

Automatic Generation of Data-Oriented Exploits

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Control Flow Attacks Are Getting Harder

- State-of-the-art exploits
 - Code injection
 - heap spray / JIT spray

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 - Defenses
 - Data Execution Prevention
 - Control Flow Integrity
- control-flow bending

- Stat-of-the-art exploits from memory errors
 - Code injection (e.g., heap spray / JIT spray)
 - Code reuse (e.g., ret2libc, ROP)
- Defenses
 - DEP, CFI, ASLR
 - Block control flow hijacking in principle

CONTROL PLANE



Data-Oriented Exploits

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setuid(0);
.....
// set normal user privilege
setuid(pw->pw_uid);
// execute user's command
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Wu-ftpd *setuid* operation*

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Data-Oriented Exploits

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// set root privilege
seteuid(0);
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// set normal user privilege
seteuid(pw->pw_uid);
// execute user's command
```

Wu-ftpd *setuid* operation*

```
//0x1D4, 0x1E4 or 0x1F4 in JScript 9,
//0x188 or 0x184 in JScript 5.8,
safemode = *(DWORD*)(jsobj + 0x188);
if(safemode & 0xB == 0) {
    Turn_on_God_Mode();
}
```

IE *SafeMode* Bypass+

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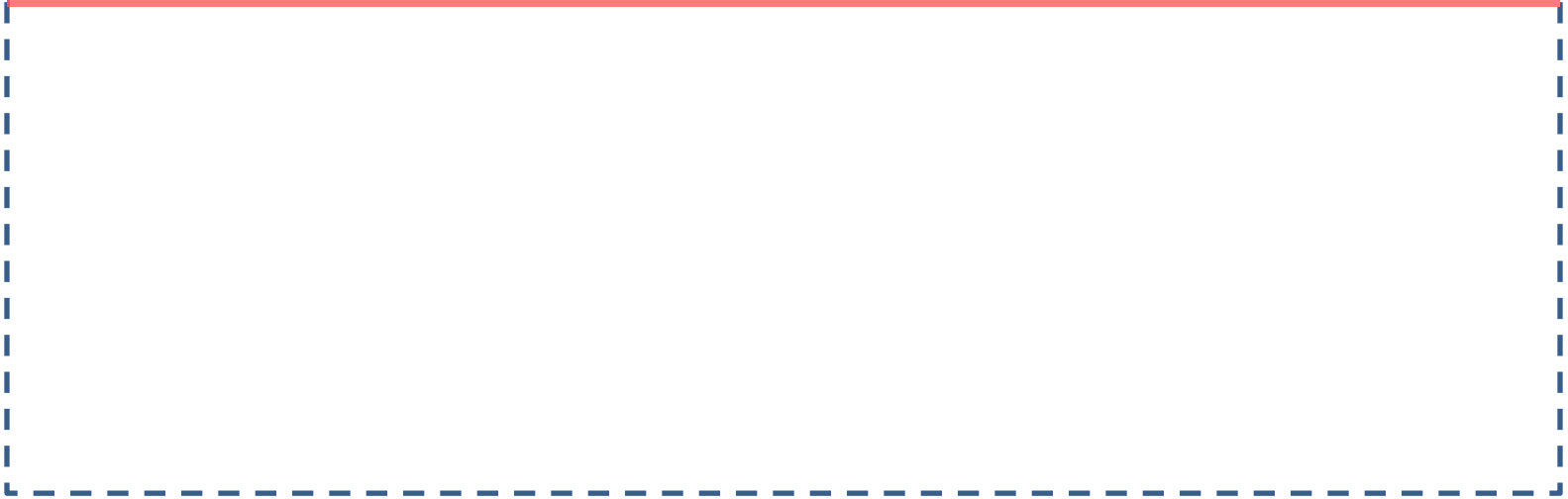
+ Yang Yu. Write Once, Pwn Anywhere. In Black Hat USA 2014

Data-Oriented Exploits

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- New class of Data-Oriented Exploits
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- Data Flow Stitching
 - Systematic search for data-oriented exploits
 - Works on binary directly

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- Data Flow Stitching
 - Systematic search for data-oriented exploits
 - Works on binary directly
- Results
 - Concrete exploits on real web/file servers
 - 19 exploits (16 new) from 8 vulnerabilities

Motivating Example

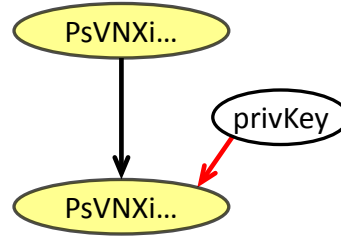
- SSL-enabled web server

```
1 int server() {
2   char *userInput, *fileName;
3   char *privKey, *result, output[BUFSIZE];
4   char fullPath[BUFSIZE]="/path/to/root/";
5
6   privKey=loadPrivKey("/path/to/privKey");
7   GetConnection(privKey, ...);
8   userInput = read_socket();
9   if (checkInput(userInput)) {
10    fileName = getFileName(userInput);
11    strcat(fullPath, fileName);
12    result = retrieve(fullPath);
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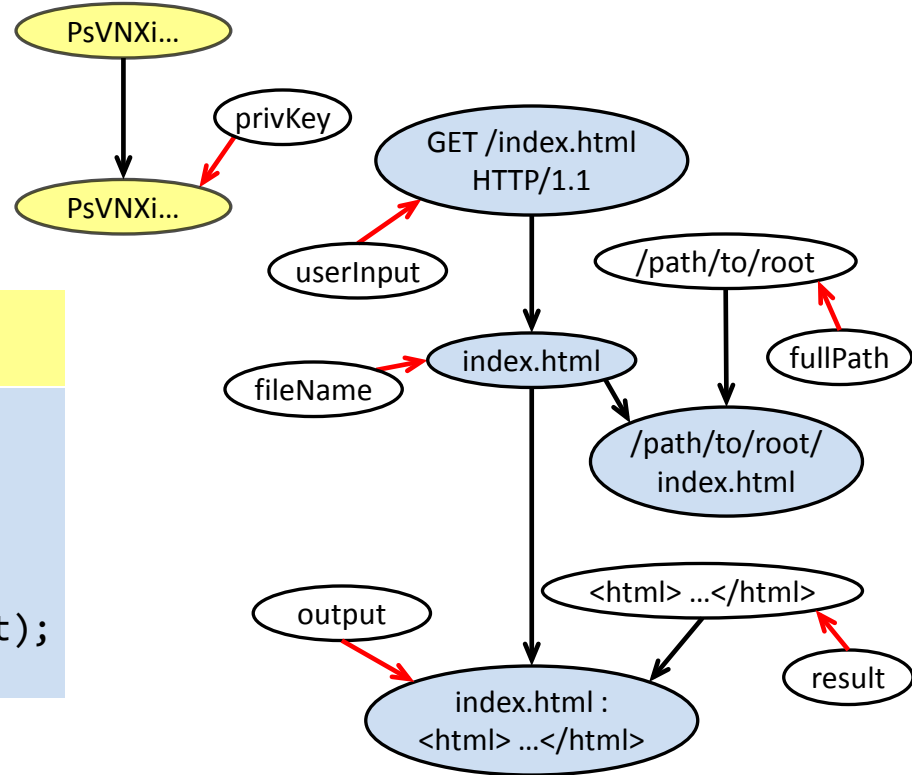
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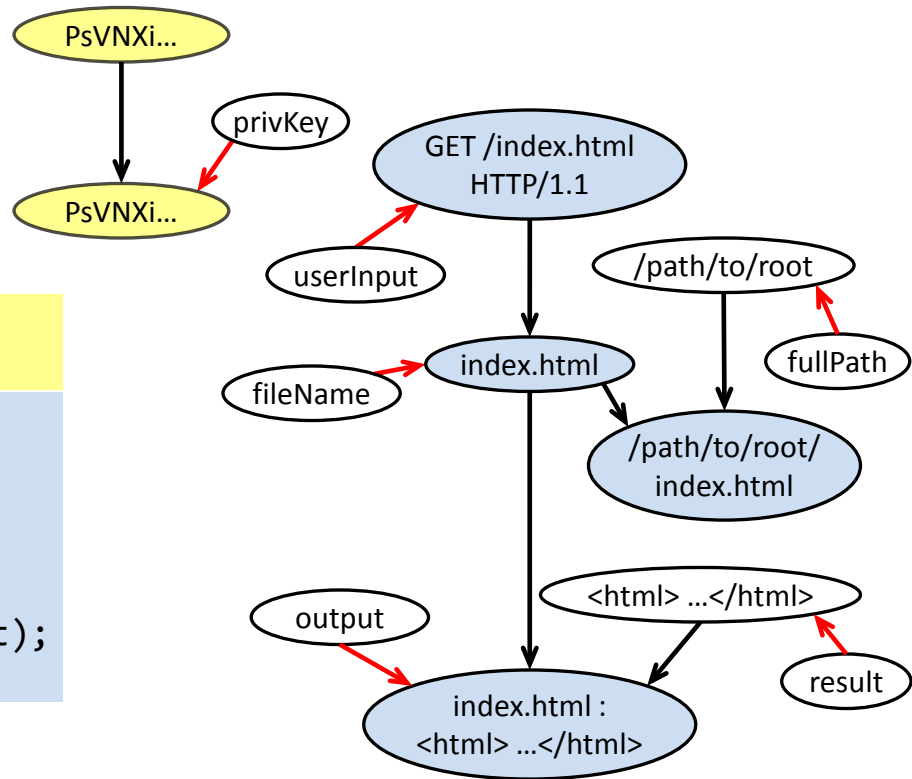
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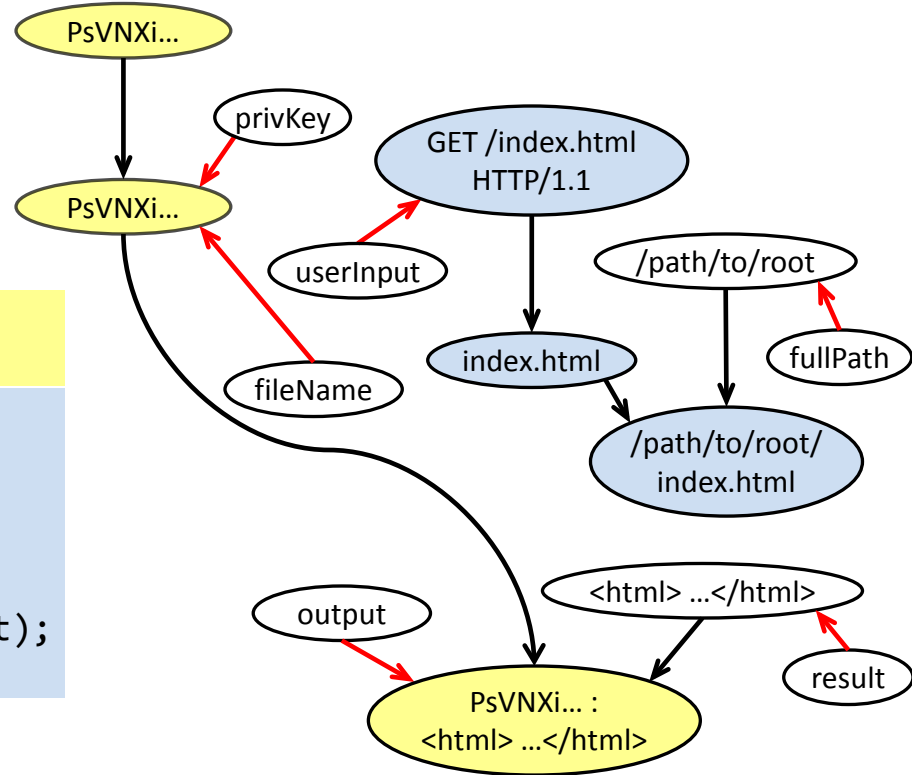
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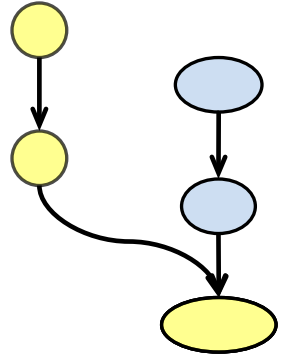
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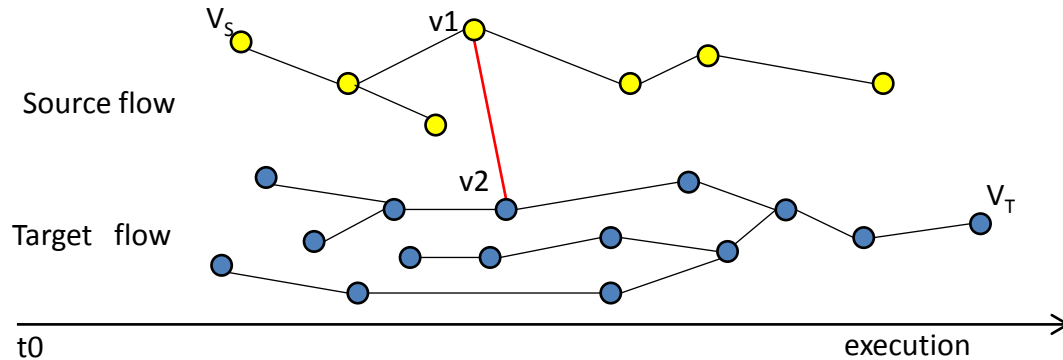
Data-Flow Stitching

- Manipulate data flows for exploits
- Enables systematic way to search for exploits
 - Input: binary & error-exhibiting input
 - Output: data-oriented exploits
- Goal:
 - Information Leakage (e.g., password, keys)
 - Privilege Escalation (e.g., setuid, access priv. files)
- Constraints:
 - Keep the control-flow same
 - Prevent abrupt termination
 - No knowledge of randomized values (CFI tags, ASLR addresses)



Challenges

- Time-consuming search
 - The search-space: Cartesian product $|\text{SrcFlow}| \times |\text{TgtFlow}|$
 - Heavy analysis for each candidate



- Our solution:
 - Filter out candidates with memory error influence
 - Use an SMT solver to verify candidates

Single-Edge Stitch

- Corrupt data vertex

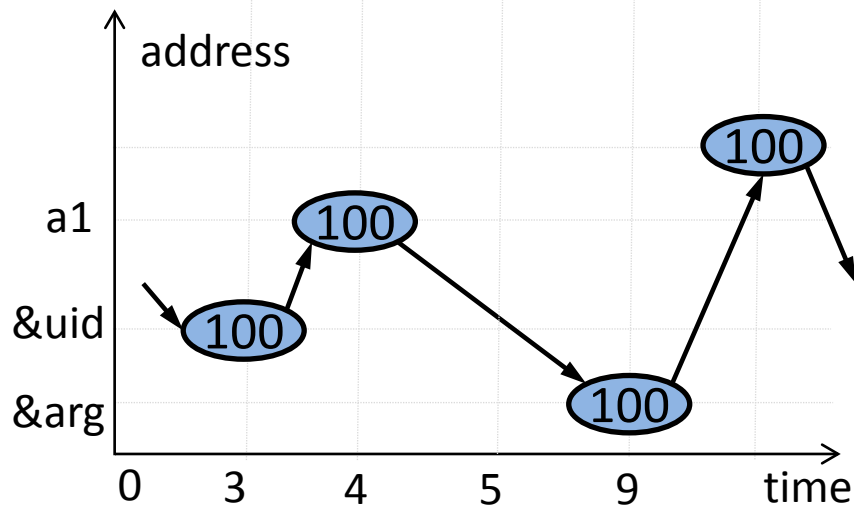
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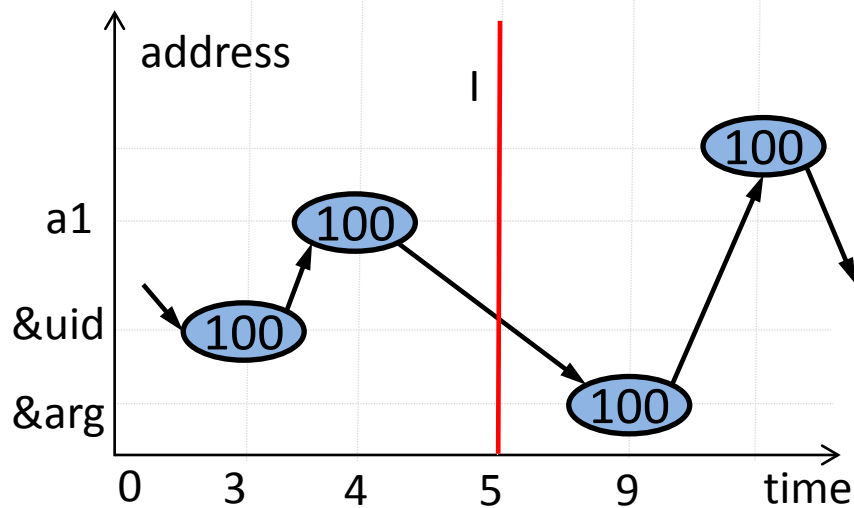
- 2D-DFG



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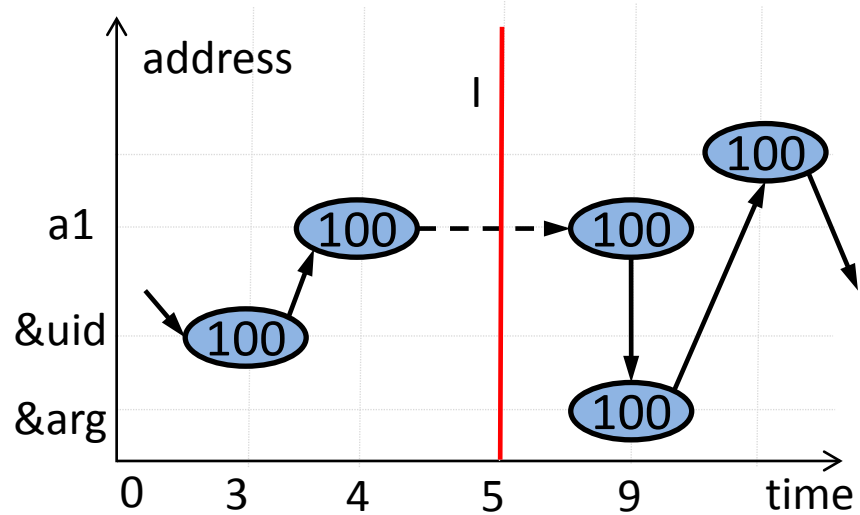
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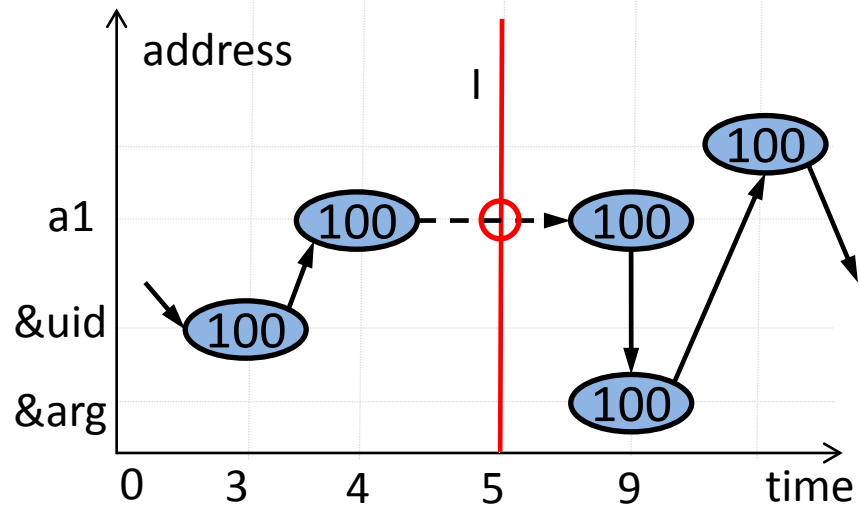
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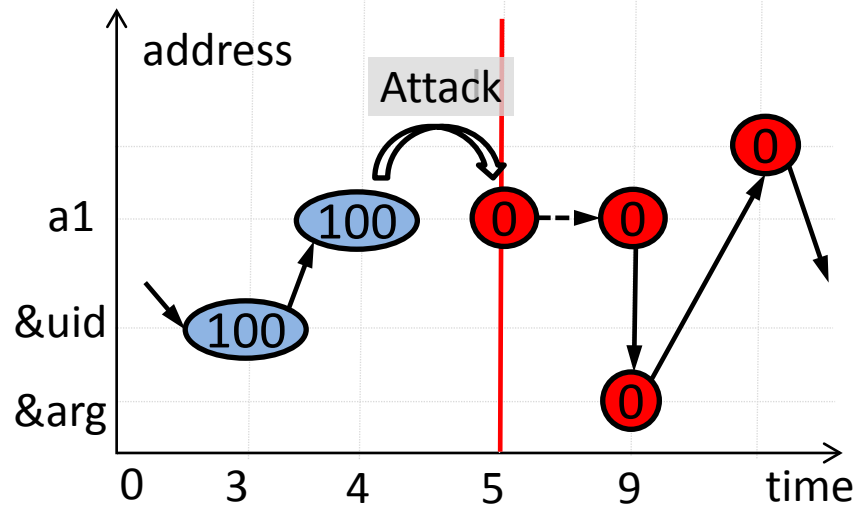
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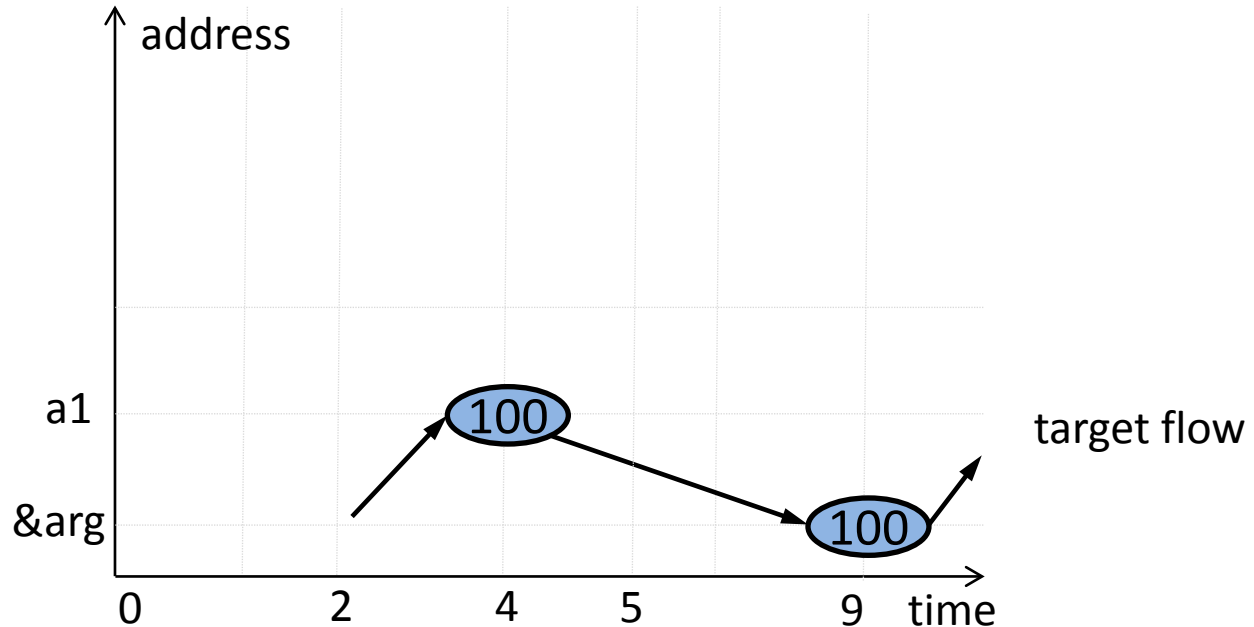
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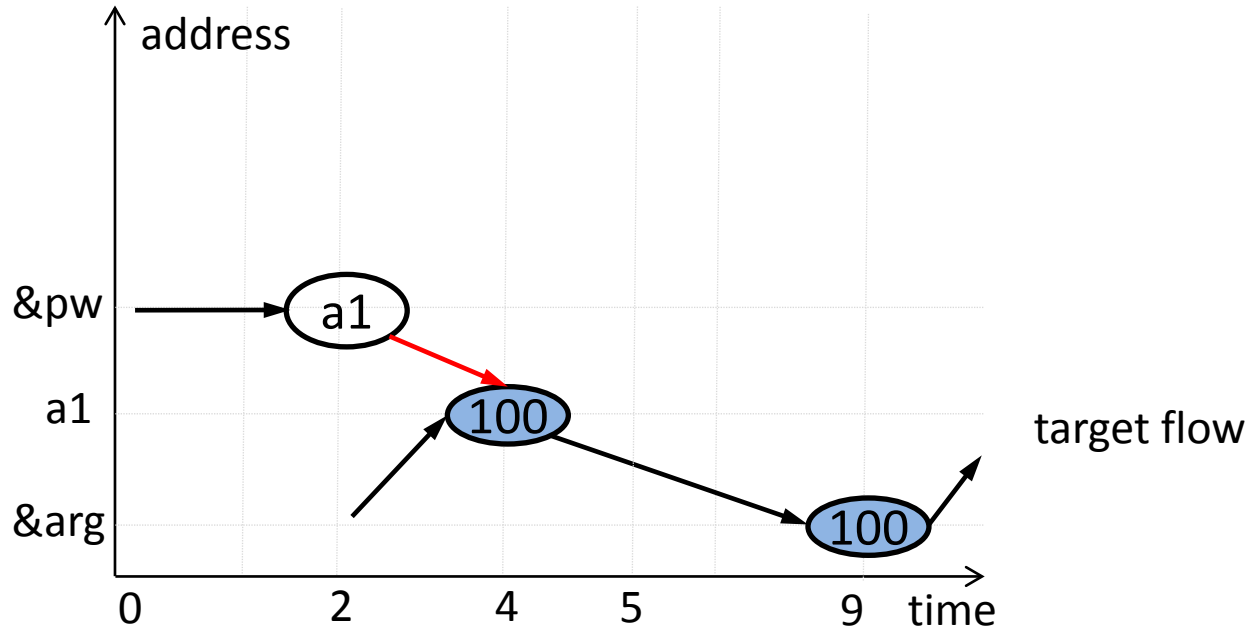
Pointer Stitch

- Corrupt pointers to connect data flows
 - Pointers decide data movement direction



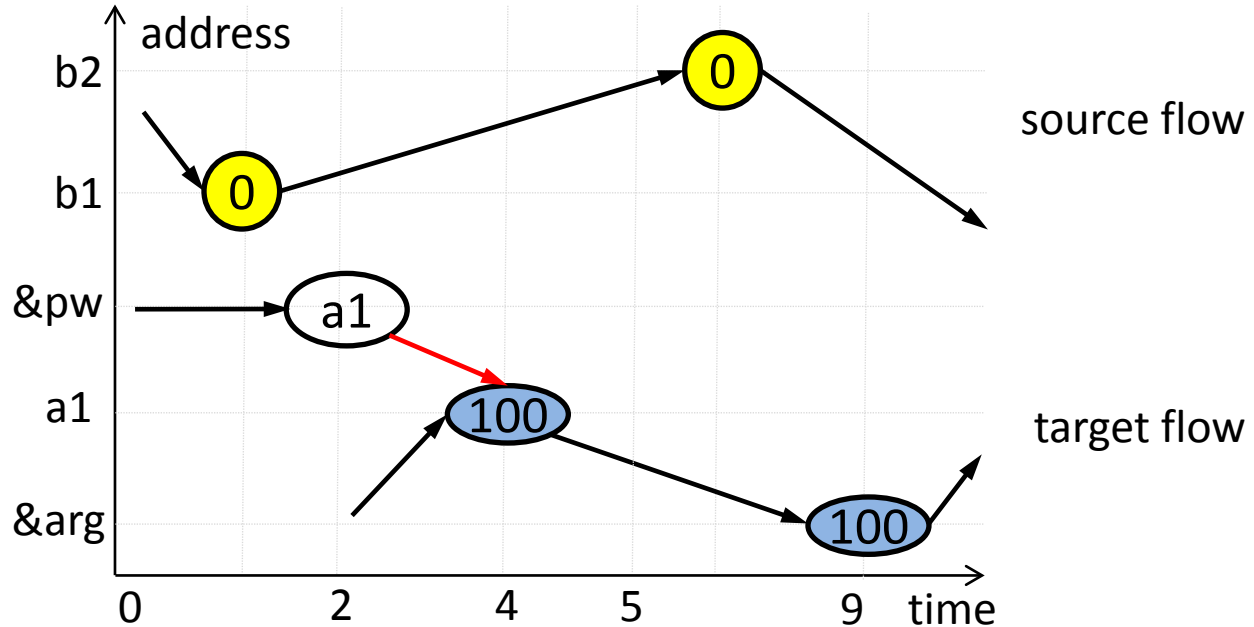
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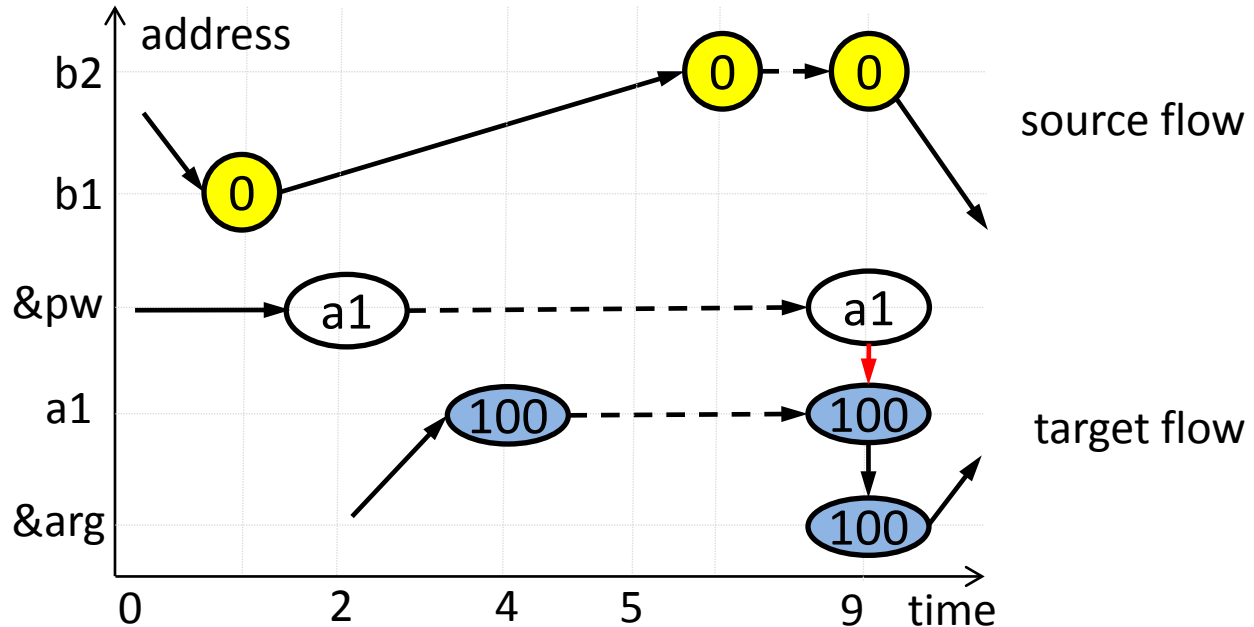
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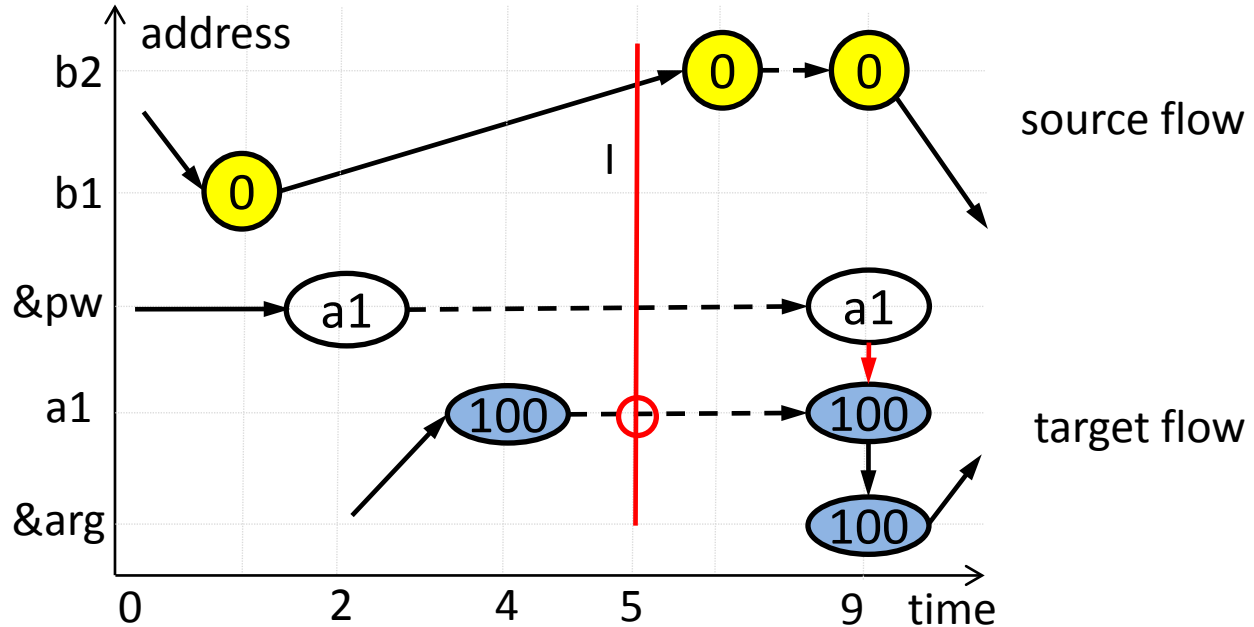
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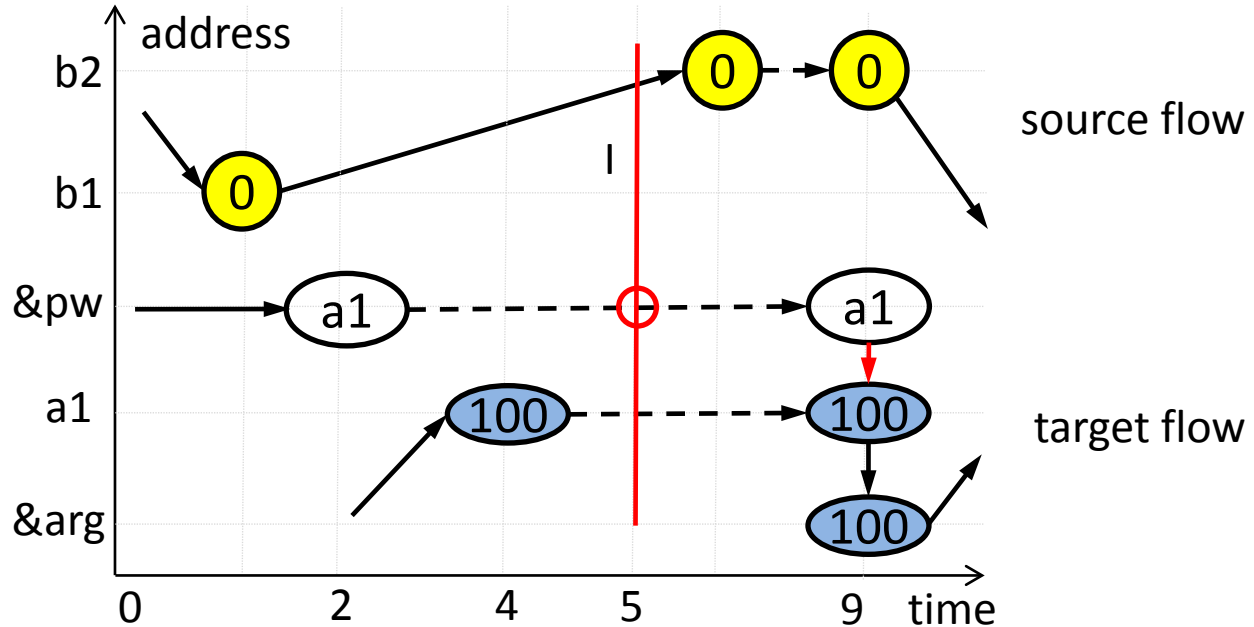
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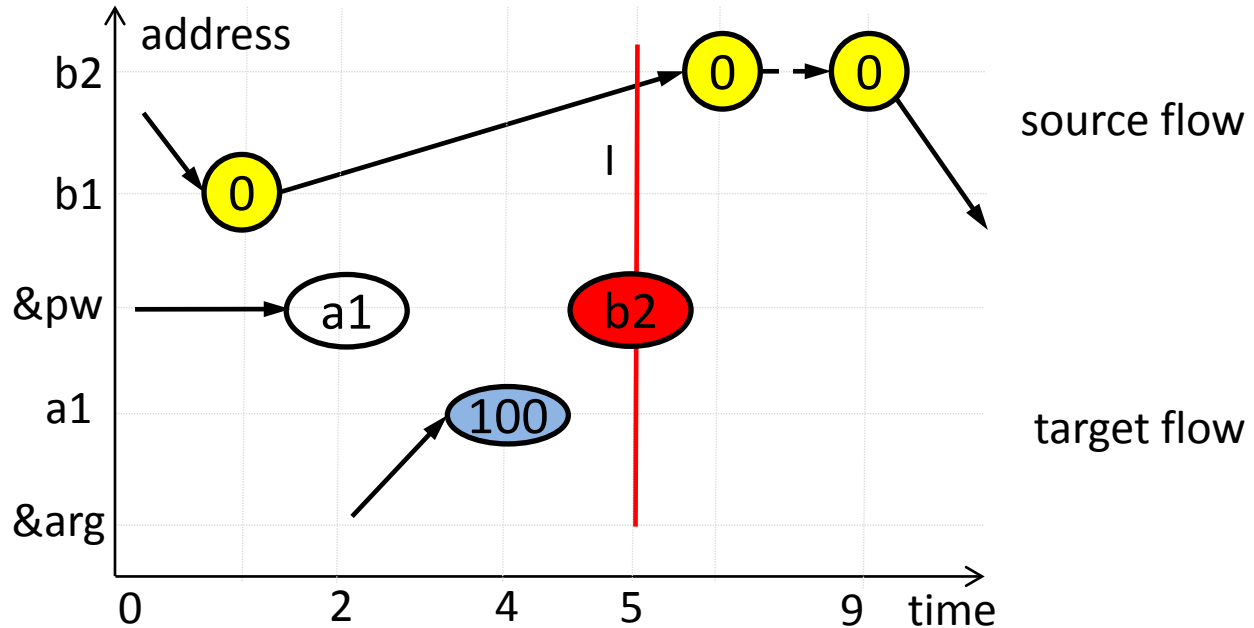
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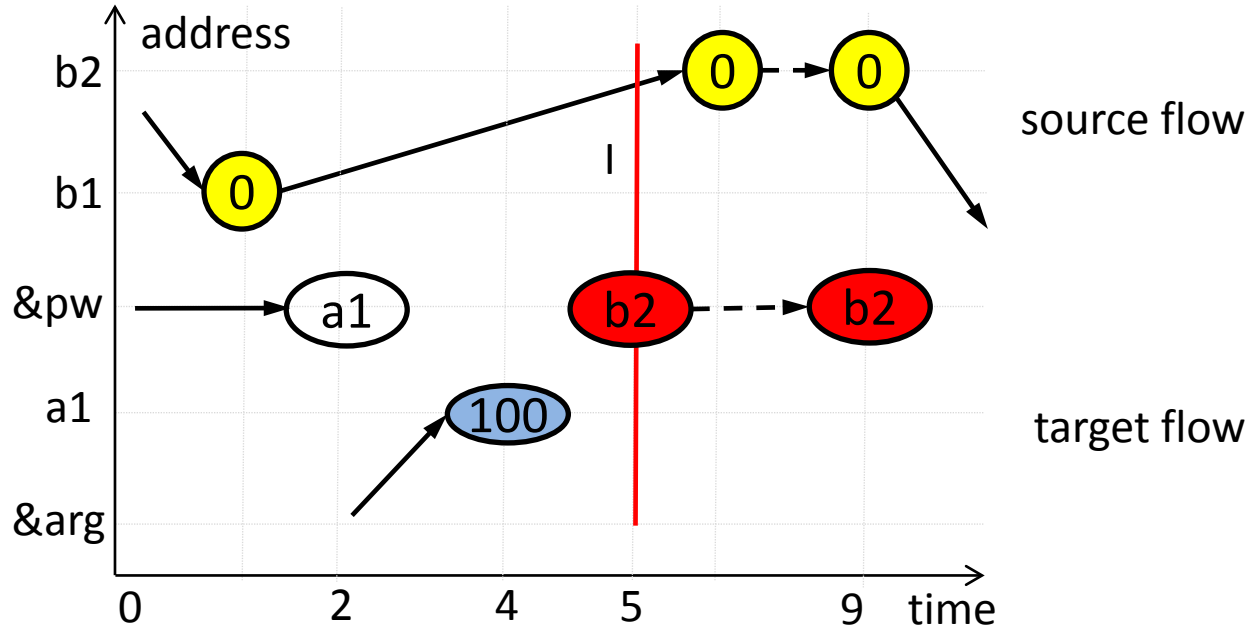
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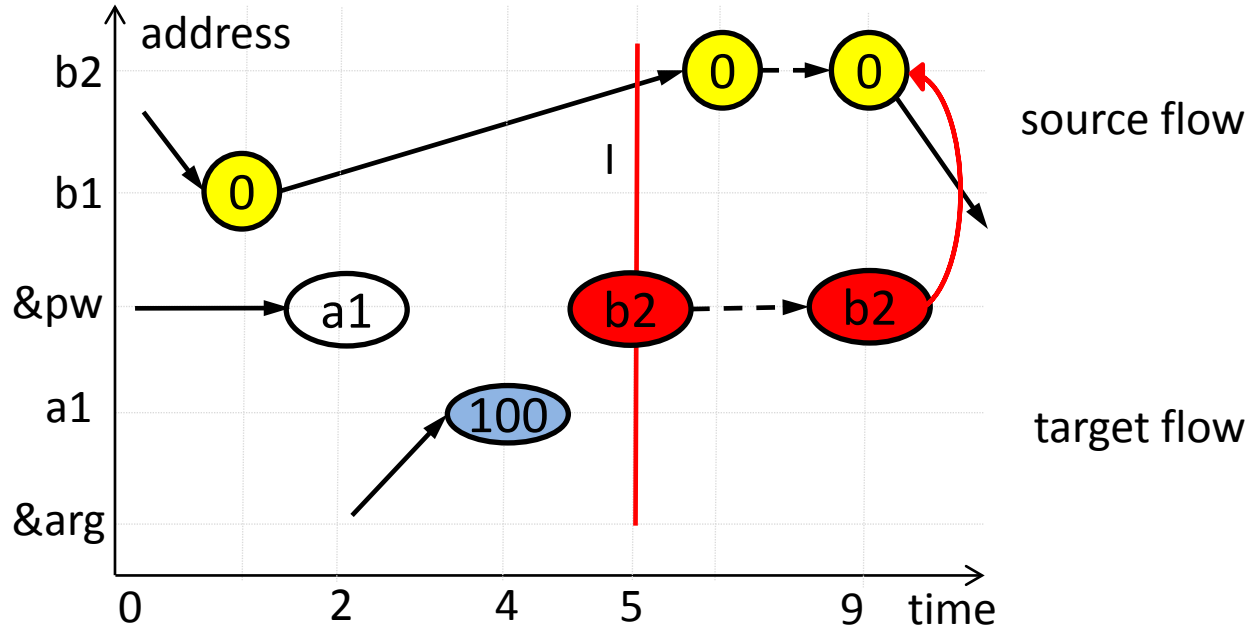
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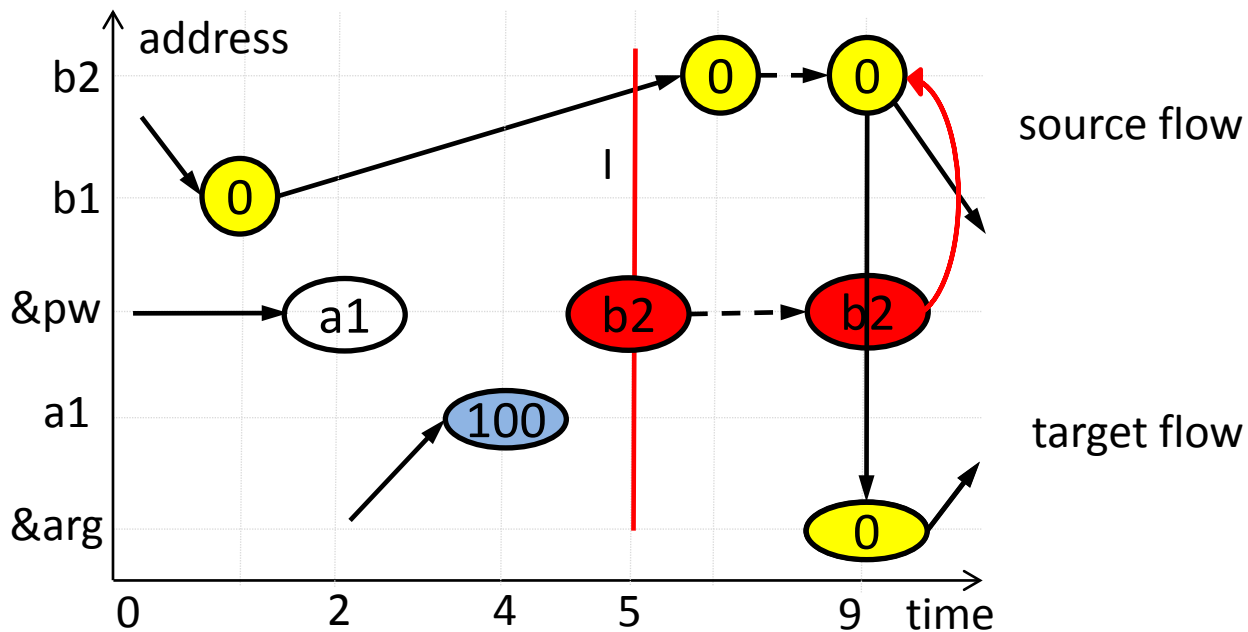
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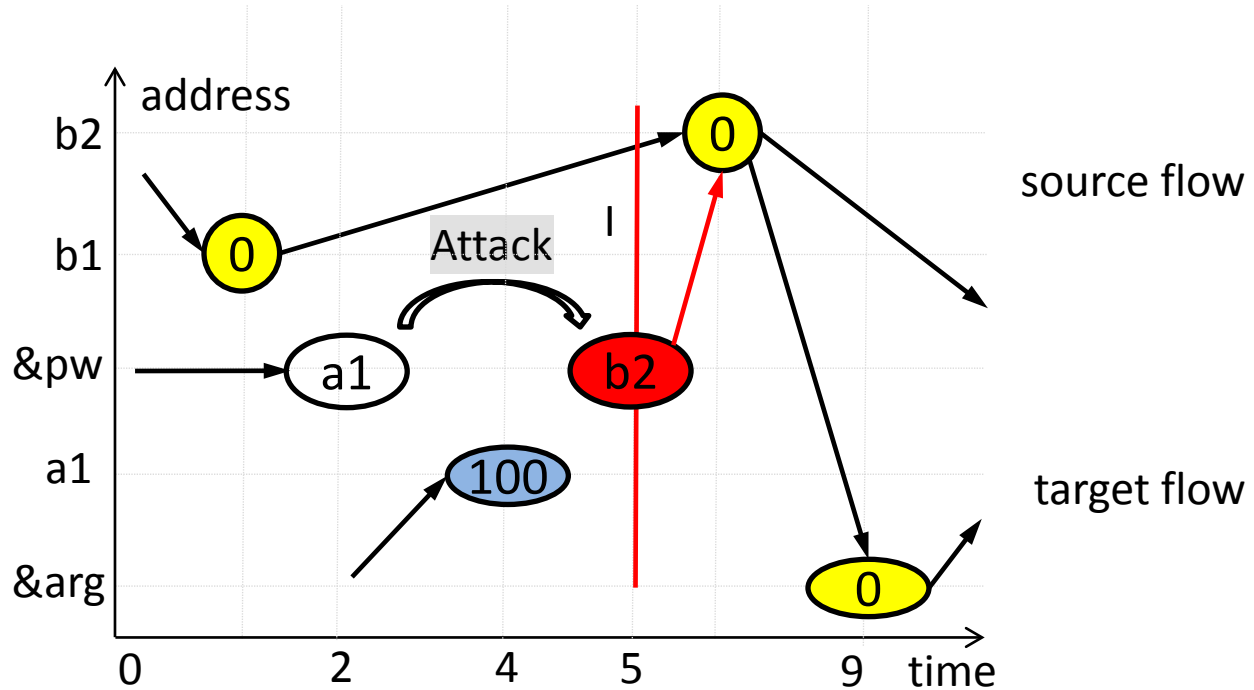
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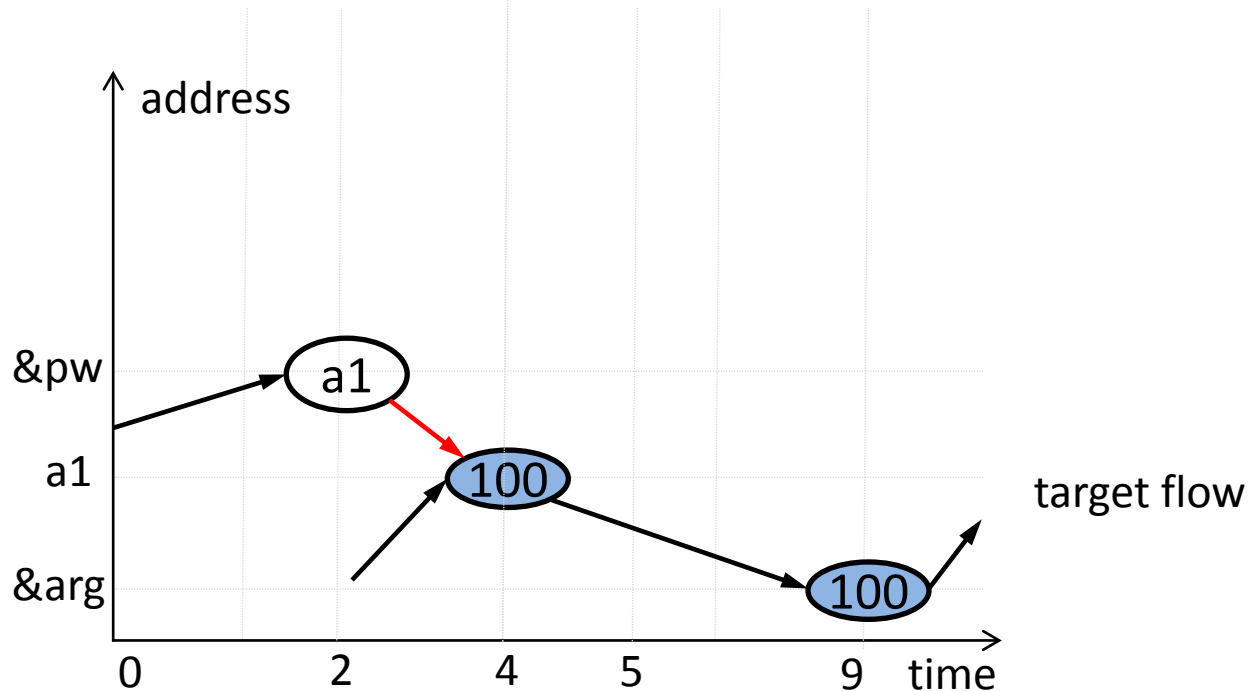
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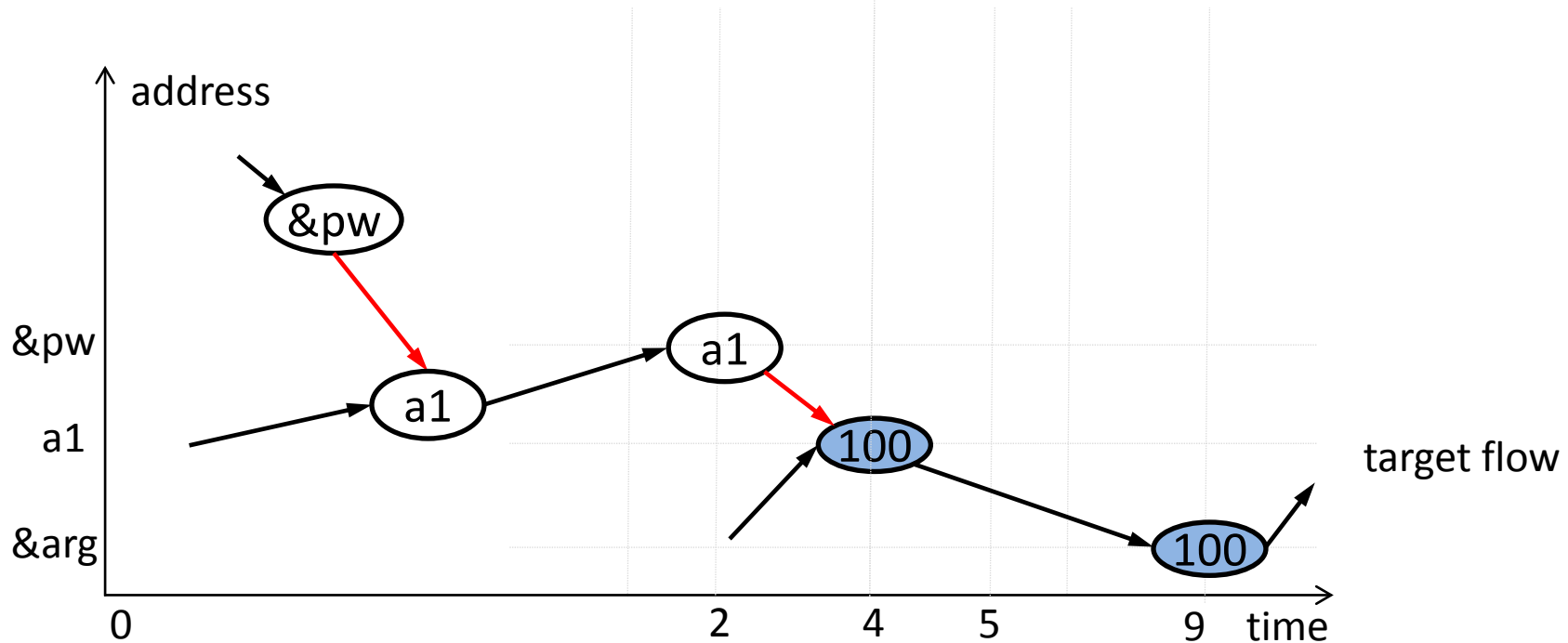
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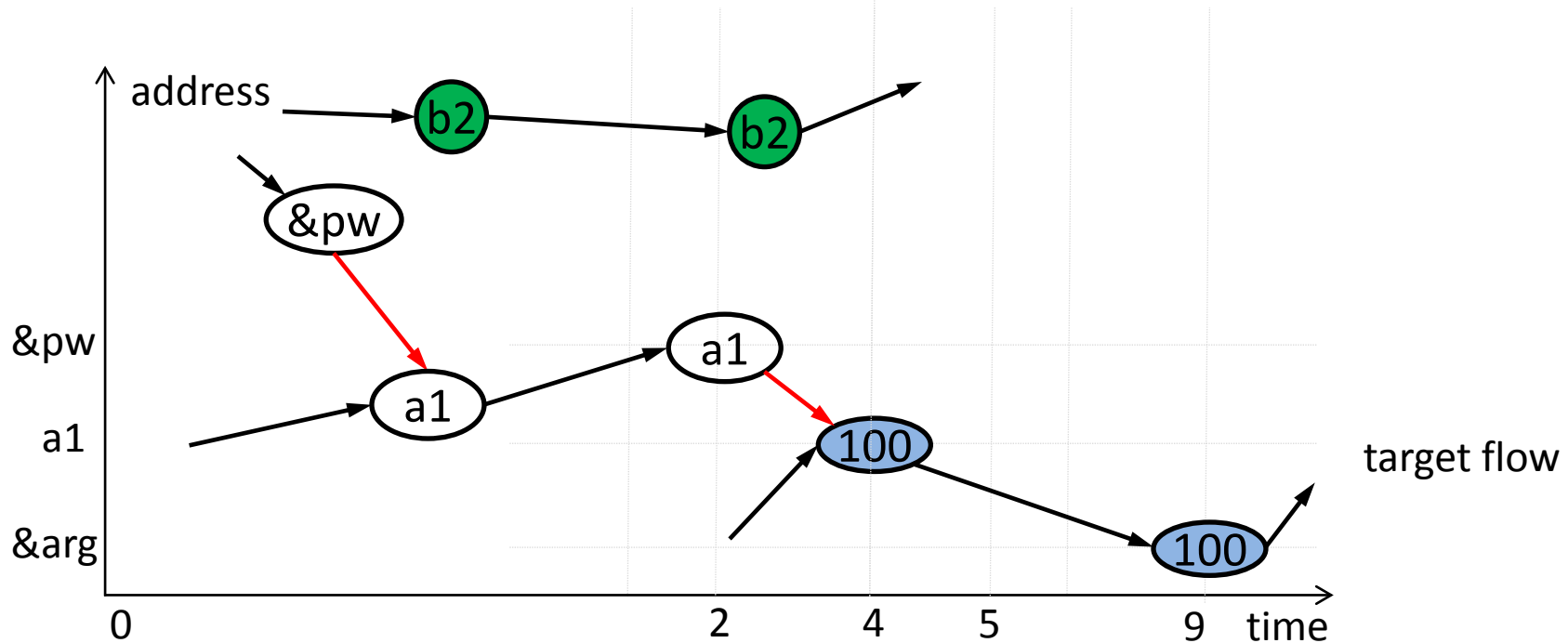
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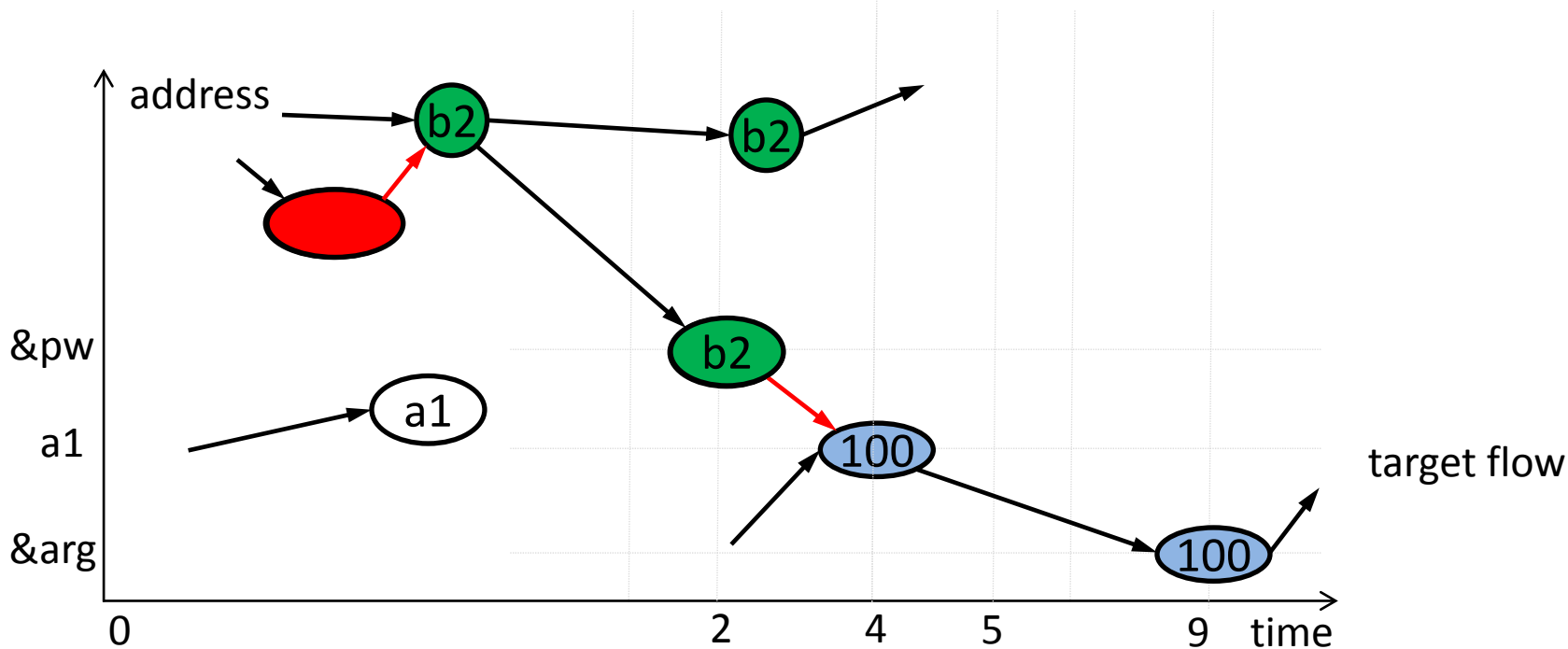
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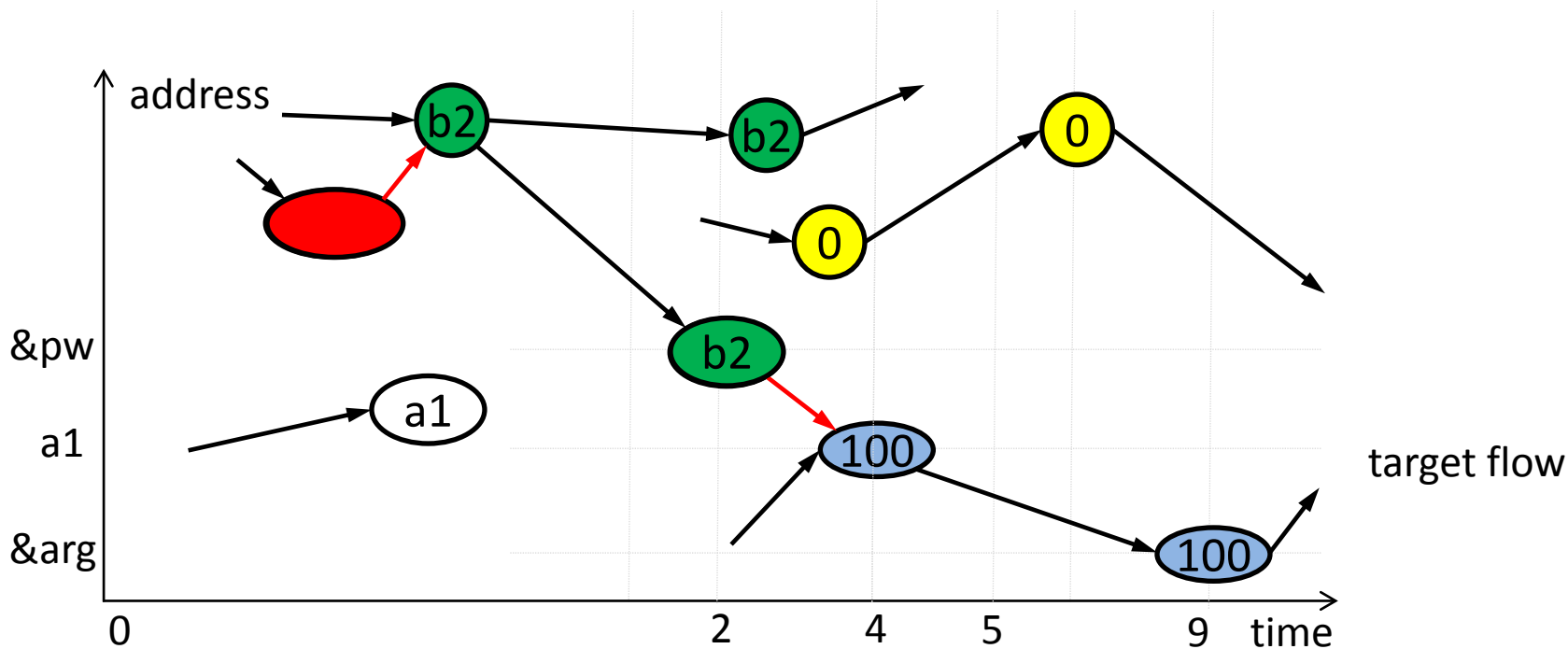
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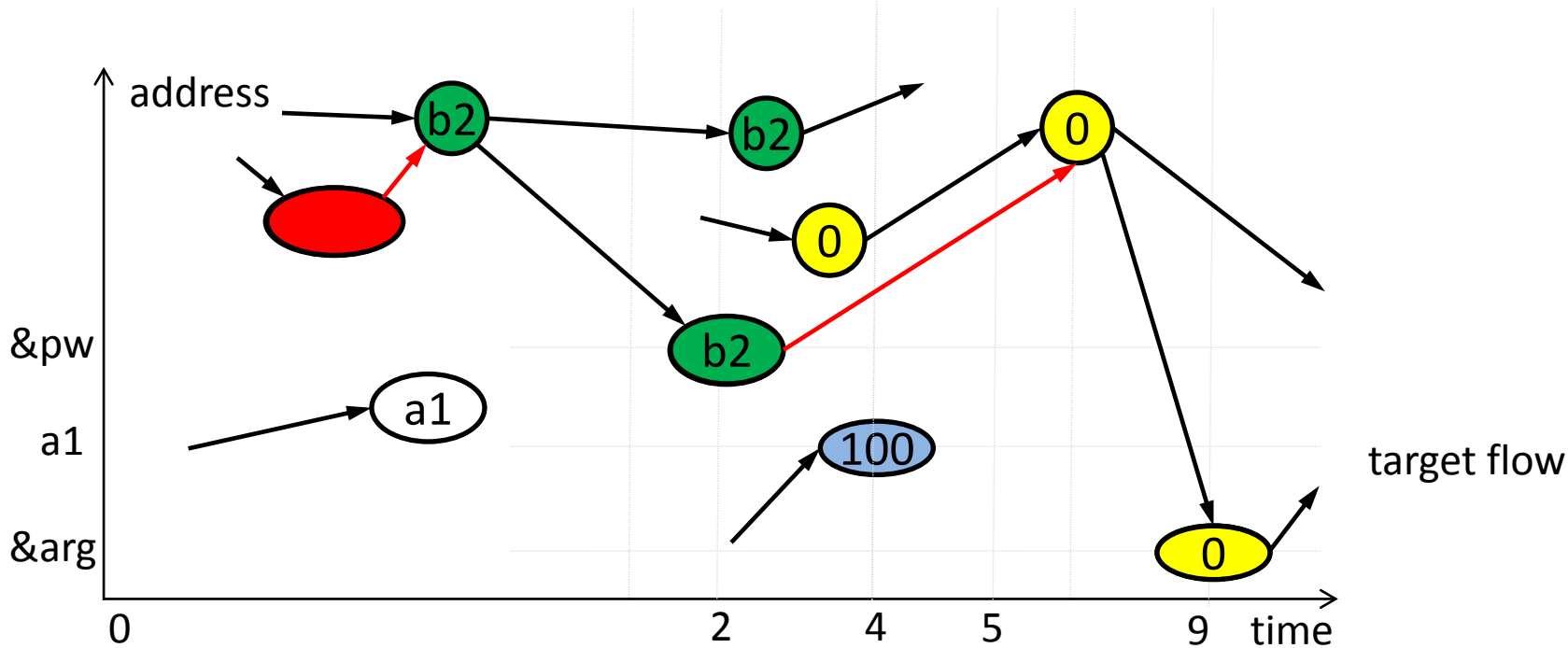
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- Multi-flow stitching
 - Intermediate data flows
 - Source flow -> flow 1 -> flow 2 -> ... -> Target flow

Defeat ASLR --- Address Reuse

- Partial reuse: offset is fixed

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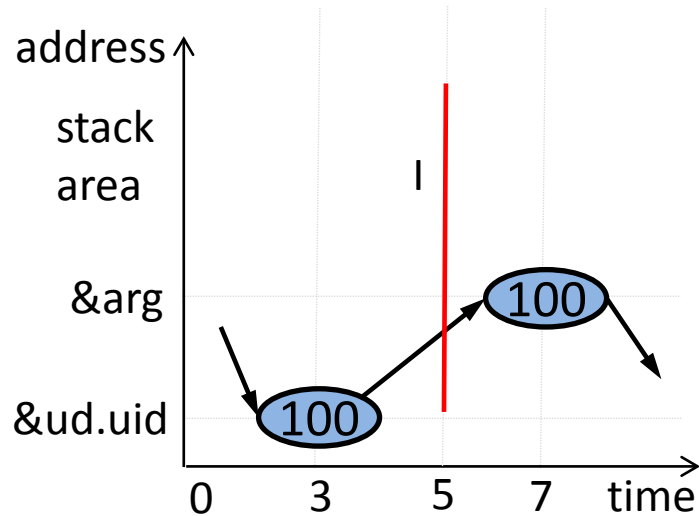
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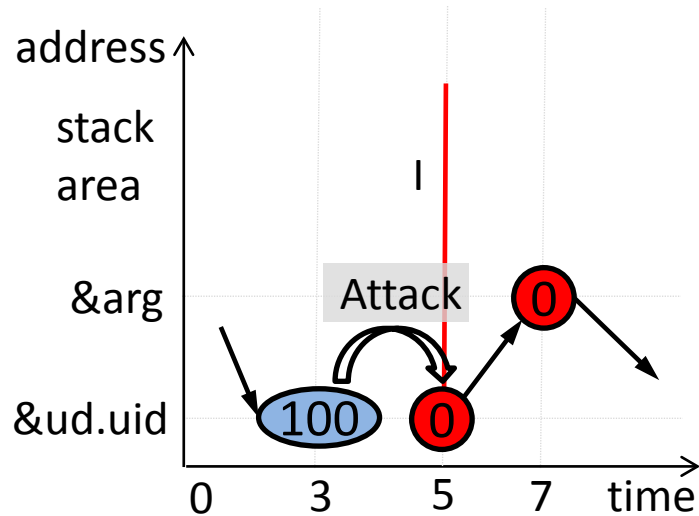
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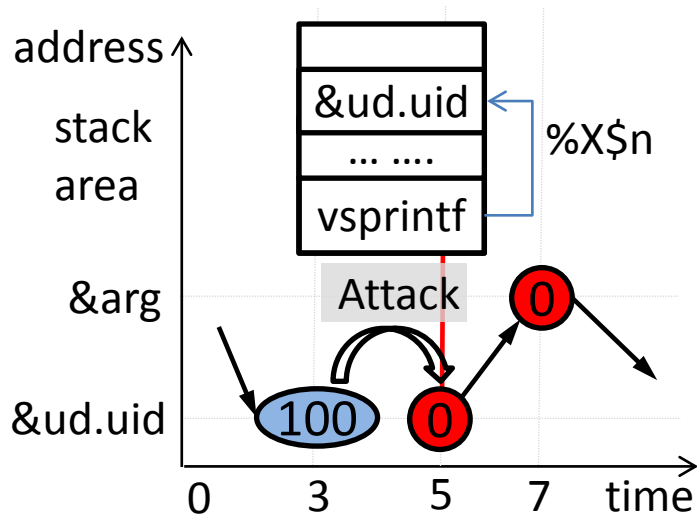
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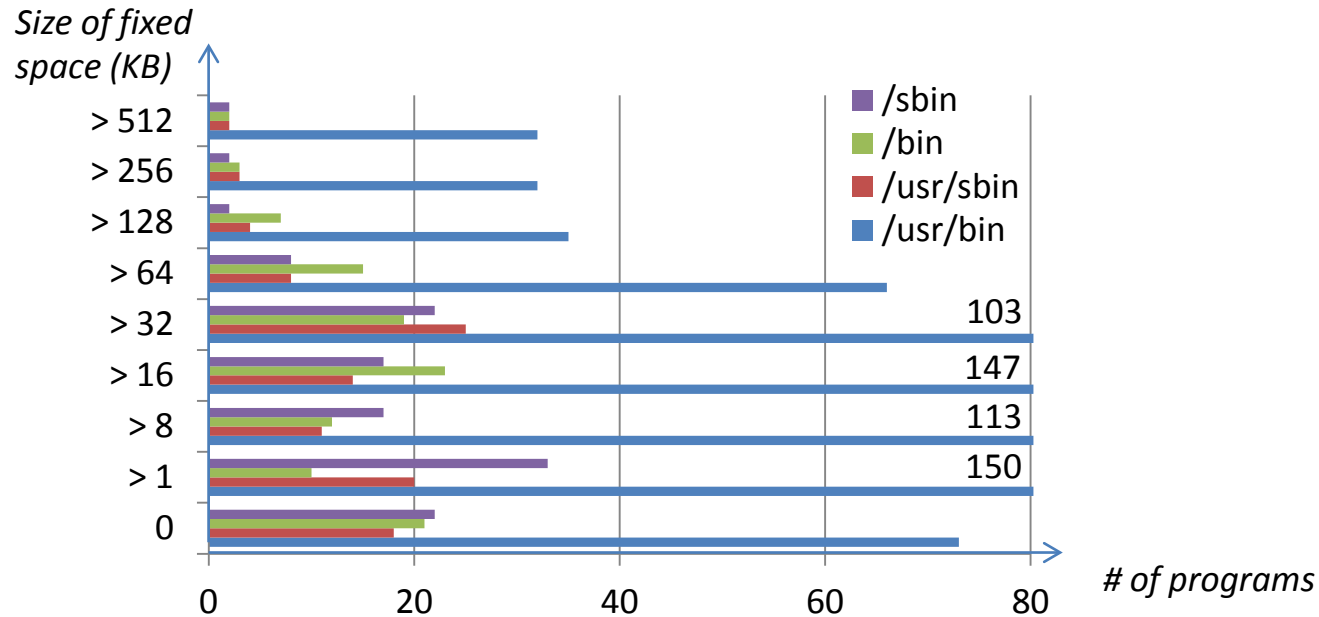
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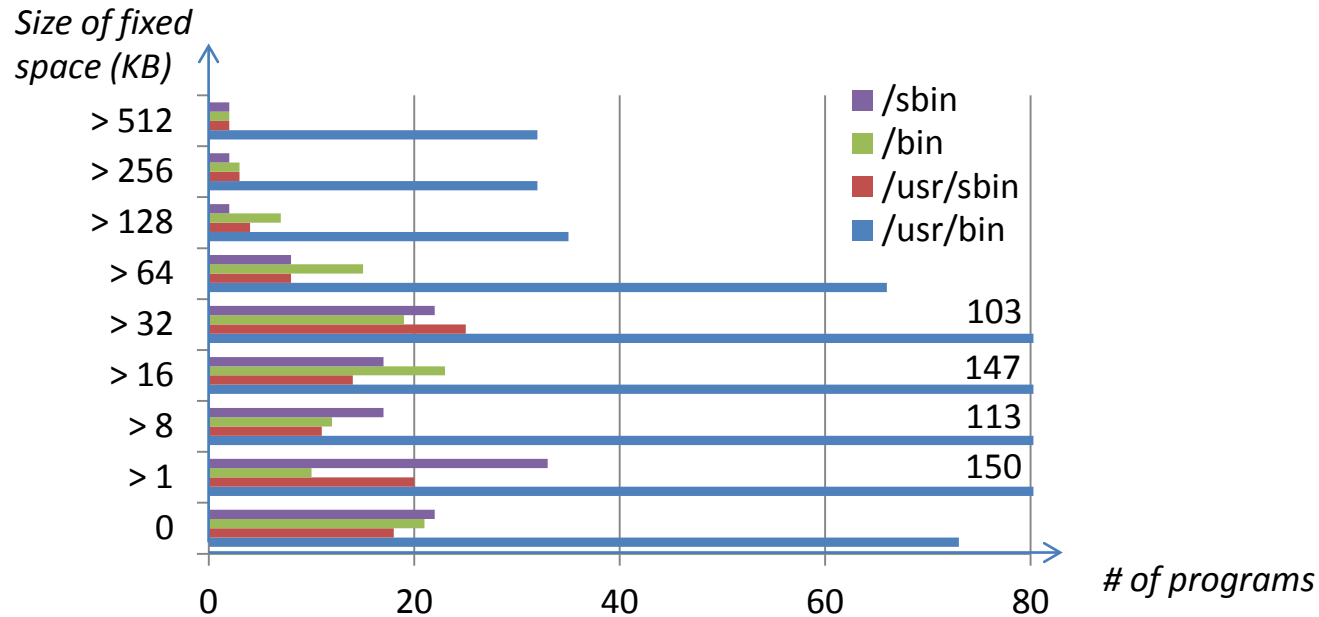
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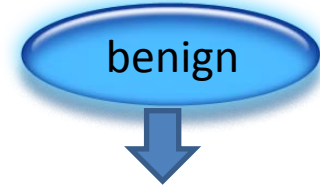
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- msvcrt71.dll, hxds.dll on Windows

FlowStitch



FlowStitch

error-exhibiting



benign



error-exhibiting trace



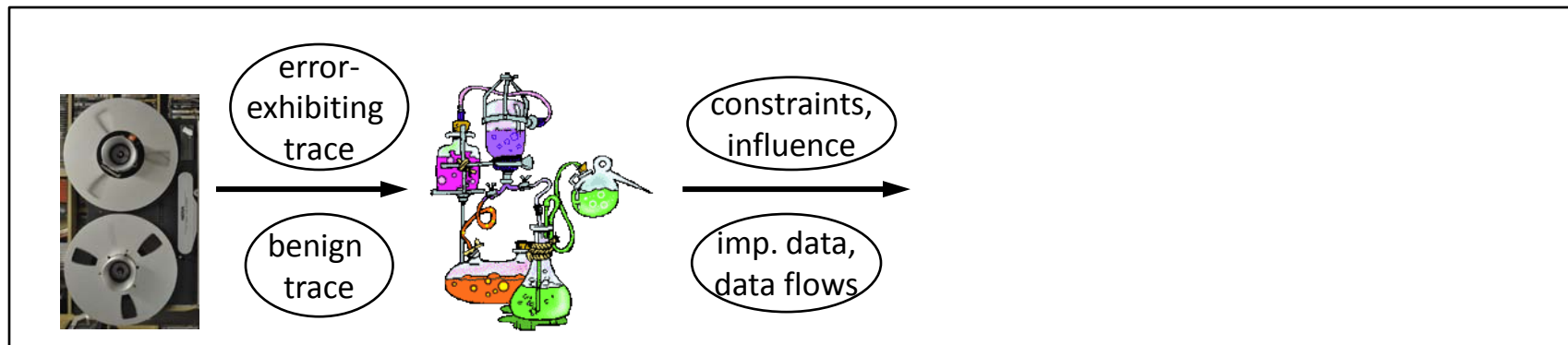
benign trace

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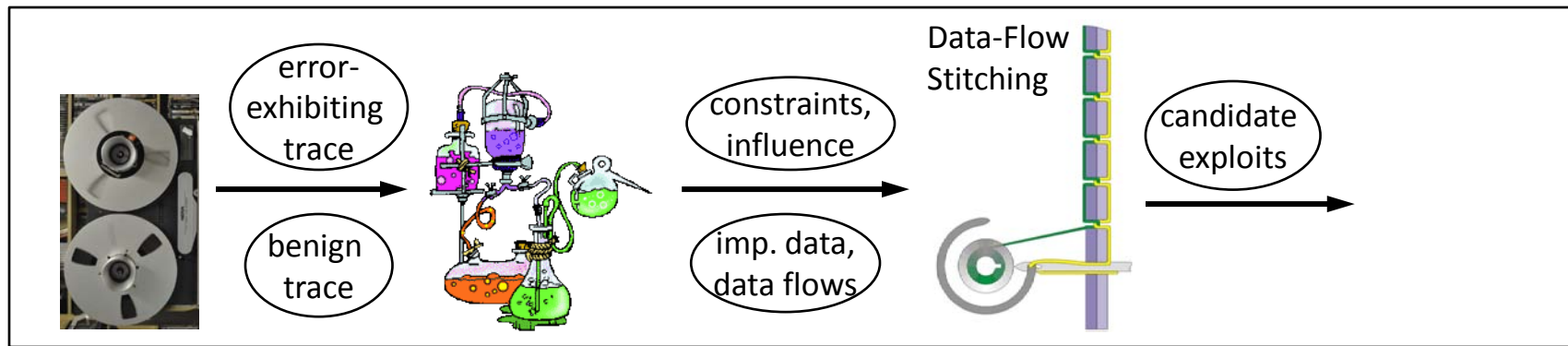
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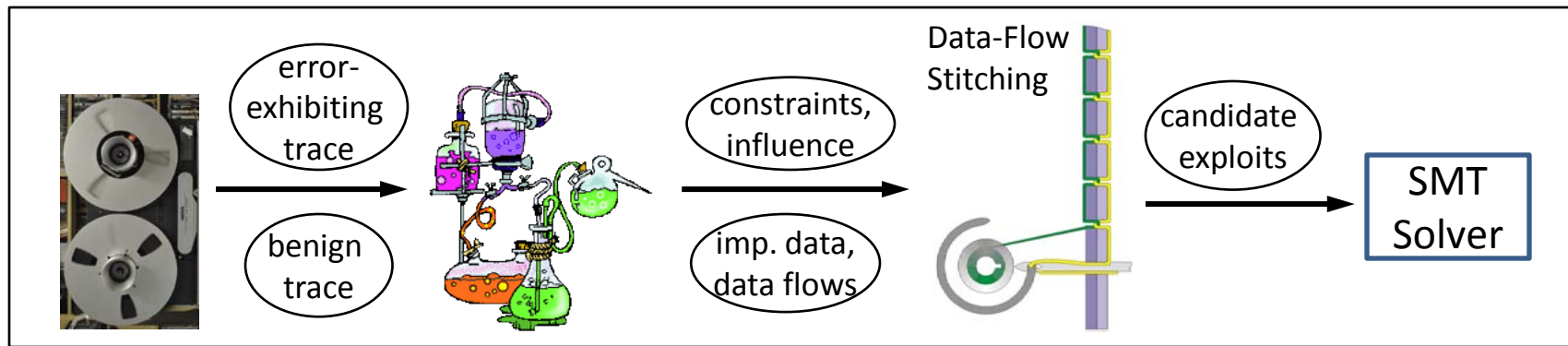
benign



FlowStitch

error-exhibiting

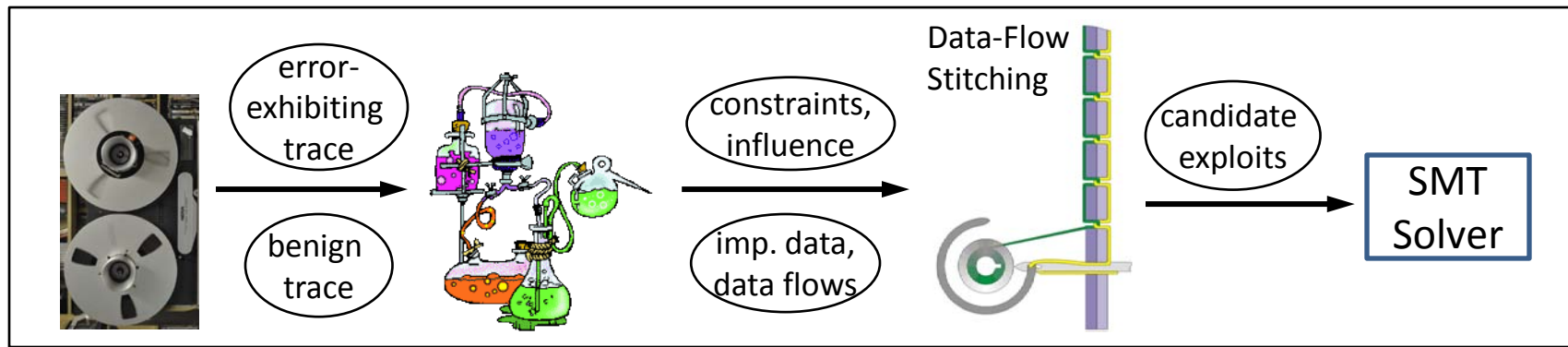
benign



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Evaluation --- Generated Exploits

ID	Vul. bin	Vulnerability	Data-Oriented Exploits	ASLR
CVE-2013-2028	nginx	Stack bof	L ₀ : private key M ₀ : http root dir	
CVE-2012-0809	sudo	Format string	M ₀ : user id	✓
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CVE-2001-0144 *	SSHD	integer overflow	L ₀ : root passwd hash	
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CVE-2000-0573 *	wu-ftpd	Format string	L ₀ : env variables	
			M ₀ : user id (single-edge)	✓
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* CVEs discussed in Shuo Chen's work [1]

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 - 2-level stitch

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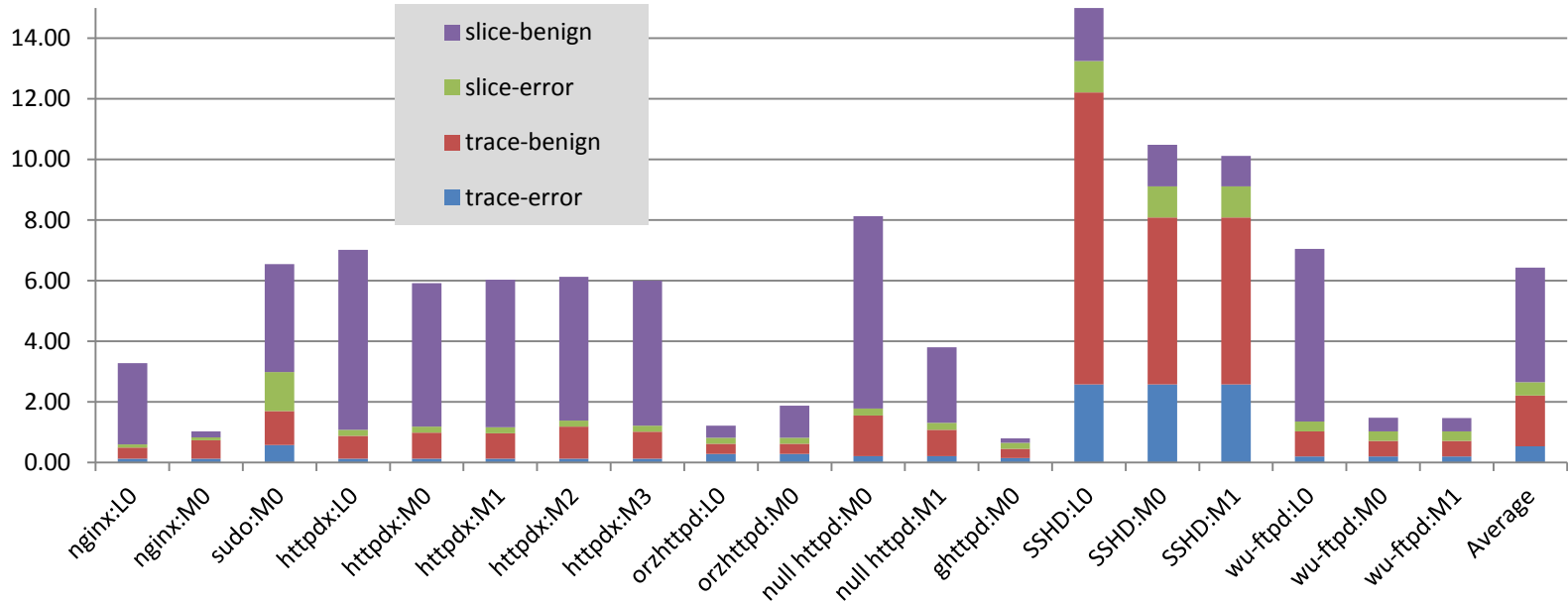
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- 7 advanced stitch
 - 2-level stitch
- 10 bypass ASLR
 - 8 fixed addresses
 - 2 address reuse

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Evaluation --- Performance



- 6.5 min/exploit
- Slice takes long
 - faster version is available (binary version)

Case Study – 2-Level Stitch

- *ghttpd* web server: stack buffer overflow

<pre>//serveconnection(): char *ptr; //URL pointer //esi is allocated for it 1: if(strstr(ptr,"/..")) reject the request; 2: log(...); 3: exec(ptr);</pre>	<pre>Assembly of log(...) push %ebp push %esi // stack overflow pop %esi pop %ebp ret</pre>	<pre>Assembly of line 3: push %esi ... call <exec@plt></pre>
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--	---	---

- Previous exploit^[1] does not work any more
 - Corrupt pointer ptr: *(ptr) -> url
- We build a 2-level stitch
 - Corrupt pointer saved ebp: (*(saved ebp)) -> *ptr -> url

Case Study – Sensitive Data Lifespan

- *SSHD* hashed key info leak
- *getspnam()* in *glibc* gets hashed key (heap copy)
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 - Overwritten by later usage
- *Challenging* to make lifespan correct!



Conclusion

- Rich Category: Data-Oriented Exploits
 - Single-edge stitch, Pointer stitch
 - N-level stitch, Multi-flow stitch
- Data Flow Stitching
 - Systematic way to generate data-oriented exploits
 - Agnostic to CFI, DEP and often ASLR
- Automatic construction is feasible

Thanks!

Hong Hu

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<http://www.comp.nus.edu.sg/~huhong/>