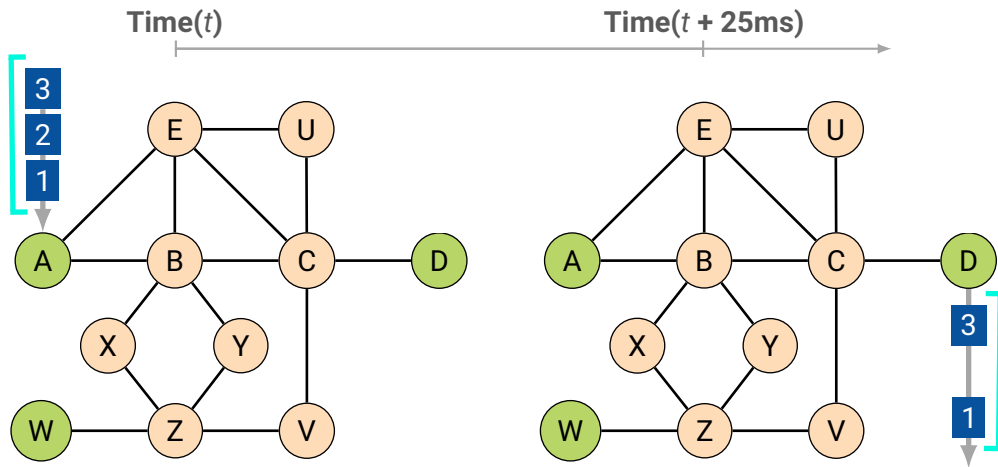
Is traffic load-balanced as expected?



Is the latency acceptable?



Are there losses?

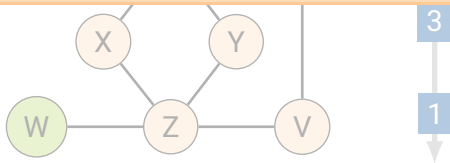
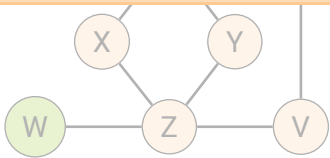


Are there losses?

Time(t)

Time($t + 25\text{ms}$)

Fine-grained data-plane performance metrics require **packet-level visibility** over individual flows



Fined-grained network monitoring is widely researched

Gigascopex [SIGMOD'03]

Planck [SIGCOMM'14]

Everflow [SIGCOMM'15]

Compiling Path Queries [NSDI'16]

Trumpet [SIGCOMM'16]

Marple [SIGCOMM'17]

Fined-grained **ISP** network monitoring poses unique and unmet challenges

- No control over end hosts

Gigascope [SIGMOD'03]

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Fined-grained **ISP** network monitoring poses unique and unmet challenges

- No control over end hosts
- Limited data-plane flexibility
- Limited monitoring bandwidth

~~Gigascopex [SIGMOD'03]~~

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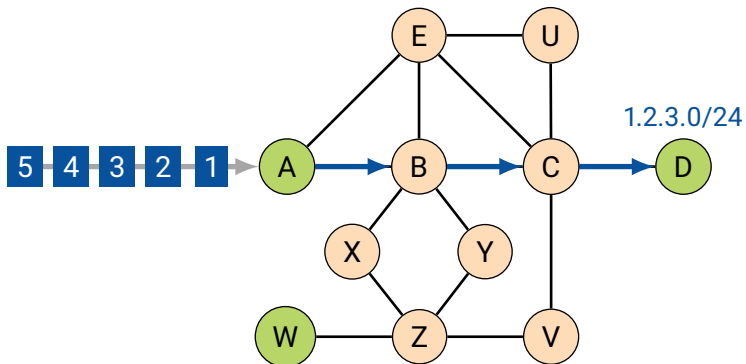
~~Marple [SIGCOMM'17]~~

Stroboscope: Declarative Network Monitoring on a Budget

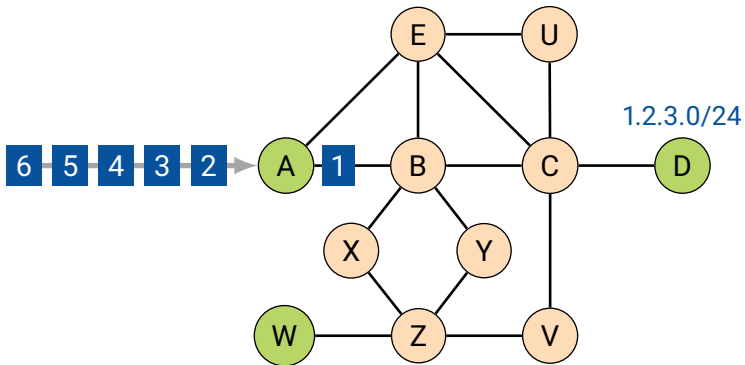


- Collecting traffic slices to monitor networks
- Adhering to a monitoring budget
- Using Stroboscope today

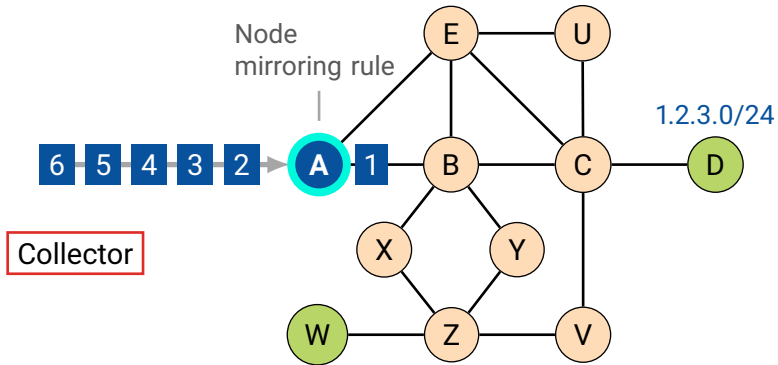
Consider the following flow of packets



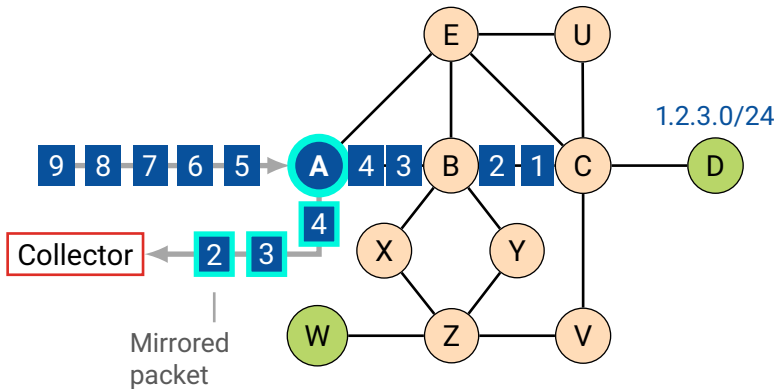
Consider the following flow of packets



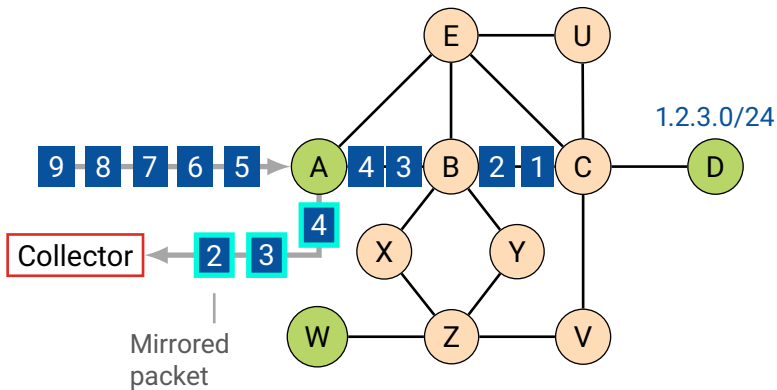
Stroboscope activates mirroring for the flow



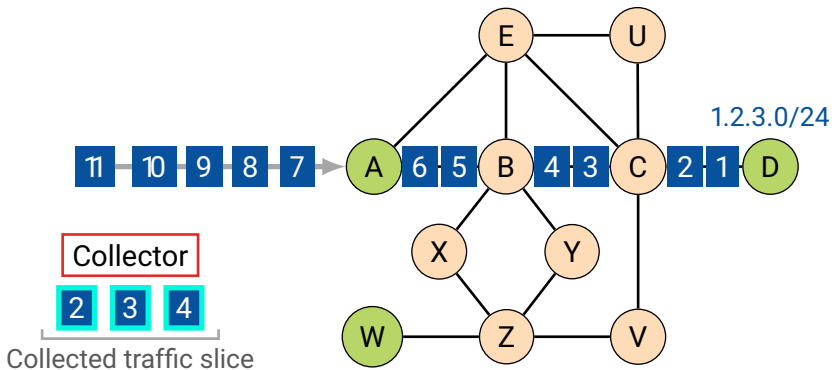
Packets are copied and encapsulated towards the collector



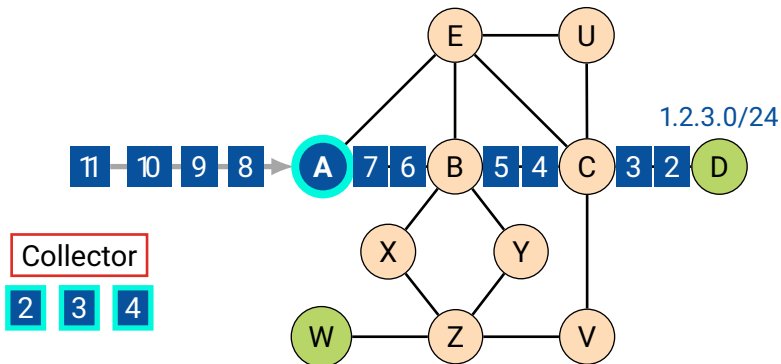
The mirroring rule is deactivated after a preset delay



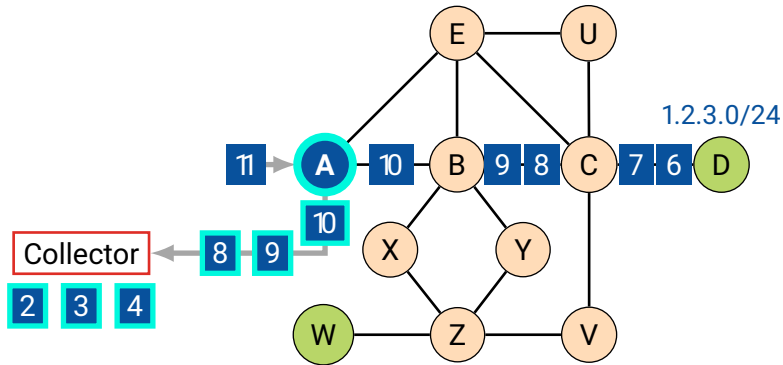
Stroboscope stores the traffic slice for analysis



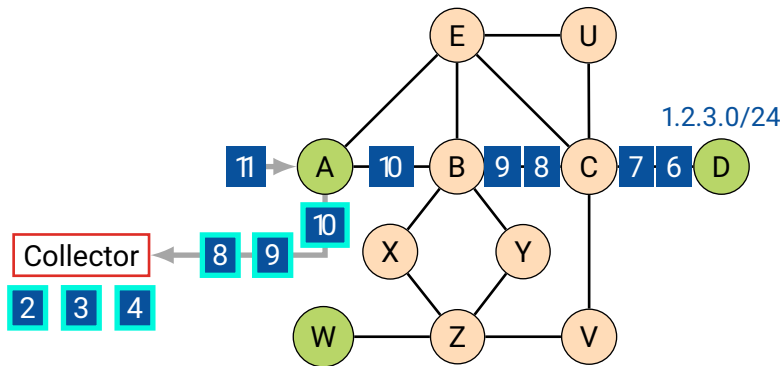
Stroboscope periodically toggles the mirroring rule



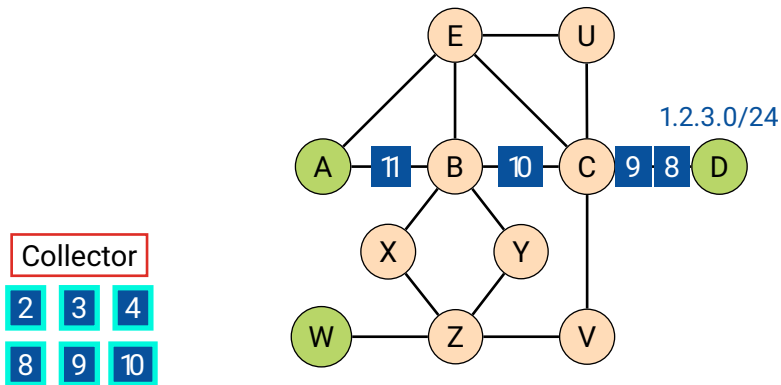
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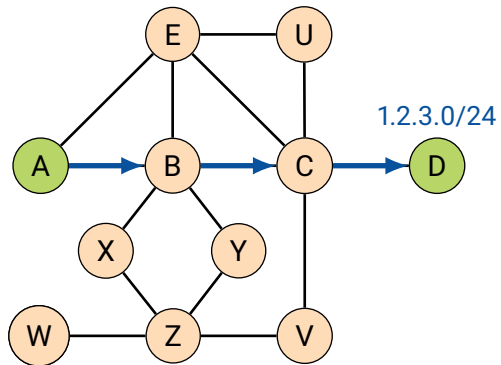
Stroboscope collects multiples traffic slices over time



Stroboscope works with currently deployed routers

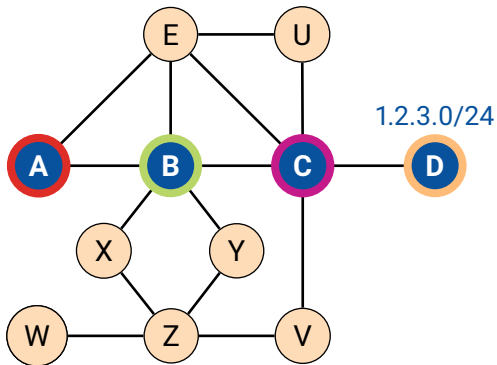
- Most vendors provide traffic mirroring and encapsulation primitives
- The collector activates mirroring for a flow by updating one ACL
- Routers autonomously deactivate mirroring rules using timers
- Traffic slices can be as small as **23 ms** on our routers (Cisco C7018)

Consider the following forwarding path

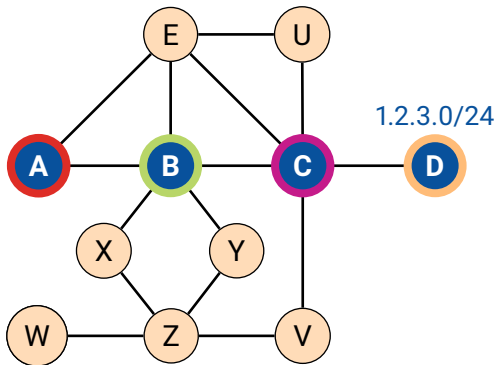
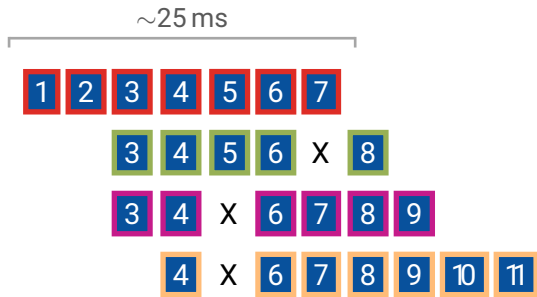


Stroboscope activates mirroring rules along a path

MIRROR 1.2.3.0/24 ON [A B C D]



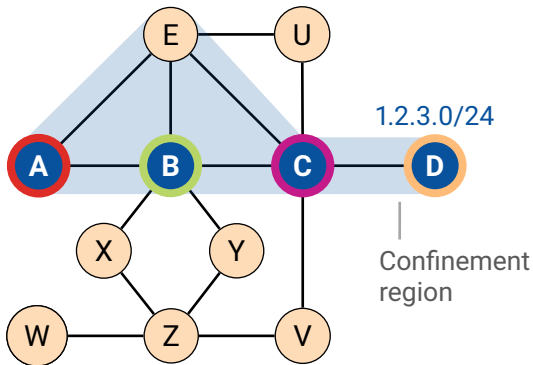
Traffic slices are collected



A CONFINE query mirrors any packet leaving a region

MIRROR 1.2.3.0/24 ON [A B C D]

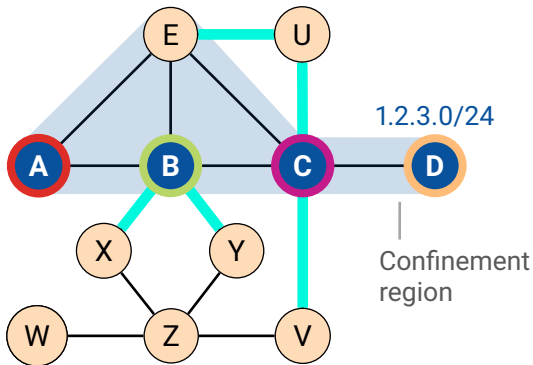
CONFINE 1.2.3.0/24 ON [A B E C D]



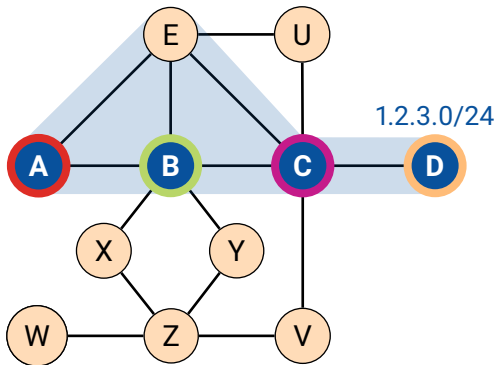
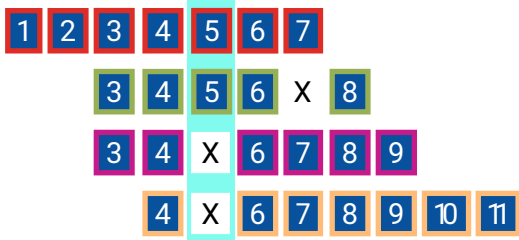
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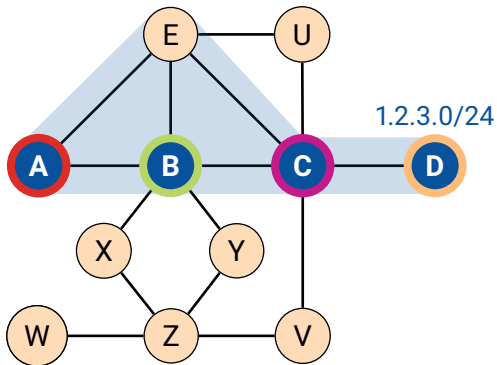
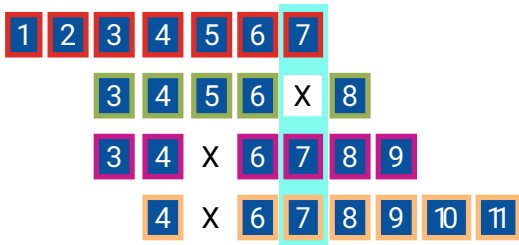
CONFINE 1.2.3.0/24 ON [A B E C D]



Counting packets missing in all last hops of a path estimates loss rates



Counting packets partially following the path estimates load-balancing ratios



Analyzing matching packets across traffic slices
enables fine-grained measurements at scale

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enables fine-grained measurements at scale

Forwarding paths discovery, timestamp reconstruction, payload inspection, ...

Stroboscope: Declarative Network Monitoring on a Budget



- Collecting traffic slices to monitor networks
- Adhering to a monitoring budget
- Using Stroboscope today

Stroboscope defines two types of queries

MIRROR

CONFINE

Stroboscope defines two types of queries



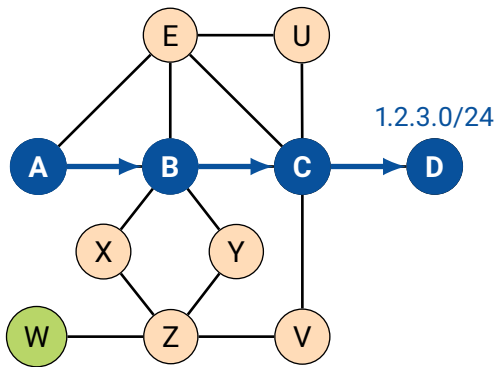
MIRROR



CONFINE

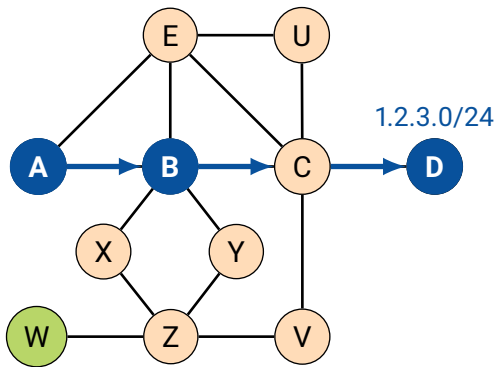
MIRROR queries reconstruct the path taken by packets

MIRROR 1.2.3.0/24 ON [A B C D]



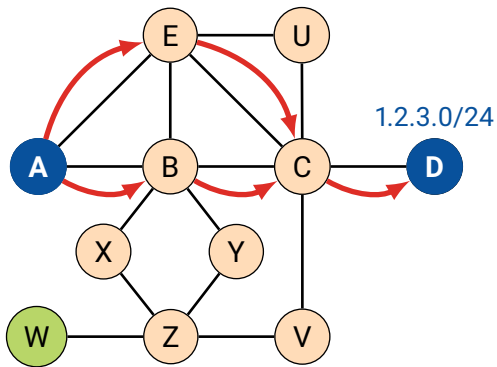
Fewer mirroring rules reduces bandwidth usage

MIRROR 1.2.3.0/24 ON [A B C D]



Too few mirroring rules creates **ambiguity**

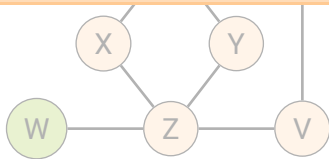
MIRROR 1.2.3.0/24 ON [A B C D]



Too few mirroring rules creates **ambiguity**

MIRROR 1.2.3.0/24 ON [A B C D]

The **Key-Points Sampling** algorithm minimizes mirroring rules and guarantees non-ambiguous reconstructed paths



Stroboscope defines two types of queries



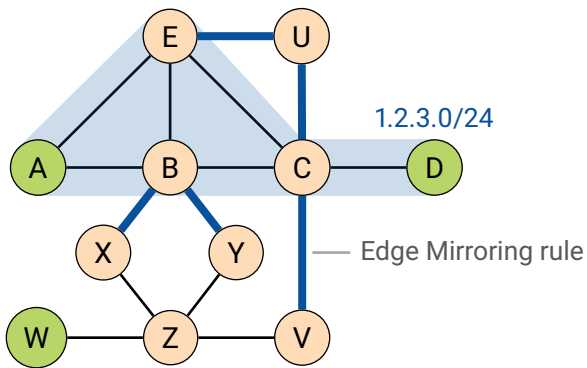
MIRROR



CONFINE

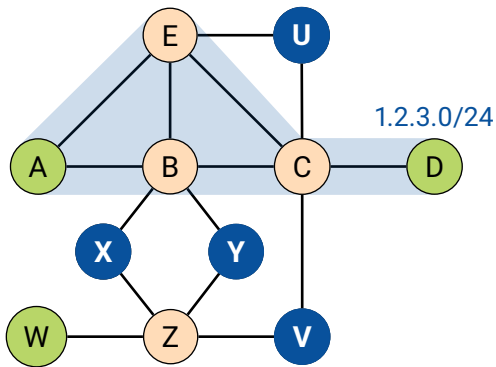
CONFINE queries mirror packets leaving a confinement region

CONFINE 1.2.3.0/24 ON [A B E C D]



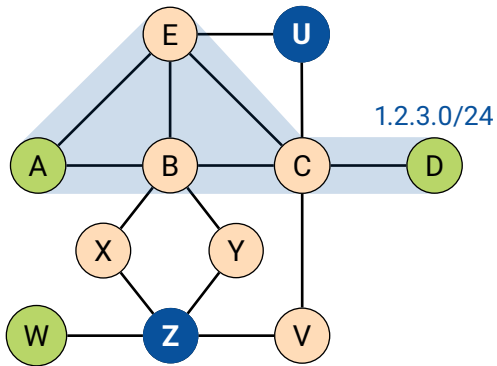
Fewer mirroring rules minimizes control-plane overhead

CONFINE 1.2.3.0/24 ON [A B E C D]



The lower bound is a multi-terminal node cut

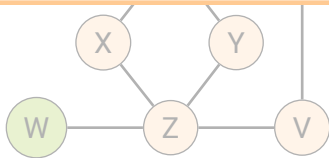
CONFINE 1.2.3.0/24 ON [A B E C D]



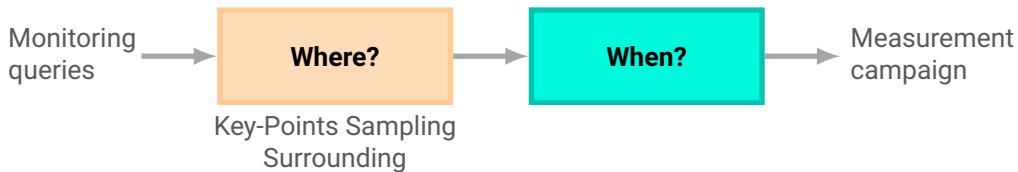
The lower bound is a multi-terminal node cut

CONFINE 1.2.3.0/24 ON [A B E C D]

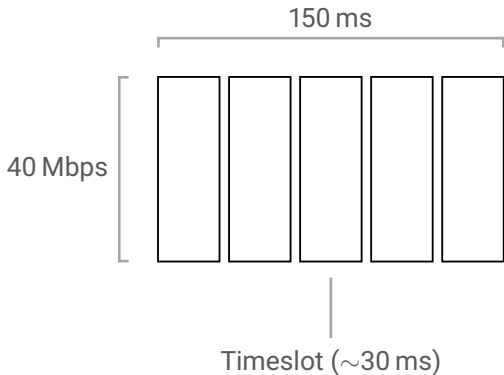
The **Surrounding** algorithm minimizes mirroring rules and guarantees to mirror any packet leaving the confinement region



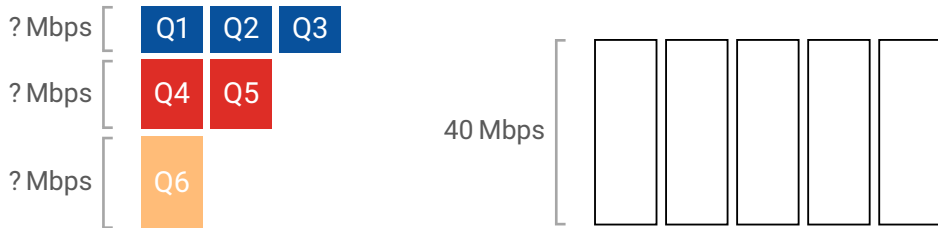
Query activations must be scheduled to meet the budget



Stroboscope divides the monitoring budget in timeslots



Stroboscope requires traffic demand estimations



Stroboscope conservatively estimates traffic demands

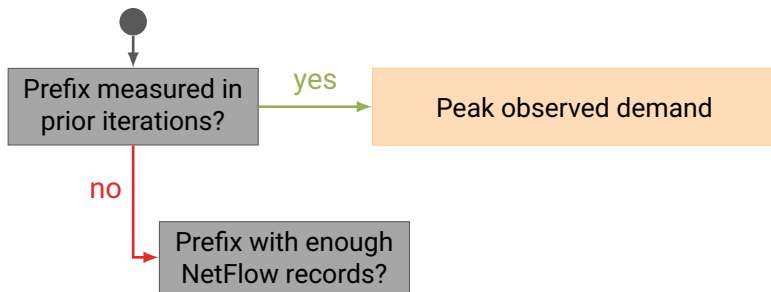


Prefix measured in
prior iterations?

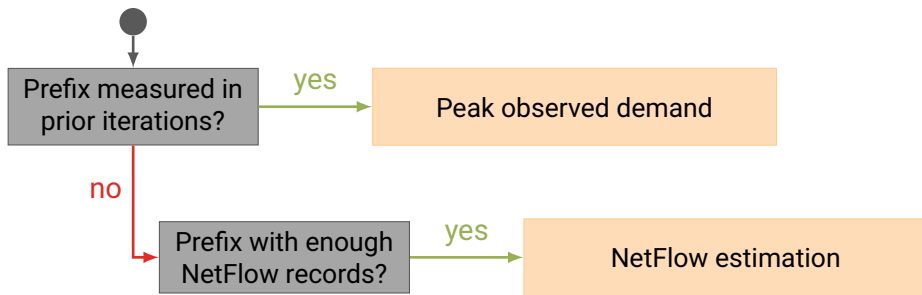
Stroboscope conservatively estimates traffic demands



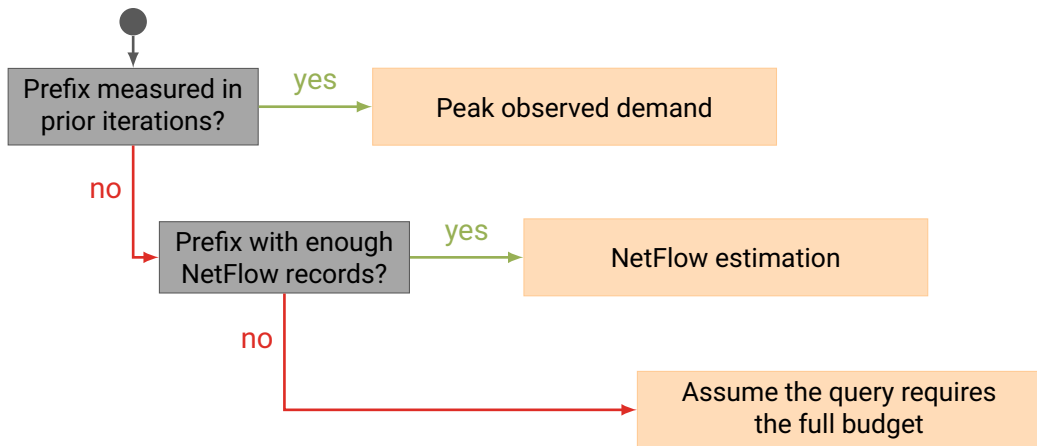
Stroboscope conservatively estimates traffic demands



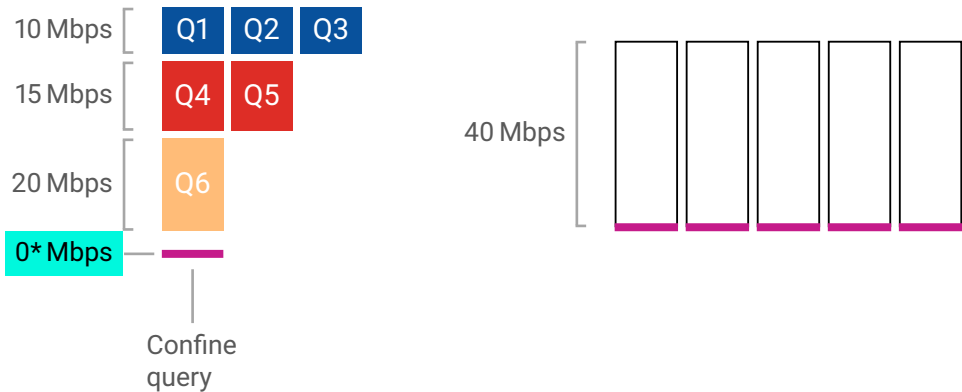
Stroboscope conservatively estimates traffic demands



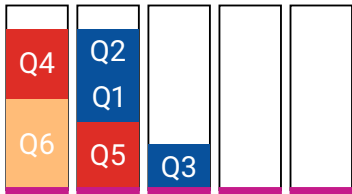
Stroboscope conservatively estimates traffic demands



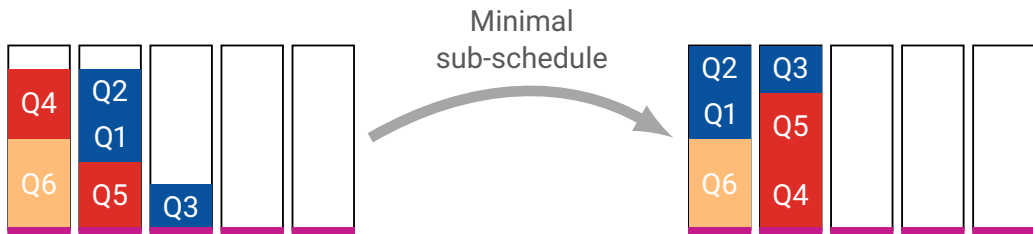
Confine queries are scheduled in all timeslots



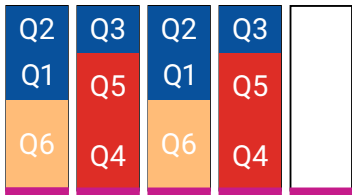
Stroboscope first approximates a minimal sub-schedule



Stroboscope first approximates a minimal sub-schedule, optionally optimizing for the optimal bin-packing solution

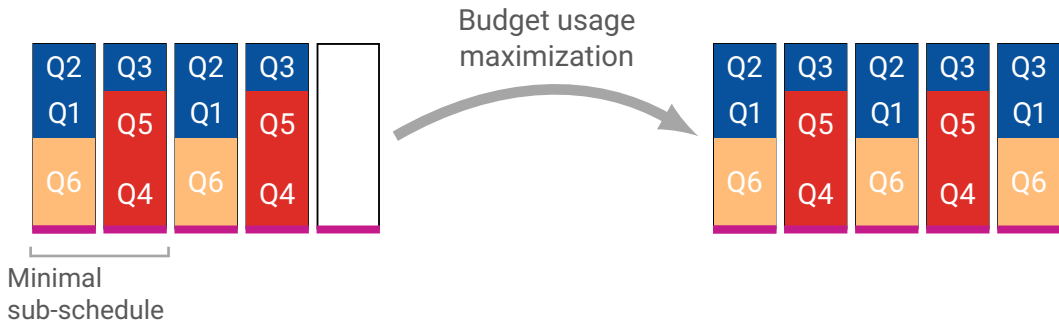


Stroboscope replicates the sub-schedule



Minimal
sub-schedule

Stroboscope replicates the sub-schedule, and minimizes budget leftovers



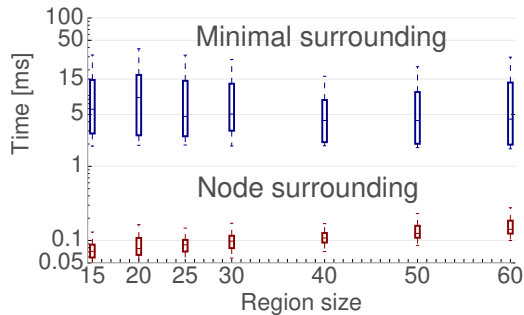
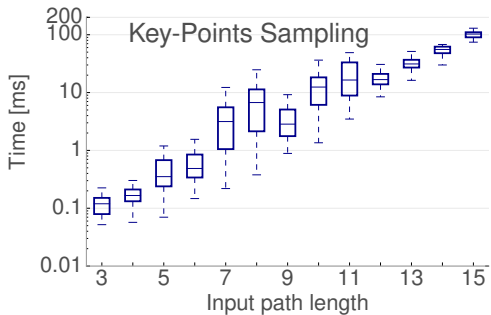
Stroboscope achieves deterministic sampling

Stroboscope: Declarative Network Monitoring on a Budget

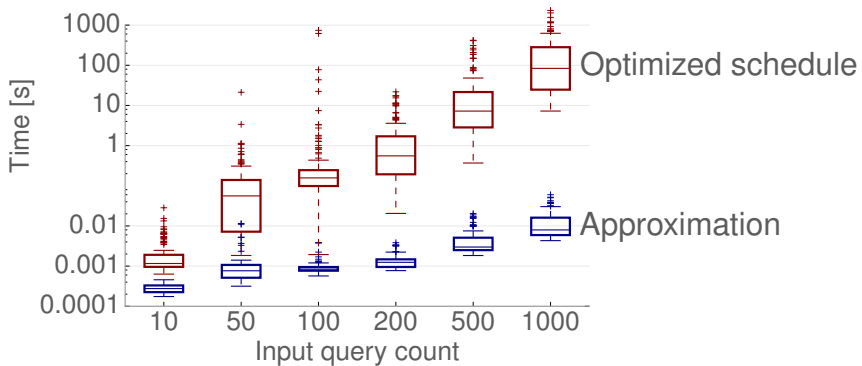


- Collecting traffic slices to monitor networks
- Adhering to a monitoring budget
- Using Stroboscope today

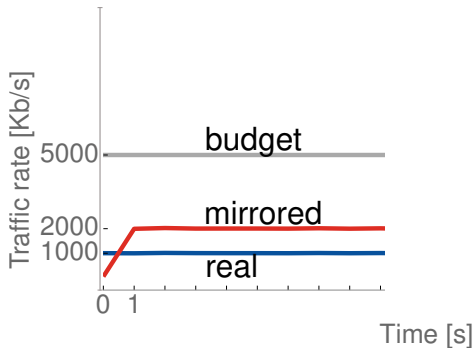
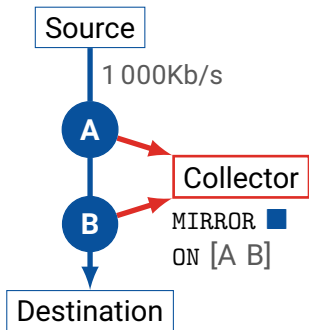
Selecting mirroring locations in realistic ISP topologies is fast



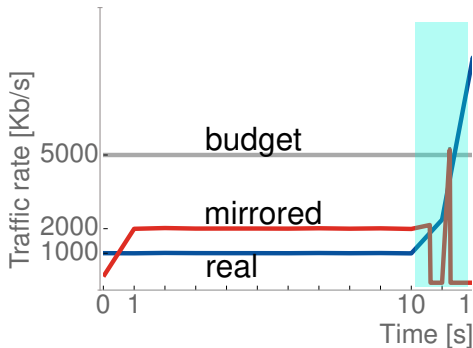
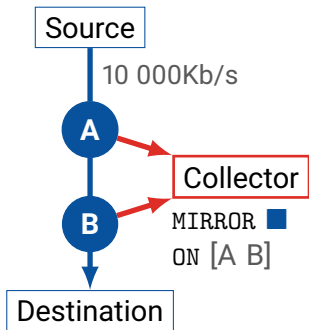
Schedules can be quickly approximated



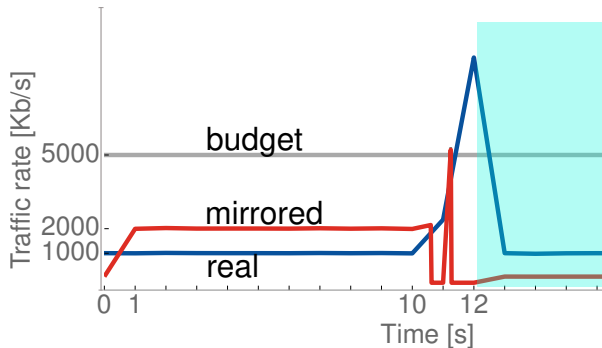
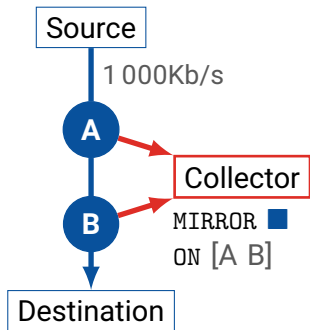
Stroboscope tracks the rate of mirrored traffic in real time



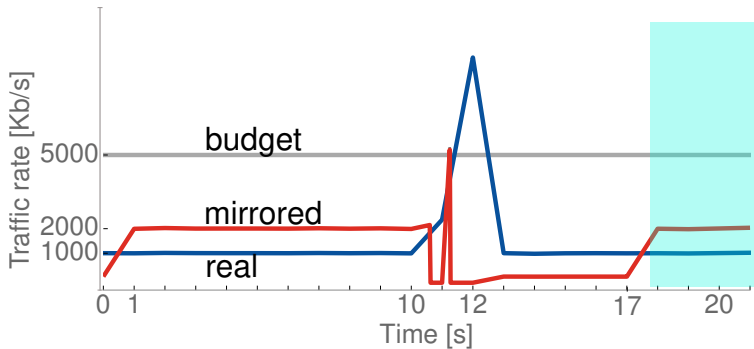
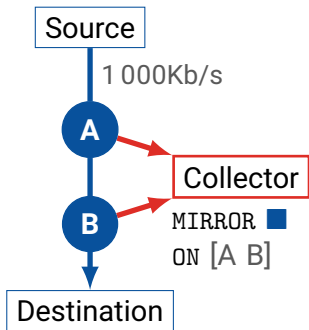
Measurement campaigns are stopped early if the estimated demand are exceeded



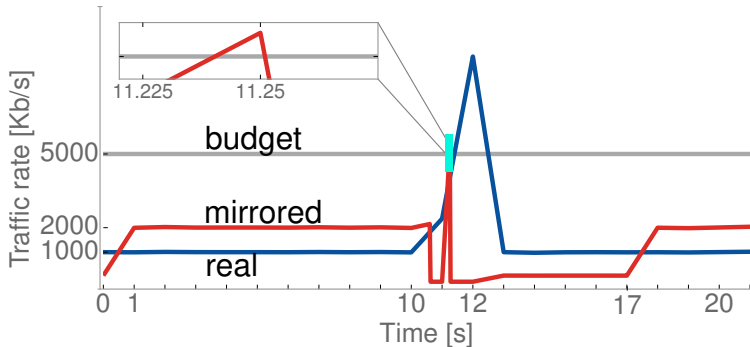
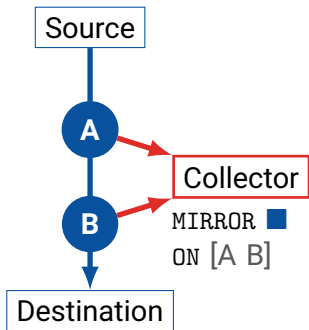
Exceeding the total budget schedules the query once per measurement campaign



Stable recorded traffic rates are used for future estimations



Stroboscope exceeds the monitoring budget for at most one timeslot



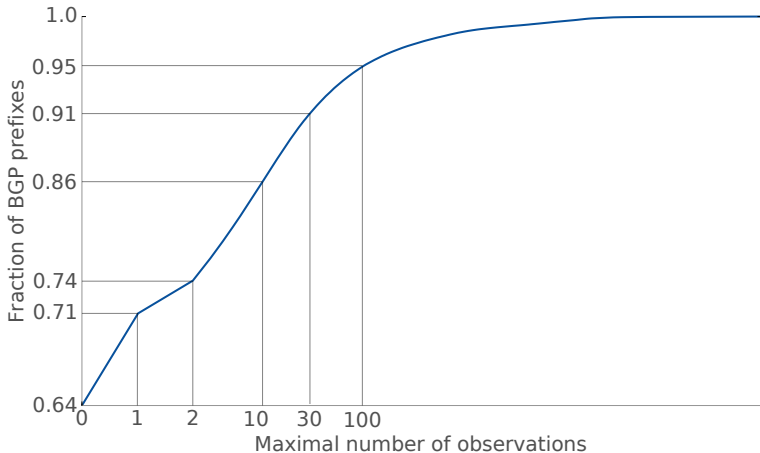
Stroboscope: Declarative Network Monitoring on a Budget



- Traffic slicing as a first-class data-plane primitive
- Strong guarantees on budget compliance and measurement accuracy
- Measurement analysis decoupled from measurement collection

Backup slides

NetFlow brings a poor visibility over traffic in ISP networks



Stroboscope defines a declarative requirement language

MIRROR 1.2.3.0/24 ON [A B C D], [A E C D]

MIRROR 1.2.3.0/24 ON [A -> D]

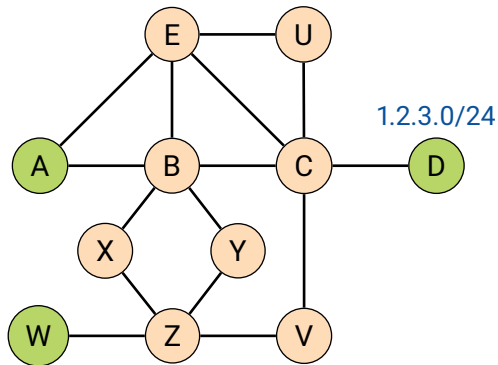
CONFINE 1.2.3.0/24 ON [A B E C D]

CONFINE 1.2.3.0/24 [A -> D]

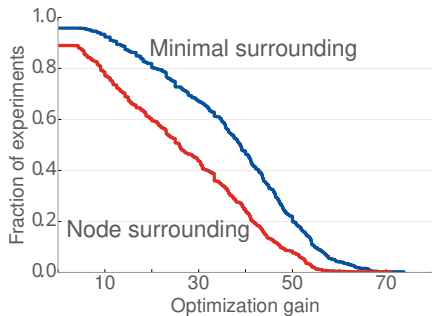
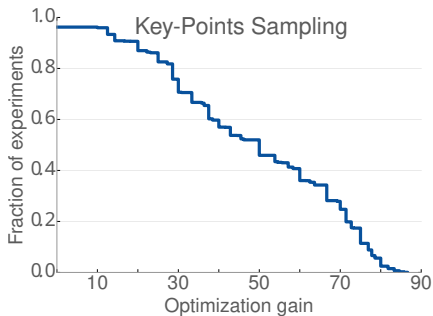
MIRROR 1.2.3.0/24 ON [-> D]

CONFINE 1.2.3.0/24 ON [-> D]

USING 15 Mbps DURING 500 ms EVERY 5 s



The placement algorithms minimize the mirroring rules



Schedules can be computed by two pipelines

