#### **USENIX Security '24**

### Query Recovery from Easy to Hard: Jigsaw Attack against SSE

Hao Nie, Wei Wang, Peng Xu, Xianglong Zhang, Laurence T. Yang, and Kaitai Liang\*

Huazhong University of Science and Technology \*Delft University of Technology



I. Motivations











• Search Pattern, which indicates whether two queries are identical.





### **Previous Similar-data Attacks**



•Liu et al. [LZWT14] use the query frequency (from the search pattern) to match queries with keywords.



### Previous Similar-data Attacks



- •Liu et al. [LZWT14] use the query frequency (from the search pattern) to match queries with keywords.
- •Pouliot et al. [PW16], Damie et al. [DHP21], and Oya et al. [OK23] use the query co-occurrence to match queries with keywords.
  - The query co-occurrence is the probability of two queries shown in the same document. It could be deduced from the search pattern and access pattern.



Co-occurrence of queries from the leakage

Co-occurrence of keywords in the similar-data

## **Our Observations**



- •1. A small number of cracked queries can pose a significant threat to the security of other queries.
  - Damie et al. [DHP21] proposed the refined score attack that achieves around 85% accuracy in recovering all queries by utilizing only 10 known queries.

## **Our Observations**



- •1. A small number of cracked queries can pose a significant threat to the security of other queries.
- •2. Queries with a high volume/frequency are much easier to recover than others.
  - In a database, the volume and frequency of keywords follows Zipf's law.
  - Queries with higher volume or frequency display larger disparities, which consequently makes it easier for attackers to recover those queries.

### **Our Observations**



•A simple attack, which just matches the queries with keywords that have the closest volume and frequency, has 75% accuracy on the HVHF quadrant.



**I**. Our attack

### Jigsaw - Module 1



#### •Identify and recover the distinctive queries:

- Calculate the distance between all queries and their nearest neighbors, and select the first BaseRec queries with biggest distance as the distinctive queries.
- Match the BaseRec queries to the keywords that have the closest volume and frequency.



### Jigsaw - Module 2



#### •Remove some ill-matched queries:

• We check whether the results of module 1 is good or not. The good ones should also match in the co-occurrence relations. We keep ConfRec matched queries in this module.



### Jigsaw - Module 3



#### •Recover all queries based on the output of module 2:



### Jigsaw – Experimental Results





Figure: Jigsaw vs RSA[DHP21] vs IHOP[OK23] vs Sap[OK21] vs Graphm[PW16] in accuracy and runtime.

### Jigsaw – Experimental Results





Figure: Jigsaw vs IHOP in accuracy with the same time limits.

### **Against Countermeasures**





Figure: Jigsaw vs RSA vs IHOP in accuracy against the padding in [CGPR15].

•We pad the attacker's database with the same method as the client to minimizing the disparity between the similar data and the padded data.

### **Against Countermeasures**





Figure: Jigsaw vs RSA vs IHOP in accuracy against the obfuscation in [CLRZ18].

•We use the similar adaptation as [OK23] to all the attacks.

**III.** Conclusion

### Conclusion



- •We propose a new similar-data attack, Jigsaw. Some distinctive queries could threaten the whole system due to an attack like Jigsaw.
- •Jigsaw could bypass some countermeasures and still has high accuracy due to that the countermeasures do not protect the distinctive queries well.
- •An effective defense should hide the distinctive queries.

# Thank you for listening!

Code available: https://github.com/JigsawAttack/JigsawAttack

Contact information:

- → niehao@hust.edu.cn
- → viviawangwei@hust.edu.cn
- → xupeng@hust.edu.cn