ALEMBIC: AUTOMATED MODEL INFERENCE FOR Stateful Network Functions

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Stateful Network Functions (NFs) in Modern Networks



Modern networks contain a wide range of complex stateful network functions from many vendors





Connection Map:





Connection Map:







Connection Map:















Network Testing and Verification



Operator



 Is the policy implemented correctly? Can we check before on-boarding?

We need network testing/verification tools (e.g.,VMN, SYMNET, BUZZ...)







Today, these NF models are handwritten based on manual investigation

Testing Verification **On-boarding**

Handwritten FW model

Network testing tool e.g., BUZZ [NSDI 16]

Error!

≠

Handwritten FW model

Network testing tool e.g., BUZZ [NSDI 16]

≠

Handwritten FW model

Real FW implementation

Network testing tool e.g., BUZZ [NSDI 16]

Handwritten FW model

Network testing tool e.g., BUZZ [NSDI 16]

Handwritten FW model

Network testing tool e.g., BUZZ [NSDI 16]

Limitation of Handwritten Model: Vendor Diversity

Vendor-specific differences

Vendors have different implementations!

Our Work: Alembic

Automatically infer a behavioral model of the NF for a configuration

Customers: **1)** BUZZ [NSDI16] 2) SYMNET [SIGCOMM16] 3) VMN [NSDI17]

Motivation and Goal

Challenges and Insights

Overall Workflow

Evaluation

High-Level Challenges

Large configuration space

Inferring NF behavior

Challenges on Large Configuration Space

Configuration *→* many rules

• Rule \rightarrow IP/port fields take large sets of values (e.g., 2³² for IPs)

• Rule \rightarrow IP/port fields can be **ranges** (e.g., /16 for IP prefixes)

Insight 1: We Can Compose Models of Individual Rules

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"compose" per rule models

Challenges on Large Configuration Space

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Challenges on Large Configuration Space

Configuration → many rules

• Rule \rightarrow IP/port fields take large sets of values (e.g., 2³² for IPs)

Rule → IP/port fields can be ranges (e.g., /16 for IP prefixes)

Insight 3: Exploit Independence to Create an Ensemble of FSMs

SRC IP:10.1.1.0/16...DST IP:15.1.1.0/16

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SRC IP:10.1.1.0/16...DST IP:15.1.1.0/16 Per-connection

Independent packet processing per connection

Conn 1 : 10.1.1.1 \rightarrow 15.1.1.1 States do not interfere Conn 2 : 10.1.1.2 \rightarrow 15.1.1.2

Insight 3: Exploit Independence to Create an Ensemble of FSMs

SRC IP:10.1.1.0/16...DST IP:15.1.1.0/16

Independent packet processing per connection

Instantiate

at runtime

(symbolic model from insight 2)

Learn

M(A, B)

Per-connection

Ensemble of FSMs

An ensemble of concrete FSMs can represent a rule with IP/port ranges

Summary of Insights to Address Large Configuration Space

A configuration is composed of many number of rules

Symbolic Model

Compositional Model

A rule contains **IP/port fields** which take **large sets of values** and **ranges**.

Back to High-Level Challenges

Large configuration space

Inferring NF behavior

Challenges on Inferring NF Behavior

Inferring the symbolic FSM

Inferring the state granularity

Handling dynamic header modification

We can use the L* algorithm!

FSM representing the blackbox

- Generates sequences (e.g., aa, aba) and probes the blackbox Builds a hypothesis FSM with input-output pairs seen so far • Queries an Equivalence Oracle (EO) for counterexamples

Practical Challenges of Applying L* for an NF

Generate input alphabet

Classify output of an NF

• Build an Equivalence Oracle

Naive solutions:

1. Exhaustively generating packets

Infeasible

2. Randomly generating packets **Does not explore the relevant state space**

To exercise the rule, we generate packets with IP/ports in the rule

Find IP/port fields that appear in the rule Generate the packet for for all interfaces using A and B

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To exercise the rule, we generate packets with IP/ports in the rule

- 1) Find IP/port fields that appear in the rule Generate the packet for for all interfaces using A and B
- 2) (Optional) Prune based on reachability
- 3) Plug in "packet types"

Practical Challenges of Applying L* for an NF

- Generate input alphabet
- Classify output of an NF
 - Configure the "timeout" to classify output
 - Translating to/from symbolic and concrete packets
- Build an Equivalence Oracle

Challenges on Inferring NF Behavior

Inferring the symbolic model (FSM)

Inferring the state granularity

Handling dynamic header modification

Different Types of State Granularity

State Granularity: the state variables (IP/ports) that the NF uses to keep state

Cross-connection

Per-source

Per-destination

Per-connection

This is like a "key" mapping to the FSM

	One FSM for all connections
	One FSM for each srcip
	One FSM for each dstip
	One FSM for every IP-port pair
"	nning to the ECM

Learning the State Granularity

Do these affect the "same" FSM?

Cross-connection

Do these affect the "same" FSM?

Per-source

Learning the State Granularity

Do these affect the "same" FSM?

Cross-connection

Do these affect the "same" FSM?

Per-source

Construct test cases for independence across connections

Alembic Workflow: Offline

Runs once per NF

Alembic Workflow: Online

Runs for every config

Evaluation Summary

- Alembic-generated models are accurate
- Case Studies: Alembic finds differences across NF implementations
- Alembic workflow is scalable
- Alembic-generated models improve the accuracy of network testing/verification tools

Evaluation Setup

- **Real NFs** we modeled :
 - PfSense (FW, static NAT, random NAT, LB)
 - Proprietary NF (FW, static NAT)
 - Untangle (FW)
 - HAproxy (LB)
- Packet types used:
 - Correct-Seq: {SYN_c, SYN-ACK_c, ACK_c, FIN-ACK_c, RST-ACK_c}
 - $\{SYN-ACK_{I}, ACK_{I}, FIN-ACK_{I}, RST-ACK_{I}\}$

Validated Alembic using Click-based NFs where we know the ground truth

Combined-Seq: extend the correct-seq set with incorrect seq and ack,

testing methodology to test the accuracy of our models

- Config generation: 1 to 100 rules in a configuration Packet generation: 20 to 300 packets in a sequence
- 1) **Iperf testing:** 100% across all settings for all NFs 2) Random Packet testing (randomly choosing IP/port):
 - 99.8% to 100% across all settings for all NFs
- 3) Rule Activation testing (choosing IP/port to activate one rule): 94.8% to 100% across all settings for all NFs

Accuracy Evaluation

Since we do not have the ground-truth, we designed complementary

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Packet sequence before the FW allows TCP traffic from an external host (B) to an internal host (A)

SYN, A→B

Firewall Case Study

Packet sequence before the FW allows TCP traffic from an external host (B) to an internal host (A)

SYN, A→B

Number of states

Firewall Case Study

Packet sequence before the FW allows TCP traffic from an external host (B) to an internal host (A)

Number of states

Default behavior

SYN, A→B

Firewall Case Study

- Implements "default allow"
- Connection-terminating

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the FW preemptively responds with ACK

- Implements "default allow"
- Connection-terminating

When A replies with ACK, the FW drops to prevent duplicates

- Implements "default allow"
- Connection-terminating

Takeaways: Vendor diversity (no common practice) The real FSMs are complex and are infeasible for humans to manually generate

• **FW:** models with incorrect seq \rightarrow large FSM (257 states for PfSense)

FW: many do not correctly handle out-of-window packets

• LB: HAproxy (connection-terminating) vs. PfSense (destination NAT)

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Number of Rule 10 100 1,000

Scalability of Alembic Online

S	Runtime
	0.075 s
	0.6 s
	5 s

Alembic can denerate concrete models in a few seconds for a large config

Limitations and Future Work

Assumption on configurations:

- Assume at most one rule is applied
- States across different state granularities (i.e., keys) are independent
- Assume that IP/port fields are treated homogeneously such that we can pick one representative sample and infer a model

Assumption on NF actions:

- Focused on modeling TCP-relevant behavior where actions are restricted to dropping and forwarding, possibly with IP/port modifications
- Do not explicitly model temporal effects
- Support the following state granularity types: per-connection, per-source, perdestination, cross-connection, and stateless

Future work:

Dealing with more complex NFs (e.g., rate-limiting NF, modeling temporal effects)

Conclusions: Alembic can accurately model stateful NFs

- Network testing and verification today need NF models
- Handwritten models: tedious, error-prone, and inaccurate
- Alembic: infers a high-fidelity NF model given a configuration
- Our evaluations show:
 - Alembic finds implementation-specific behavior of NFs
 - Alembic-generated models increase the accuracy of testing/ verification
 - Alembic is scalable and accurate

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