DICE*: A Formally Verified Implementation of DICE Measured Boot

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Microsoft[®]

Establishing Trust in a Remote Device

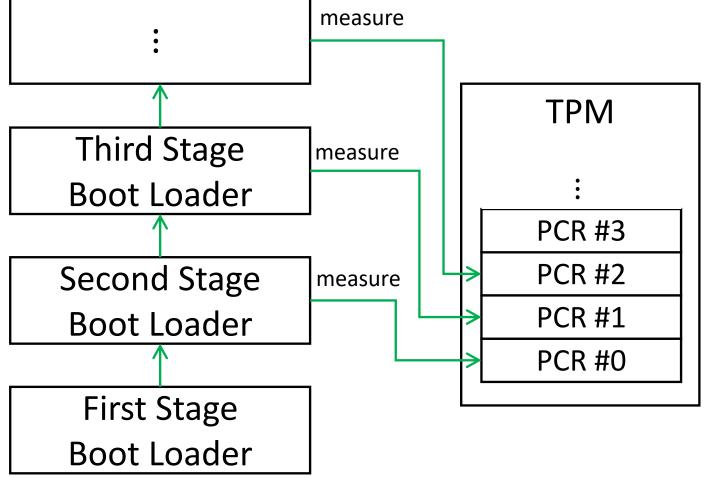
Send sensitive data to ML accelerators on the cloud



How do we verify that a device is running expected code?

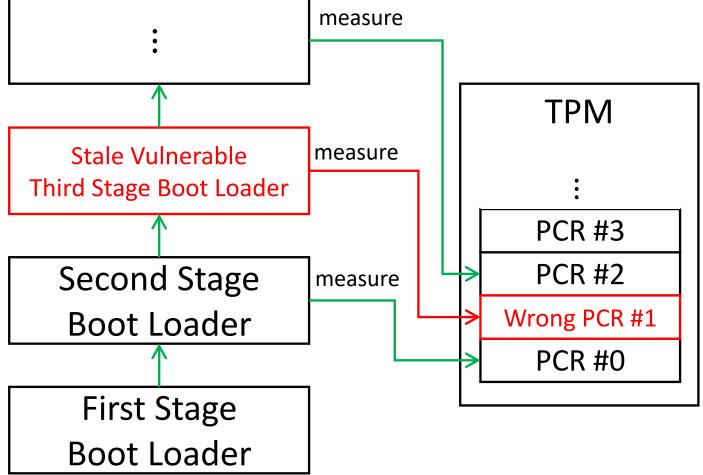
Prevents an attacker from loading unexpected code in boot by

- measuring the boot sequence
- recording the measurements for later attestation



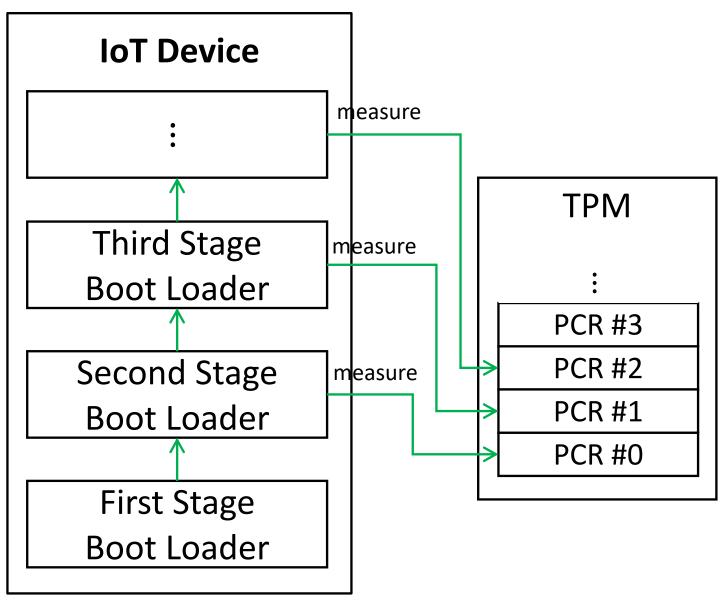
TPM Trusted Platform Module

Unexpected code results in wrong measurement and fails attestation



TPM Trusted Platform Module

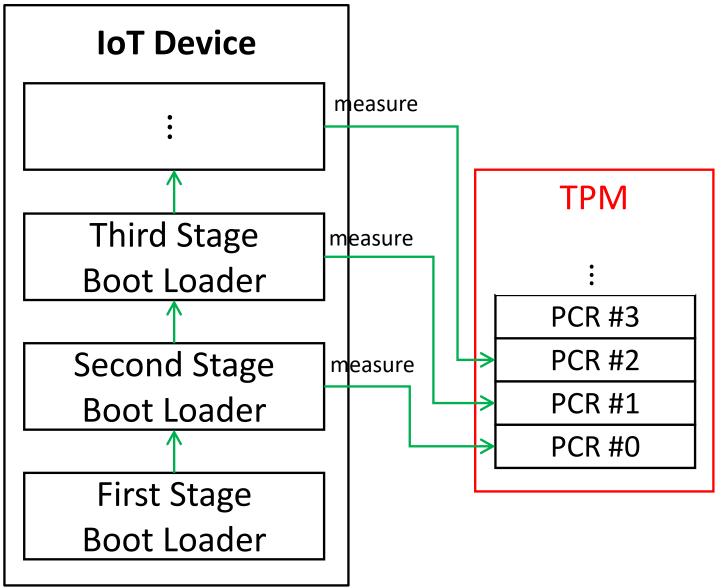
But traditional measured boot protocols are not applicable to IoT devices



TPM Trusted Platform Module

But traditional measured boot protocols are not applicable to IoT devices

Because they require a dedicated chip like TPM, which is too expensive in terms of cost, power, or real estate

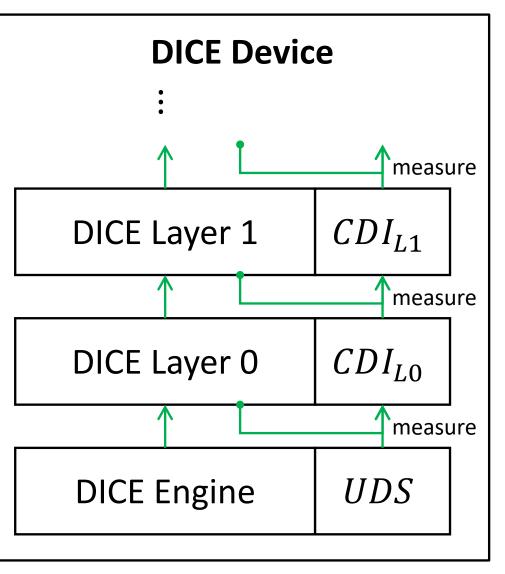


TPM Trusted Platform Module

Lightweight measured boot for IoT devices proposed by **Trusted Computing Group**.

DICE is becoming important. TCG members like Microsoft, STMicro, Microchip, Micron, NXP, etc. are behind its effort.

DICE is general for scenarios beyond IoT devices, like servers.

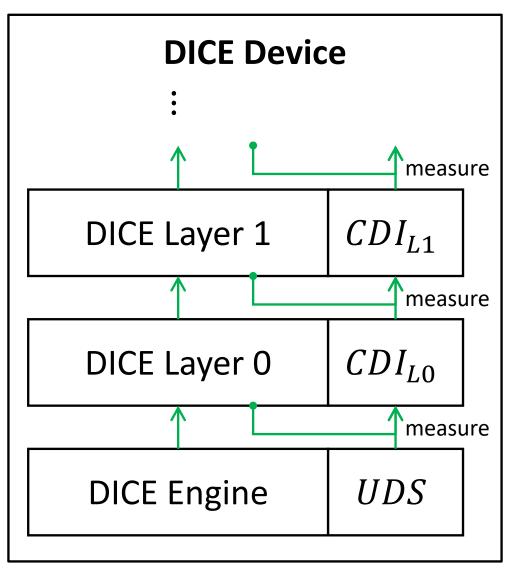


https://trustedcomputinggroup.org/work-groups/dice-architectures/

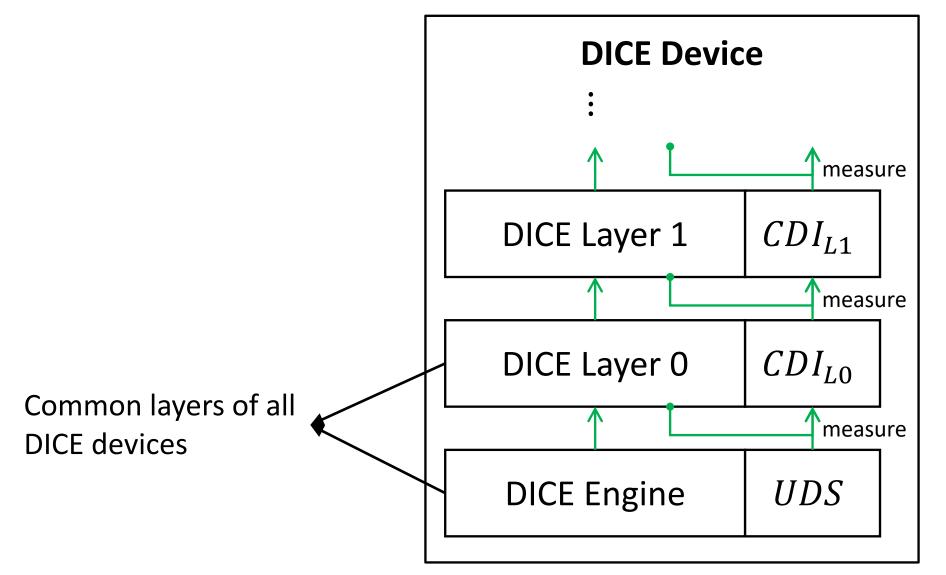
UDS Unique Device Secret

DICE implicitly captures TCB as secrets (*CDI*) derived during boot

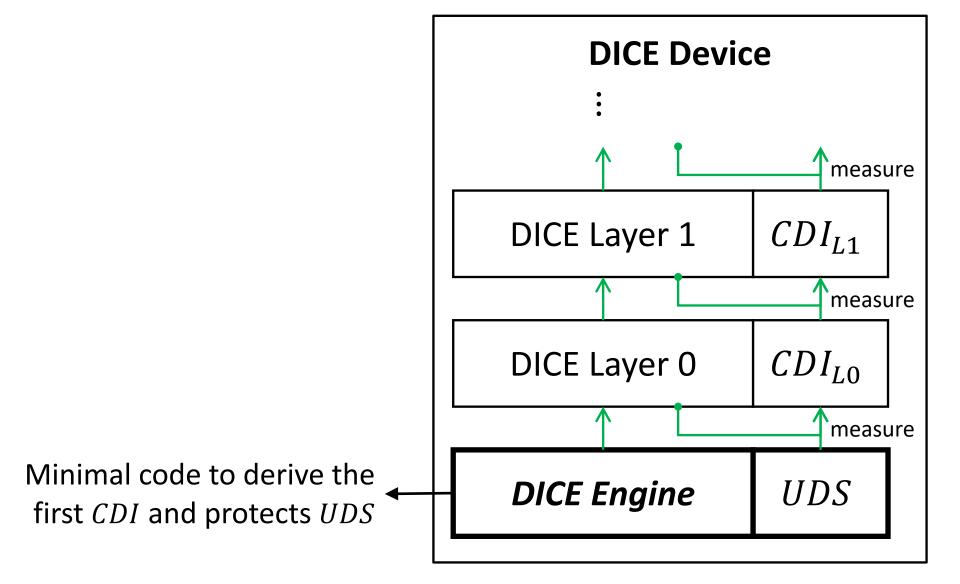
Layered structure: each layer extends the TCB by measuring the upper layer and deriving a new *CDI*.



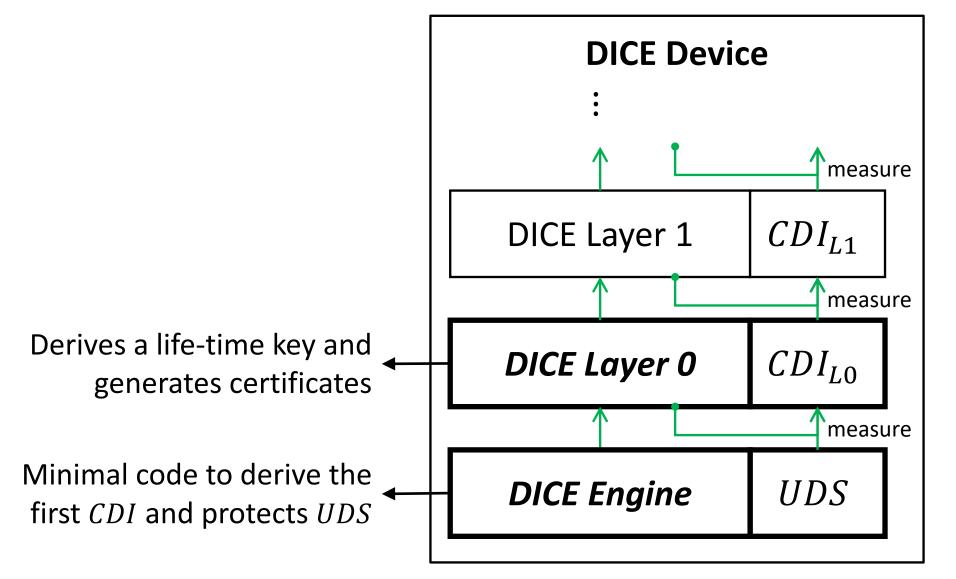
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DICE implementation is hard to get right because of key derivation, hashes, signatures and X.509 certificates — complex piece of code



National Security Agency | Cybersecurity Advisory

Patch Critical Cryptographic Vulnerability in Microsoft Windows **Clients and Servers**



Summary

NSA has discovered a critical vulnerability (CVE-2020-0601) affecting Microsoft Windows®1 crystocraphic functionality The certificate validation vulnerability allows an attacker to undermine how Windows verifies cry enable remote code execution. The vulnerability affects Windows 10 and Windows Server 2016 applications that rely on Windows for trust functionality. Exploitation of the vulnerability allows a network connections and deliver executable code while appearing as legitimately trusted entitie validation of trust may be impacted include:

- HTTPS connections
- Signed files and emails
- Signed executable code launched as user-mode processes

The vulnerability places Windows endpoints at risk to a broad range of exploitation vectors. NSA

Critical crypto bug leaves Linux, hundreds of apps open to eavesdropping

This GnuTLS bug is worse than the big Apple "goto fail" bug patched last week.

进CVE-2016-2108 Detail

Current Description

The ASN.1 implementation in OpenSSL before 1.0.10 and 1.0.2 before 1.0.2c allows remote attackers to execute arbitrary code or cause a denial of service (buffer underflow and memory corruption) via an ANY field in crafted serialized data, aka the "negative zero" issue.

Severity CVSS Version 3.x CVSS Version 2.0

CVSS 3.x Severity and Metrics:

Base Score: 9.8 CRITICAL

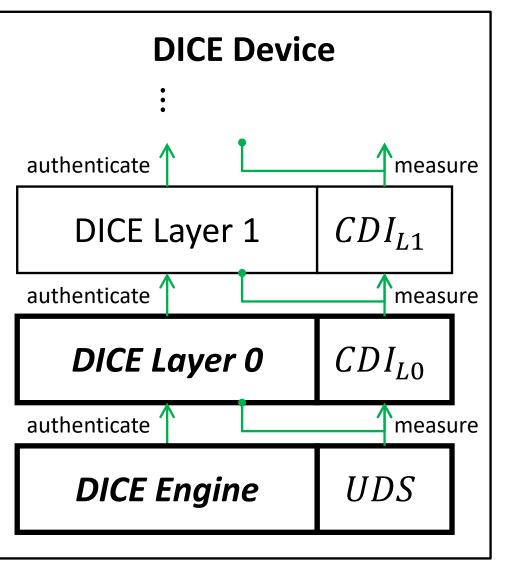
Vector: CVSS:3.0/AV:N/AC:L/PR:N/UI:N/S:U/C:H/I:H/A:H

ctor strings and CVSS scores. We also display any CVSS information provided within the

DICE is hard to get right because of the complex code and libraries

And bugs like memory errors, misuse of secrets, malleability attacks on X.509, side-channels may leak secrets allowing impersonation attack

But patching the first two layers is either impossible or extremely expensive

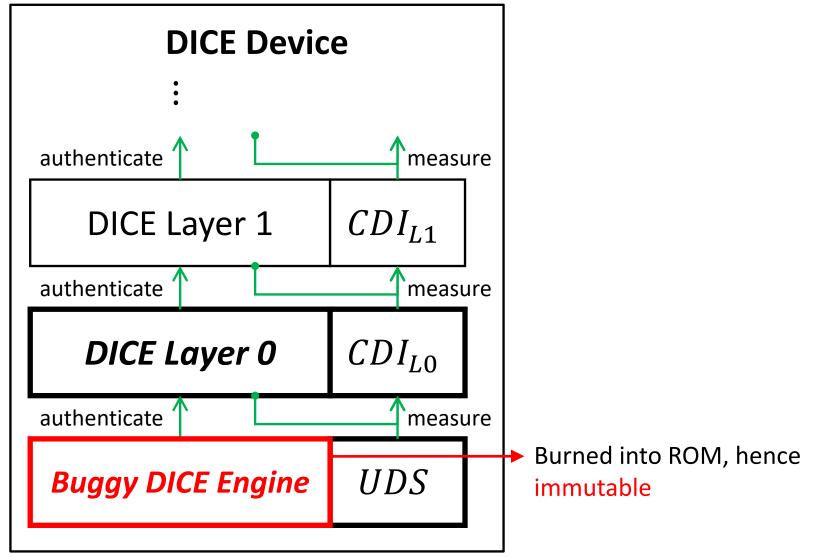


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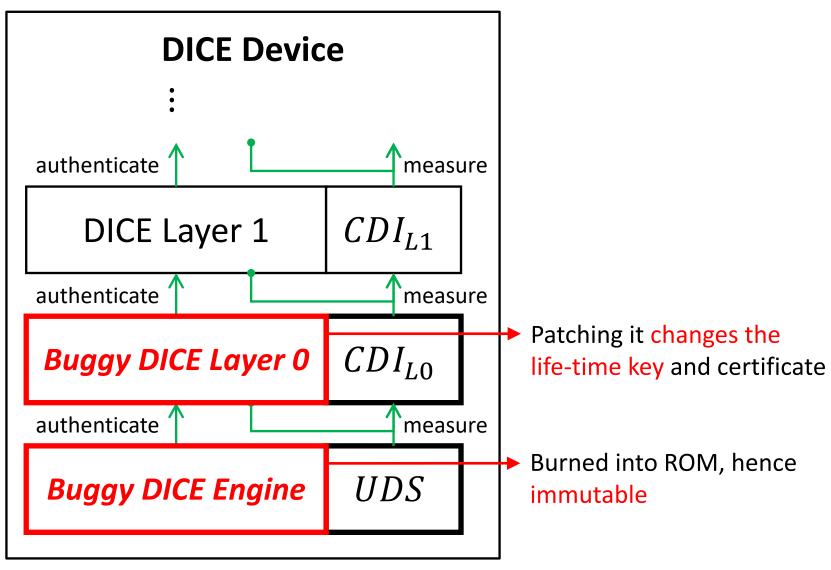


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UDS Unique Device Secret

DICE*: A Formally Verified DICE implementation

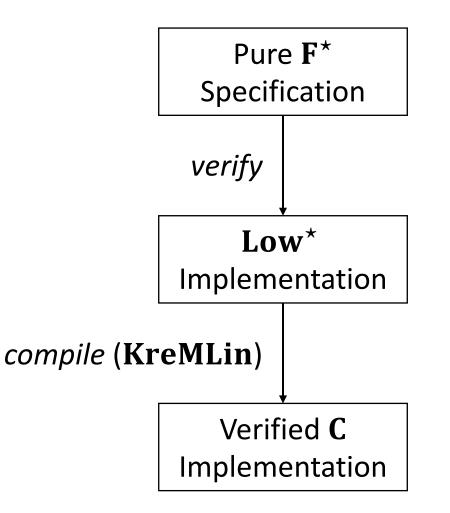
- Formally specify DICE specification
- Present a verified DICE Engine with a platform-agnostic interface
 - Users can focus on analyzing the platform-specific components
- Present a verified DICE Layer 0
 - Including a verified library for a subset of X.509 which can be extended and reused
- Generate verified C implementation and evaluate it on STM32H753ZI
 - Comparable to unverified hand-written code in terms of boot time and binary size
- Available at https://github.com/verified-HRoT/dice-star

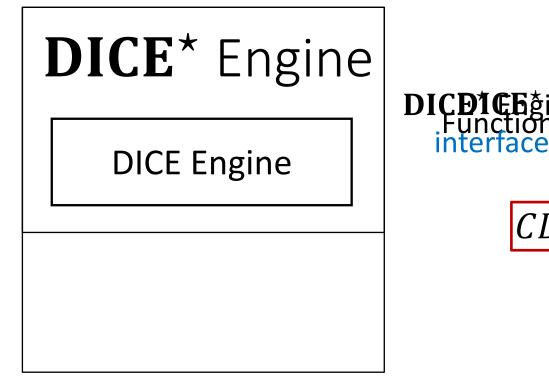
Verified Properties

- Functional correctness
 - Secrets, keys and certificates are derived as per specification
- Memory safety
 - Buffer overflows, use-after-free, no null dereferences, dangling pointers, etc.
- Confidentiality
 - No secret leakage via outputs, memory, etc.
- Side-channel resistance
 - Free of certain timing- and cache-based side channels
- X.509 certificate security
 - No malleability attacks

Verification Toolchain

- **F***:
 - Functional language with effects
 - Dependent type
 - Semi-automated proof via SMT solvers
- Low*: a shallow embedding of C in F^{\star}
 - C-like memory model
 - First Order
 - C-compatible types
- KreMLin: a Low*-C compiler





DICDICE gibegineckeluse blattchchanghlasih Functional Specification of DICE Engine interface for UDS in different platforms

$$CDI_{L0} = HMAC(UDS, Hash(L0))$$

 Platform-Agnostic

 Interface

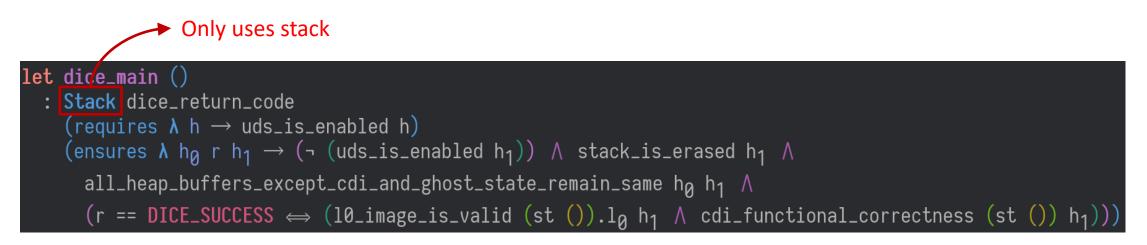
 • Read UDS

 • Latch UDS

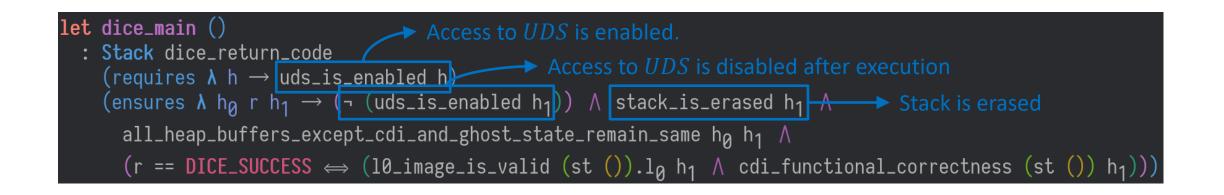
 • Erase stack

Enforces the folition of the the folition of the state in the state of the state of

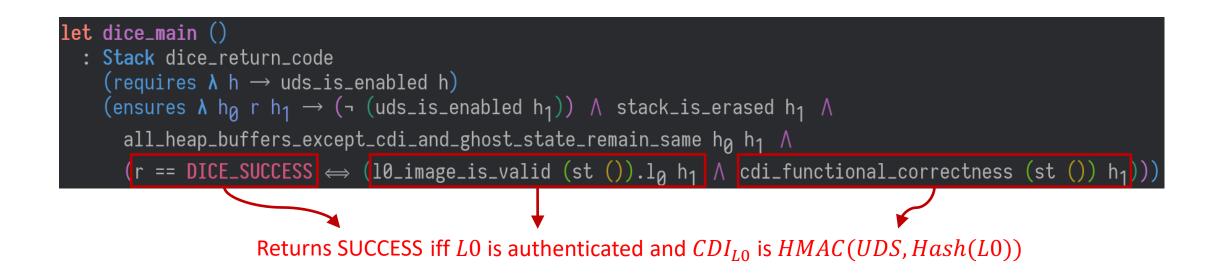
- Cannot read UDV after yatafting UDS
- Must latch UDS Poter for in the stack
- Must erase stackrøpfogeaphicalingSecure
 - Side-Channel Resistant
 - UDS Unique Device Secret
 - CDI Compound Device Identifier



UDS Unique Device Secret



UDS Unique Device Secret



UDS Unique Device Secret

let dice_main () : Stack dice_return_code (requires λ h → uds_is_enabled h) (ensures λ h₀ r h₁ → (¬ (uds_is_enabled h₁)) ∧ stack_is_erased h₁ ∧ all_heap_buffers_except_cdi_and_ghost_state_remain_same h₀ h₁ ∧ (r == DICE_SUCCESS ⇔ (10_image_is_valid (st ()).l₀ h₁ ∧ cdi_functional_correctness (st ()) h₁)))

Low^{*} allows us to specify the following properties about memory:

- All *heap* buffers,
- which were *alive* at the initial state h_0 ,
- and are *disjoint* with the *CDI* buffer,
- are still *alive* at the final state h_1
- and are *not modified*.

UDS Unique Device Secret

Verifying **DICE*** Engine: Side-Channel Resistance

- DICE* follows the secret independent coding discipline by reusing the secret integer model from HACL*
- **HACL**^{*} defines secrets as abstract, constant-time integers which
 - can not be used as array indexes
 - can not be branched on because no Boolean comparison operators for them

DICE* Layer 0

Derive Asymmetric Key Pairs

 $DeviceID_{pub}, DeviceID_{priv} = KDF(CDI_{L0})$

 $AliasKey_{pub}, AliasKey_{priv} = KDF(CDI_{L0}, L1)$

Generate Certificates

 $CSR_{DeviceID} = Sign(CreateCSR(DeviceID_{pub}), DeviceID_{priv})$ $Crt_{AliasKey} = Sign(CreateCrt(AliasKey_{pub}), DeviceID_{priv})$

KDF Key Derivation Function*CSR* Certificate Signing Request*Crt* Certificate

DICE* Layer 0

Generate Certificates

$$CSR_{DeviceID} = Sign(CreateCSR(DeviceID_{pub}), DeviceID_{priv})$$

$$Crt_{AliasKey} = Sign(CreateCrt(AliasKey_{pub}), DeviceID_{priv})$$

X.509 Certificate Generation Library provides verified serializer primitives and combinators for

- (Most of) ASN.1 constructs
- (A fragment of) X.509 messages

- *KDF* Key Derivation Function
- CSR Certificate Signing Request
- Crt Certificate

DICE* Layer 0

Derive Asymmetric Key Pairs

 $DeviceID_{pub}, DeviceID_{priv} = KDF(CDI_{L0})$ $AliasKey_{pub}, AliasKey_{priv} = KDF(CDI_{L0}, L_{1})$

Generate Certificates

 $CSR_{DeviceID} = Sign(CreateCSR(DeviceID_{pub}), DeviceID_{priv})$ $Crt_{AliasKey} = Sign(CreateCrt(AliasKey_{pub}), DeviceID_{priv})$



KDF Key Derivation Function*CSR* Certificate Signing Request*Crt* Certificate

- We reuse the secure parser and serializer model from LowParse for specification
- We verify properties such as our serializers are injective

$$\forall m_1, m_2. s(m_1) = s(m_2) \Rightarrow m_1 = m_2$$

• X.509 certificates are encoded into ASN.1

• X.509 certificates are encoded into ASN.1 Tag-

Tag

• X.509 certificates are encoded into ASN.1 Tag-Length-

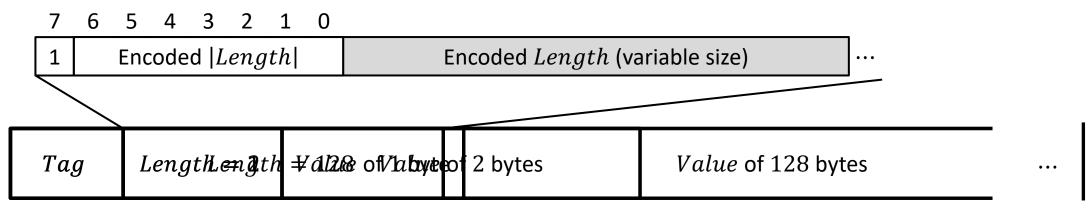
Tag

• X.509 certificates are encoded into ASN.1 Tag-Length-Value (TLV) format

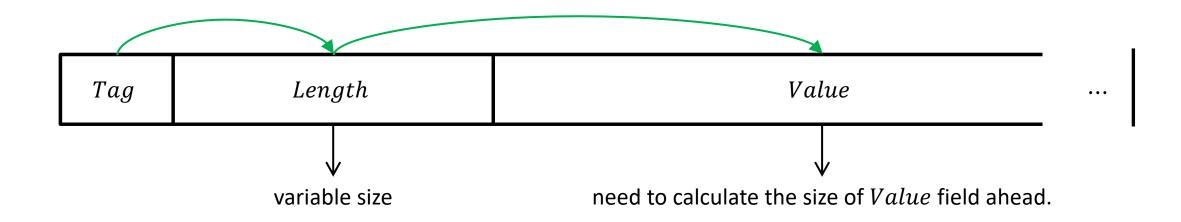
Tag	Length	Value
-----	--------	-------

- X.509 certificates are encoded into ASN.1 Tag-Length-Value (TLV) format
 - where the length field specifies the size of the value field
- But the length field is also variable size!

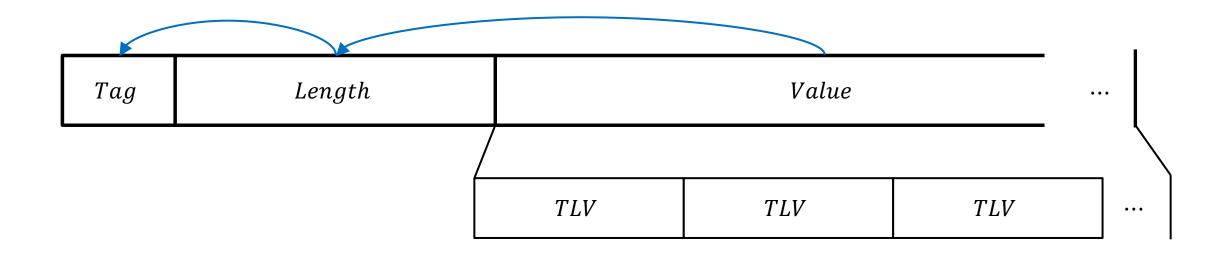
When $Length \ge 128$



- The low-level forward serializer from LowParse needs to calculate the size of value field ahead.
- Hence needs multiple passes to serialize an ASN.1 message, which is inefficient.



- We implement a verified low-level backward serializer, which serializes an ASN.1 message in one pass
 - even in the presence of nested TLV messages.



DICE* Implementation



7,677 **F*** LoC 5,051 **C** LoC X.509 Certificate Generation Library

16,564 **F*** LoC

DICE* Engine

533 **F*** LoC 205 **C** LoC

Platform-Agnostic Interface ~25k lines of \mathbf{F}^* code and proof ~5K lines of generated \mathbf{C} code

DICE* Implementation

dice_return_code dice_main()

```
HWState_state s = st();
bool b = authenticate_10_image(s.10);
dice_return_code r;
if (b)
 KRML_CHECK_SIZE(sizeof (uint8_t), uds_len);
 uint8_t uds[uds_len];
  memset(uds, 0U, uds_len * sizeof (uint8_t));
 read_uds(uds);
 uint8_t uds_digest[32U];
  memset(uds_digest, 0U, (uint32_t)32U * sizeof (uint8_t));
 uint8_t 10_digest[32U];
  memset(10_digest, 0U, (uint32_t)32U * sizeof (uint8_t));
 Hacl_Hash_SHA2_hash_256(uds, uds_len, uds_digest);
 Hacl_Hash_SHA2_hash_256(s.10.10_binary, s.10.10_binary_size, 10_digest);
  Hacl_HMAC_compute_sha2_256(s.cdi, uds_digest, (uint32_t)32U, 10_digest, (uint32_t)32U);
  zeroize(uds_len, uds);
  r = DICE_SUCCESS;
else
  r = DICE_ERROR;
disable_uds();
platform_zeroize_stack();
return r;
```

https://github.com/verified-HRoT/dice-star/tree/main/dist



 We show that the C implementation generated from DICE* is comparable to the unverified hand-written one



 We show that the C implementation generated from DICE* is comparable to the unverified hand-written one in terms of both boot time

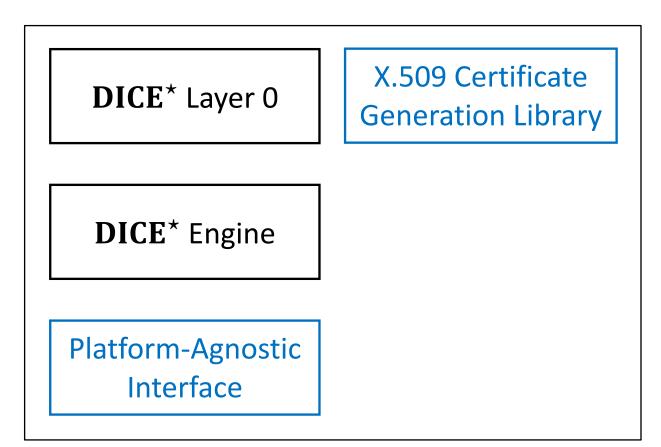
Lover	Boot time (ms)			
Layer	Unverified DICE	DICE*		
DICE Engine	786	689		
DICE Layer 0	313	208		



 We show that the C implementation generated from DICE* is comparable to the unverified hand-written one in terms of both boot time and binary size.

Layer	Boot time (ms)		Size (KB)	
	Unverified DICE	DICE*	Unverified DICE	DICE*
DICE Engine	786	689	72	68
DICE Layer 0	313	208	92	92

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Thank you!

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https://github.com/verified-HRoT/dice-star

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