

Express: Lowering the Cost of Metadata-Hiding Communication with Cryptographic Privacy

Saba Eskandarian, Henry Corrigan-Gibbs, Matei Zaharia, Dan Boneh
Stanford MIT CSAIL Stanford Stanford

How Can We Protect Whistleblowers?



Private Communication Tools

End to end encrypted messaging apps

E.g. Signal, WhatsApp

Problem: **metadata**



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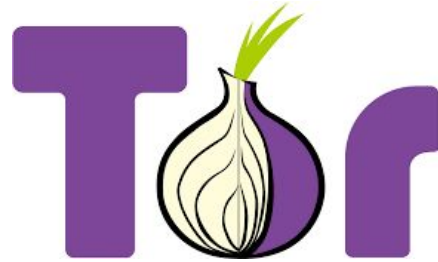
Problem: **metadata**



Anonymizing proxy

E.g. Tor, SecureDrop

Problem: **global adversaries**



Private Communication Tools

Metadata-hiding communication systems with cryptographic privacy

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Drawback: **heavy requirements placed on clients**

- Requirement to run in synchronized rounds
- High communication costs

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Fundamental issue: whistleblowing tools need *cover traffic*, which must be possible for clients to generate at minimal cost

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Fundamental issue: whistleblowing tools need *cover traffic*, which must be possible for clients to generate at minimal cost

Can we get around high client costs?

Express: Practical Metadata-Hiding Whistleblowing

Qualitative improvement: users do not access the system in synchronized rounds

Asymptotic improvements:

Client computation: $O(1)$

Communication: $O(1)$

Prior work: $O(\sqrt{N})$

Practical improvements:

6x faster server

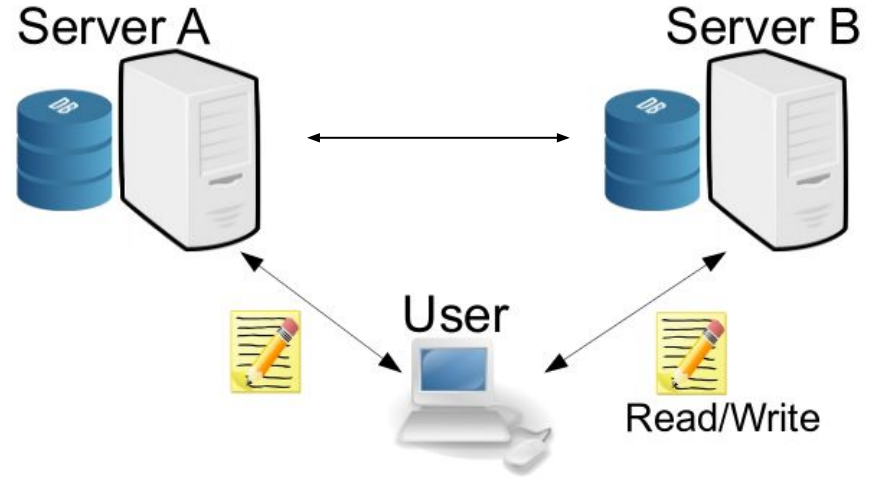
8x faster client

>10x communication reduction

6x reduction in dollar cost to run

Express Overview

2 server system, secure against:
Up to one corrupt server
Arbitrarily many corrupt users



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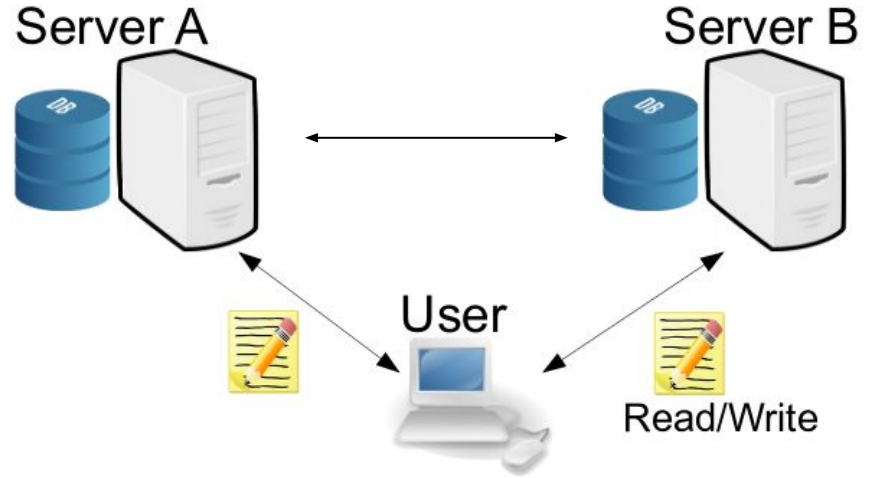
- Arbitrarily many corrupt users

Supported operations:

- Register mailbox

- (Private) write to mailbox

- Read from mailbox



Express Overview

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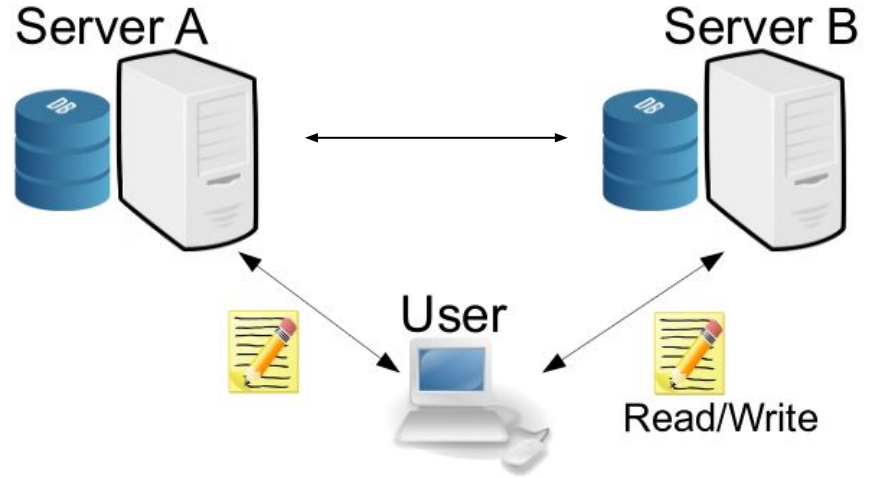
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Supported operations:

- Register mailbox

- (Private) write to mailbox

- Read from mailbox



Security: can't tell who the *recipient* of a message is

Tool: Private Writing



I want to write
"Hi!" to address 3

Addr	Data
0	0
1	0
2	0
3	0
4	0



Addr	Data
0	0
1	0
2	0
3	0
4	0

Tool: Private Writing



Addr	Data
0	0
1	0
2	0
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x	f(x)
0	0
1	0
2	0
3	"Hi!"
4	0



Addr	Data
0	0
1	0
2	0
3	0
4	0

Distributed Point Functions and their Applications, Niv Gilboa, Yuval Ishai, *Eurocrypt'14*.
Private Information Storage, Rafail Ostrovsky, Victor Shoup, *STOC'97*

Tool: Private Writing



Addr	Data
0	0
1	0
2	0
3	0
4	0



x	$f_1(x)$
0	"abc"
1	"xf\$"
2	"^tg"
3	"!7≈"
4	"jhV"

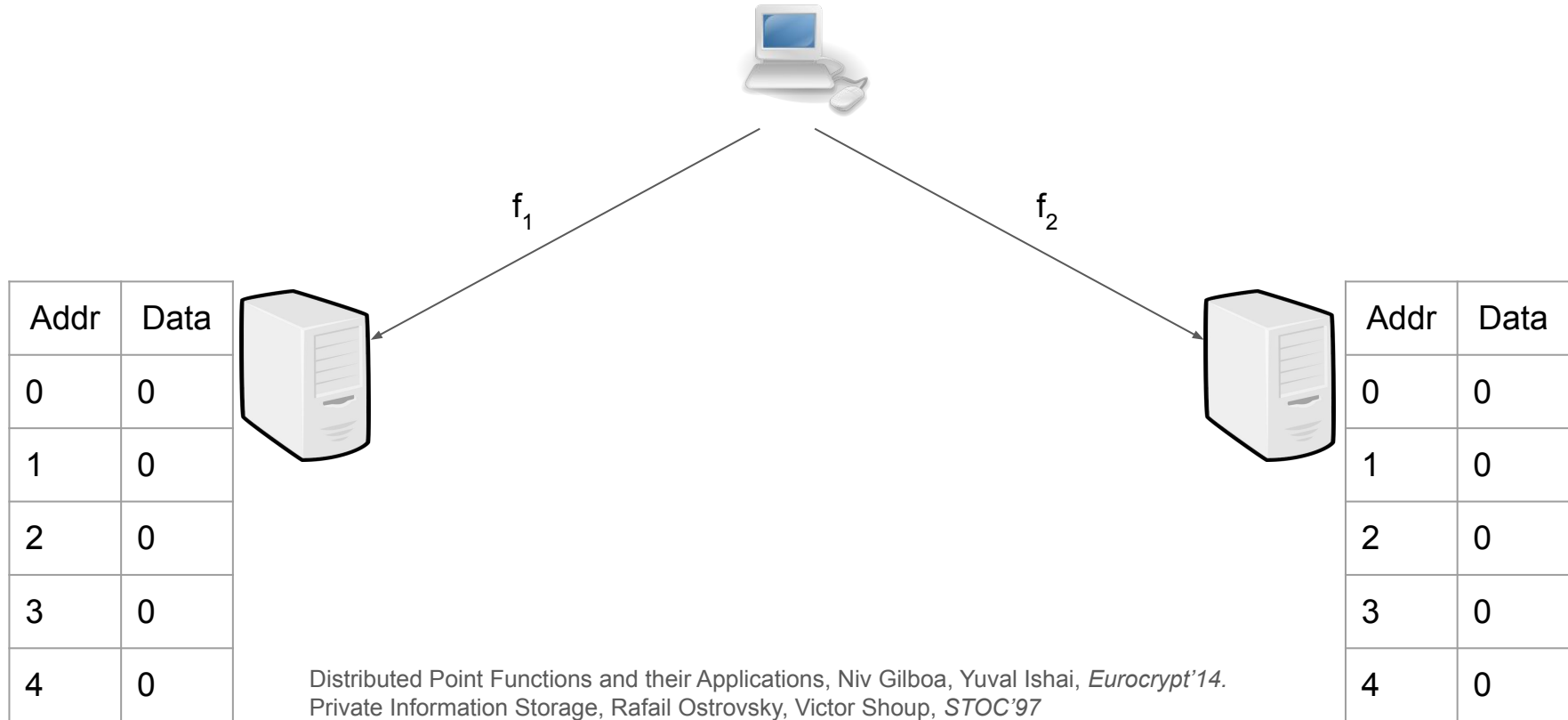
x	$f_2(x)$
0	"abc"
1	"xf\$"
2	"^tg"
3	"!2!)"
4	"jhV"



Addr	Data
0	0
1	0
2	0
3	0
4	0

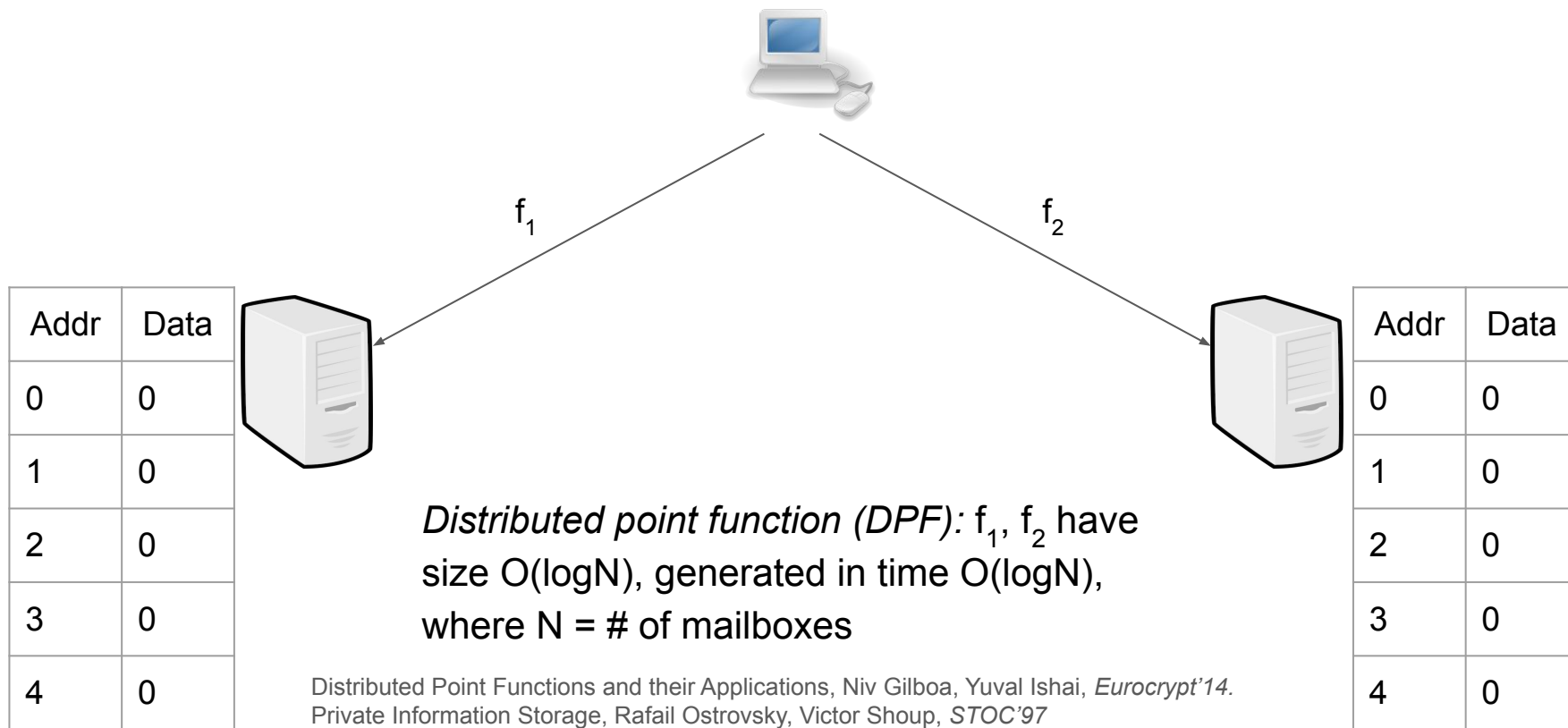
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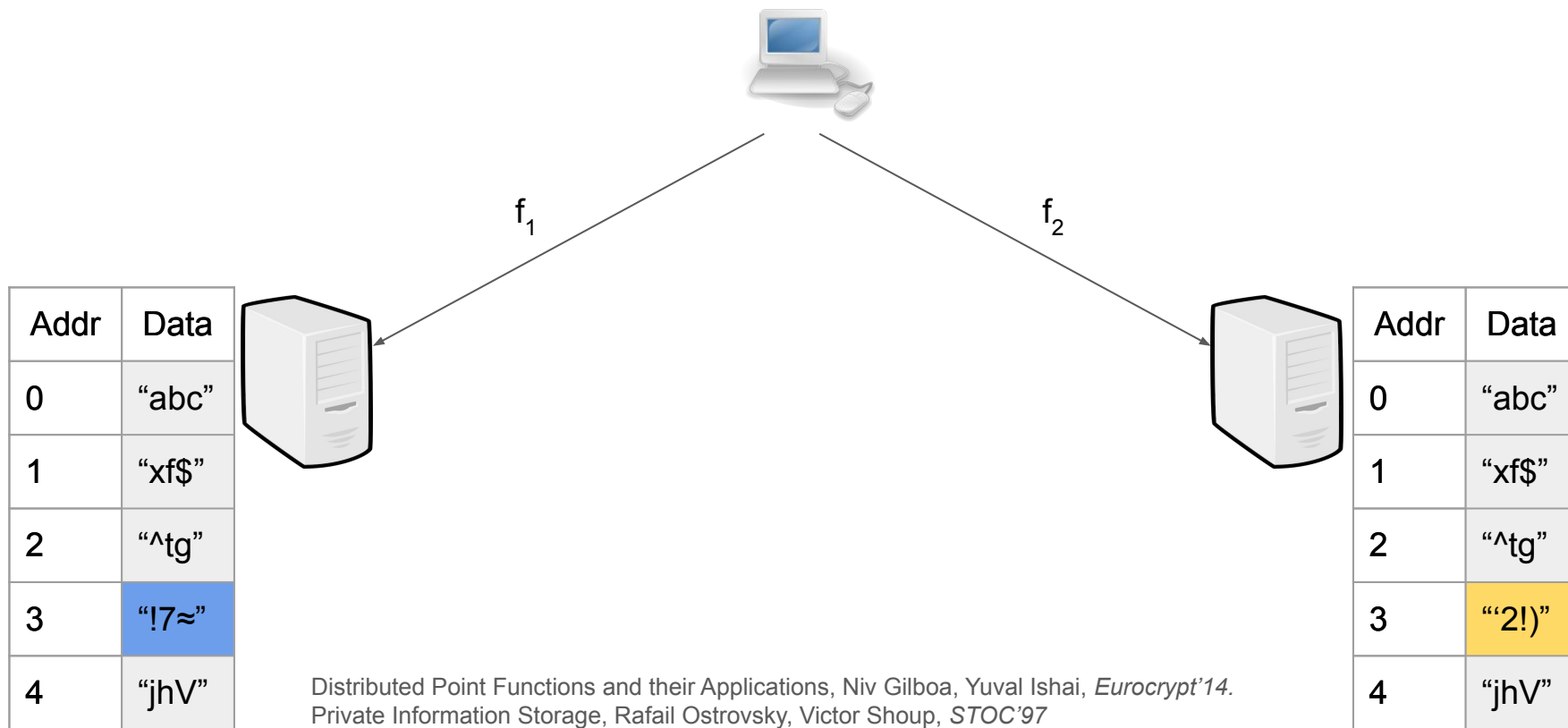


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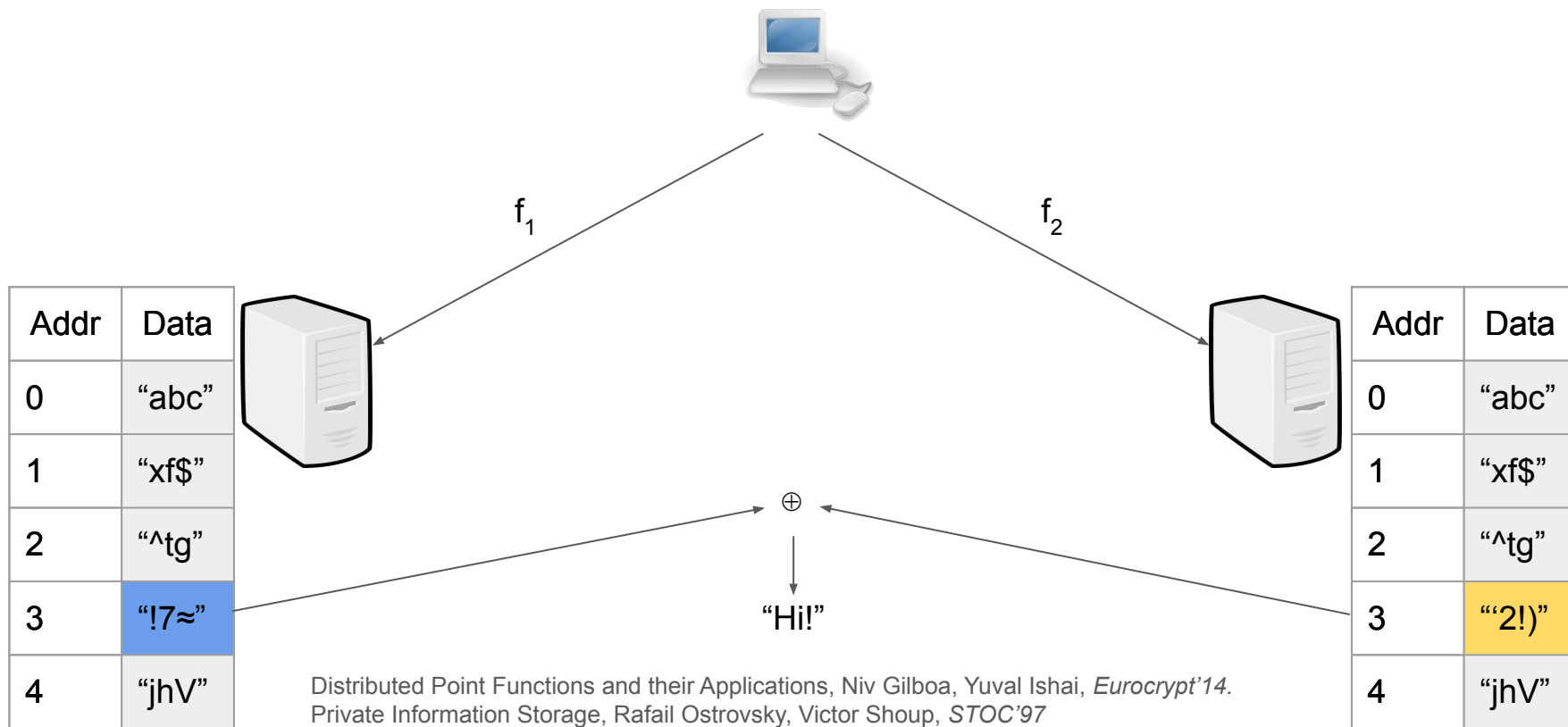


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Tool: Private Writing



Filtering out Malformed DPFs

Problem: disruptive user sends malformed message to corrupt mailboxes



x	f(x)
0	989f4
1	dDf73
2	08dji3
...	...
N	89hfif

Filtering out Malformed DPFs

Problem: disruptive user sends malformed message to corrupt mailboxes

Solution: servers blindly *audit* all incoming write requests

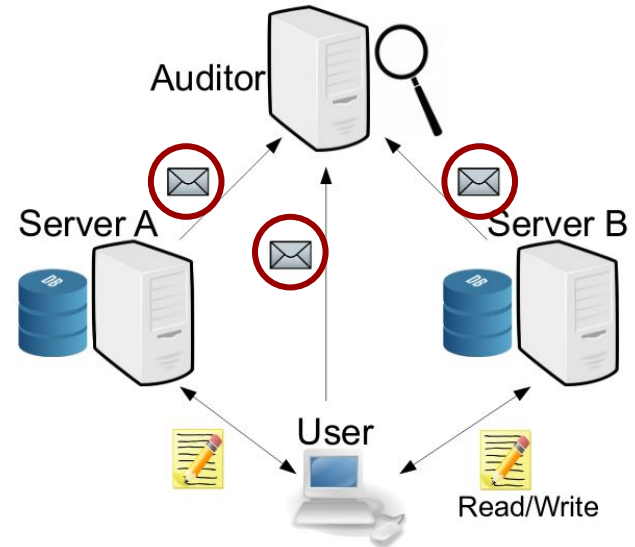
Filtering out Malformed DPFs

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Solution: servers blindly *audit* all incoming write requests

Prior work: third server audits requests

- $O(\sqrt{N})$ communication
- $O(\sqrt{N})$ client/auditor computation



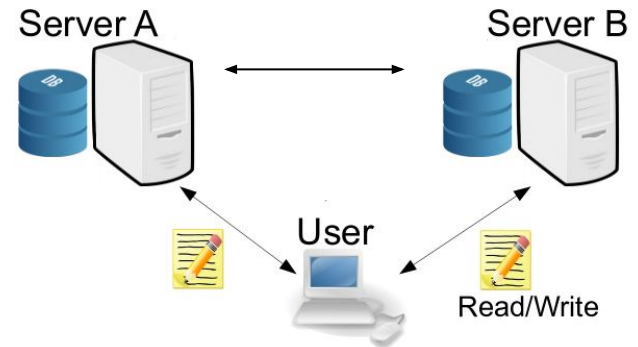
Filtering out Malformed DPFs

Problem: disruptive user sends malformed message to corrupt mailboxes

Solution: servers blindly *audit* all incoming write requests

New auditing protocol:

- $O(1)$ communication
- $O(1)$ client computation
- No additional server!



Auditing

Goal: check that values held by servers only differ at one point



Server A



Server B

Auditing

Goal: check that values held by servers only differ at one point

Prior work has a semihonest solution where servers use a cheap MPC (only 2 multiplications) to verify this property.



Server A



Server B

Auditing

Goal: check that values held by servers only differ at one point

Prior work has a semihonest solution where servers use a cheap MPC (only 2 multiplications) to verify this property.



Server A



Server B

Issue: malicious server can guess and check the differing entry

Auditing

Tool: secret-shared non-interactive proofs (SNIPs)

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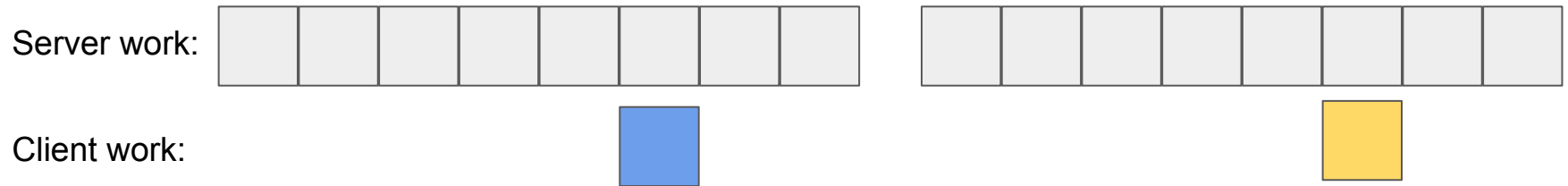
Idea: client sends SNIP proof to servers that honest evaluation of the semihonest protocol accepts

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Key Insight: client knows the message index, only needs $O(1)$ work to prove facts about computation that would take servers $O(N)$ work

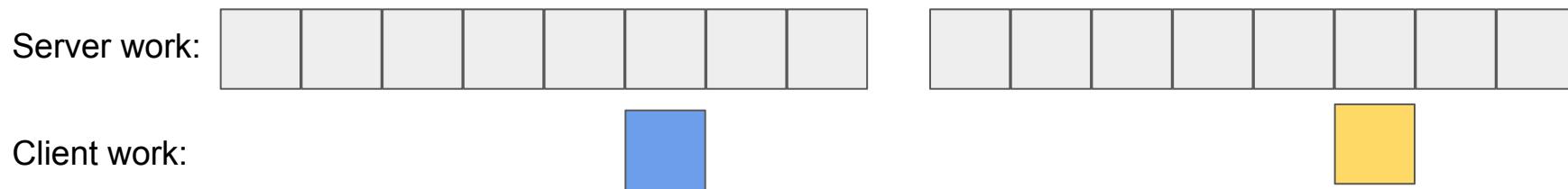


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See paper for details

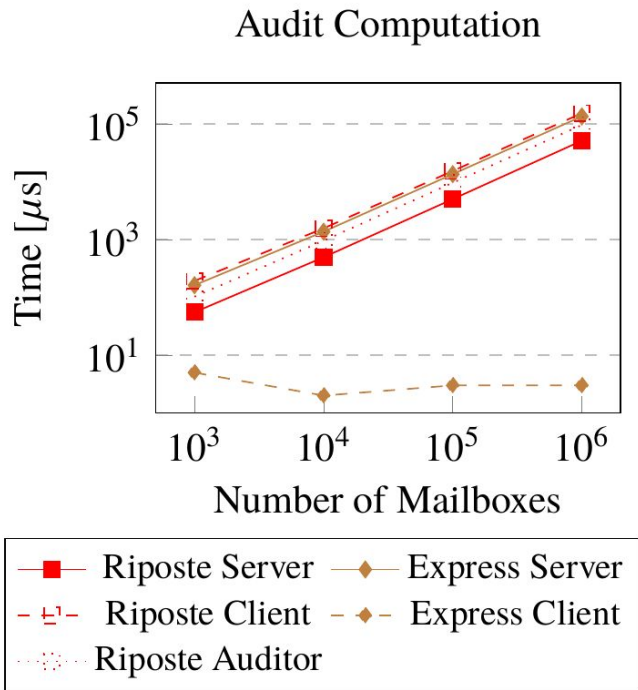
Evaluation: Auditing Protocol

Client runs in under 5 *microseconds*

55,000x faster than Riposte for 1M mailboxes

Enables 8x reduction in overall client computation (now 20ms)

Comparable on server, where auditing is not the bottleneck



Express

Metadata-hiding communication system with application to private whistleblowing

Asymptotic speedup from $O(\sqrt{N})$ to $O(1)$ for auditing

Speedup of 8x on client, up to 6x on server (compared to Riposte)

6x lower dollar cost to operate system

13-7,000x or more reduction in communication costs

Code: <https://github.com/SabaEskandarian/Express>

Contact: saba@cs.stanford.edu