## GhostRace: Exploiting and Mitigating Speculative Race Conditions

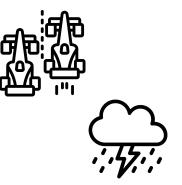
Hany Ragab, Andrea Mambretti, Anil Kurmus, and Cristiano Giuffrida



• Speculative Race Condition



- Speculative Race Condition
- Inter-Process Interrupt Storming



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- Inter-Process Interrupt Storming
- Gadget Scanner (1200+)



- Speculative Race Condition
- Inter-Process Interrupt Storming
- Gadget Scanner (1200+)
- PoC Exploit (12KB/s)

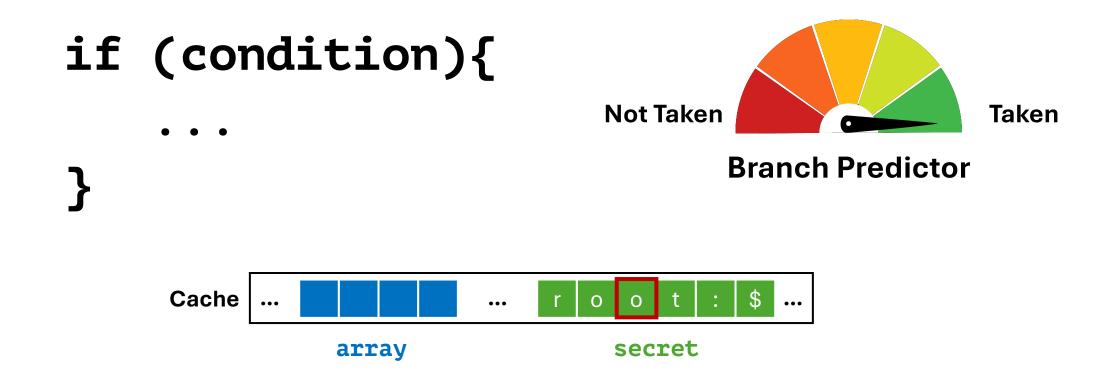


- Speculative Race Condition
- Inter-Process Interrupt Storming
- Gadget Scanner (1200+)
- PoC Exploit (12KB/s)
- Mitigation (5%)



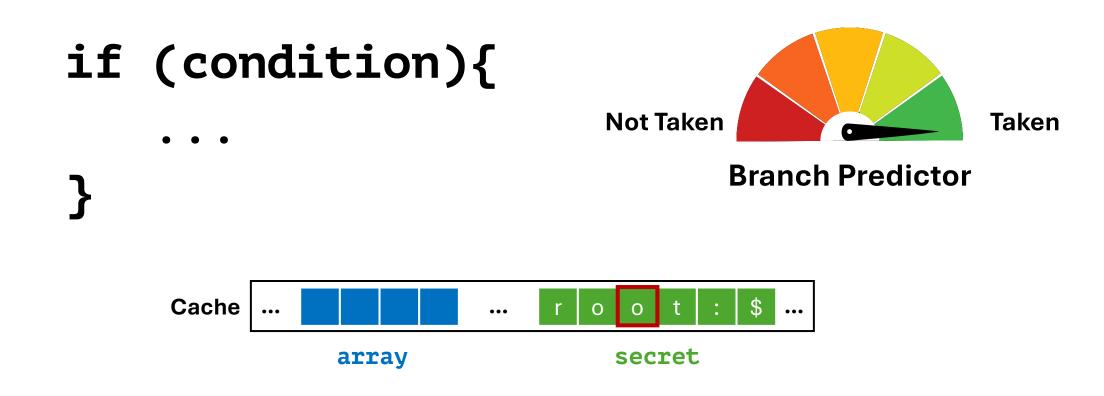


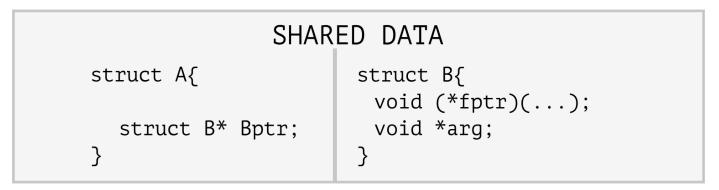
## Background **Spectre v1**

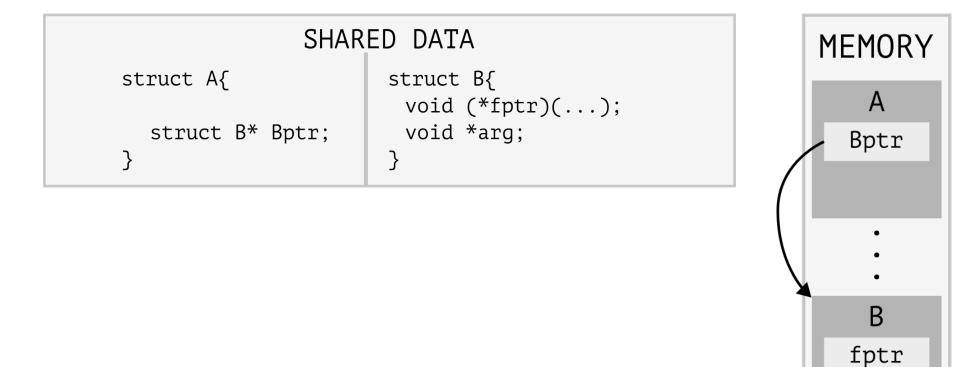


## Background **Spectre v1**

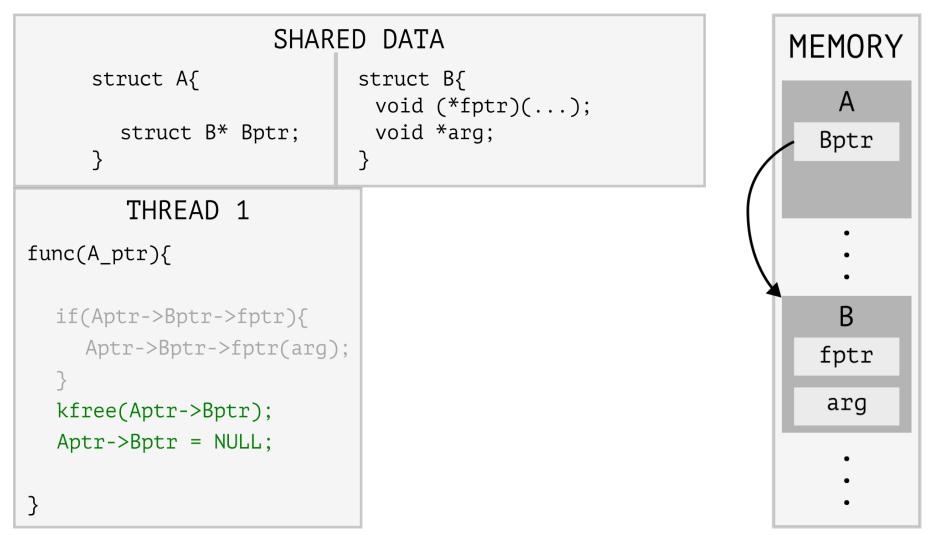
Conditional Branch Misprediction

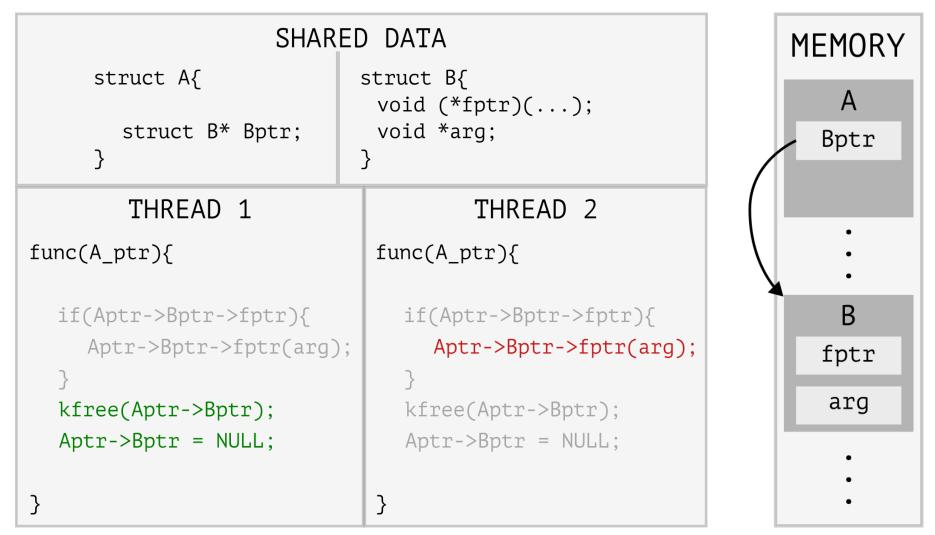


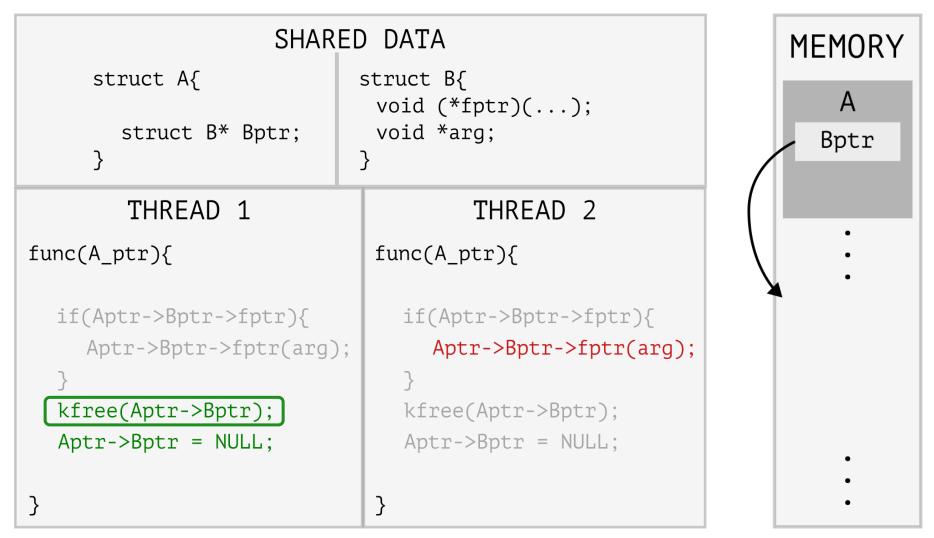


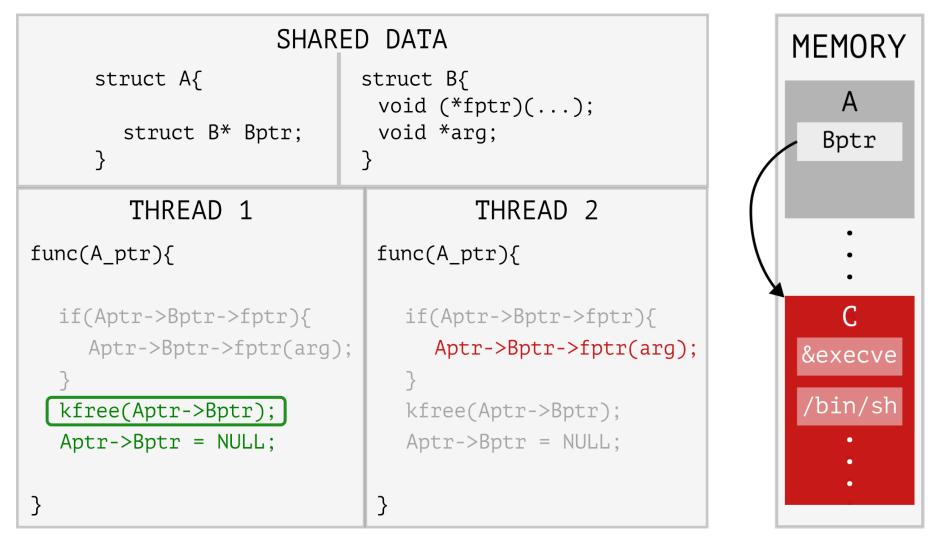


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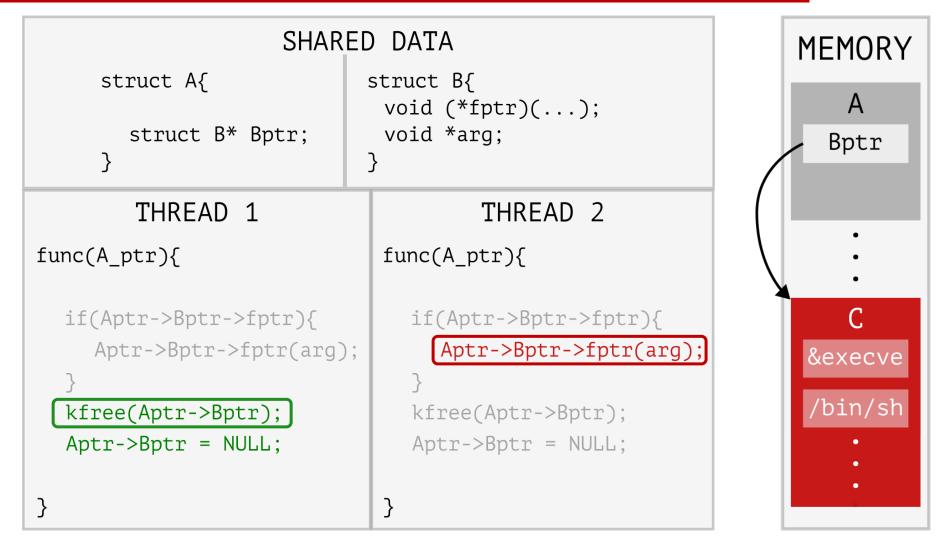




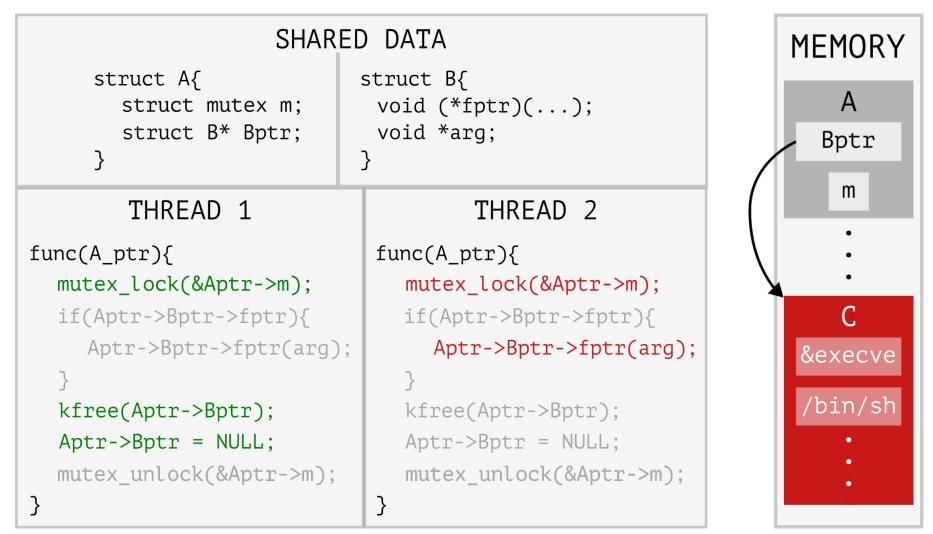


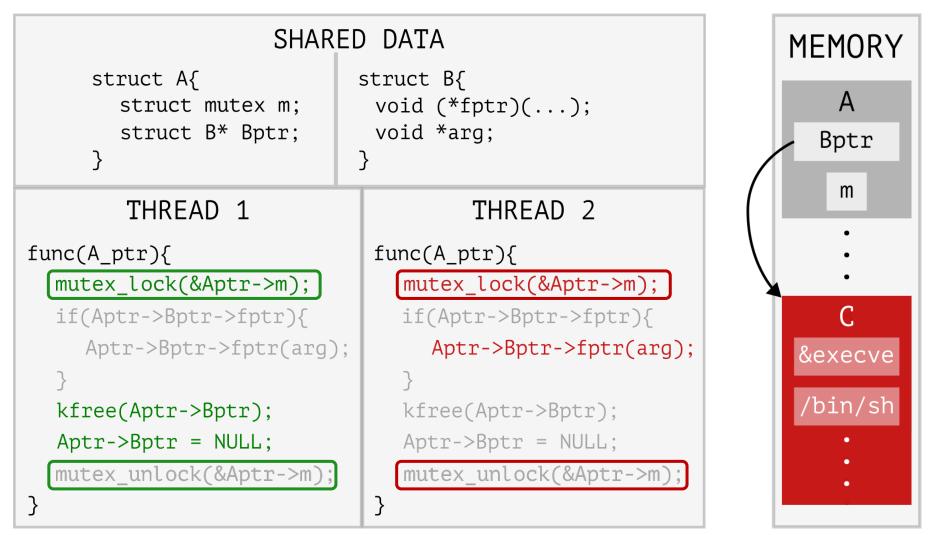


#### **Concurrent Use-After-Free Attack**

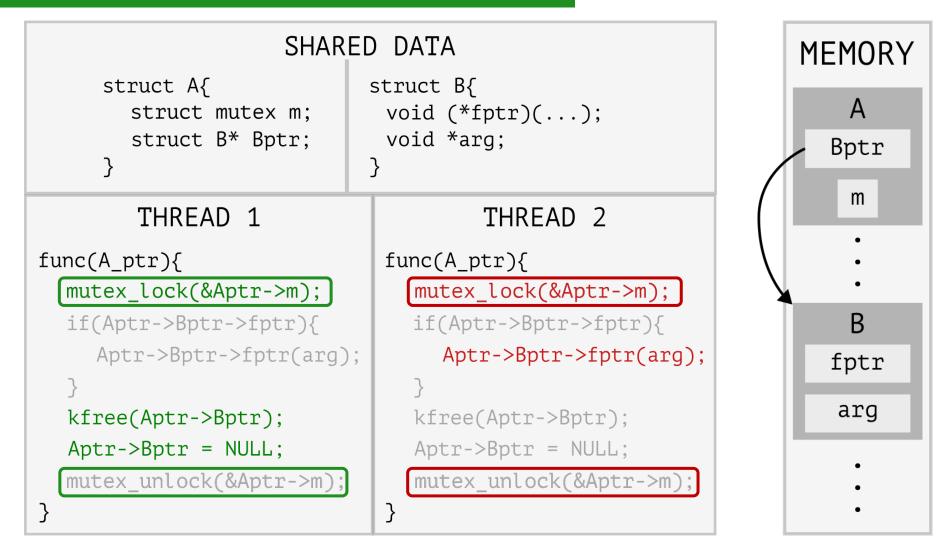


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#### **Race Free Critical Region**



## Real Example **NFC Gadget**

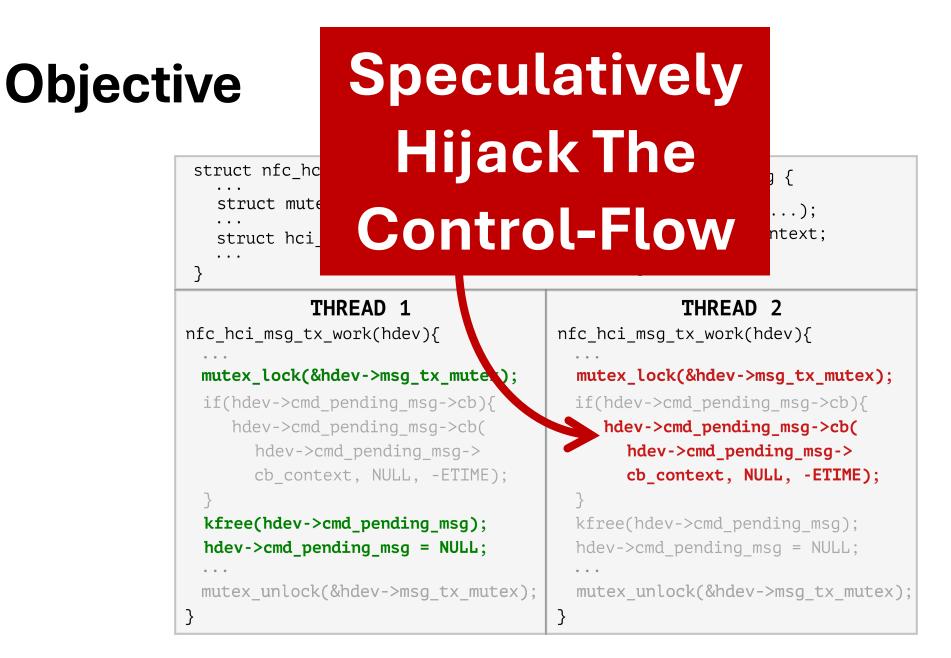
```
struct nfc hci dev {
                             SHARED DATA
                                              struct hci msg {
   struct mutex msg_tx_mutex;
                                                void (*cb)(...);
                                                void *cb context;
   struct hci msg* cmd pending msg;
                                                  THREAD 2
            THREAD 1
nfc hci msg_tx_work(hdev){
                                      nfc hci msg tx work(hdev){
 mutex_lock(&hdev->msg_tx_mutex);
                                        mutex_lock(&hdev->msg_tx_mutex);
 if(hdev->cmd pending msg->cb){
                                       if(hdev->cmd pending msg->cb){
    hdev->cmd pending msg->cb(
                                          hdev->cmd pending msg->cb(
      hdev->cmd pending msg->
                                             hdev->cmd pending msg->
      cb context, NULL, -ETIME);
                                             cb context, NULL, -ETIME);
 kfree(hdev->cmd pending msg);
                                       kfree(hdev->cmd pending msg);
 hdev->cmd pending msg = NULL;
                                       hdev->cmd pending msg = NULL;
 mutex unlock(&hdev->msg tx mutex);
                                       mutex unlock(&hdev->msg tx mutex);
}
                                      }
```

How do synchronization primitives behave during speculative execution? How do synchronization primitives behave during speculative execution?

What are the security implications for modern operating systems?

### Objective

struct nfc hci dev { SHARED DATA struct hci msg { struct mutex msg\_tx\_mutex; void (\*cb)(...); void \*cb context; struct hci msg\* cmd pending msg; **THREAD** 1 **THREAD** 2 nfc hci msg tx work(hdev){ nfc hci msg tx work(hdev){ mutex\_lock(&hdev->msg\_tx\_mutex); mutex\_lock(&hdev->msg\_tx\_mutex); if(hdev->cmd pending msg->cb){ if(hdev->cmd pending msg->cb){ hdev->cmd pending msg->cb( hdev->cmd pending msg->cb( hdev->cmd pending msg-> hdev->cmd pending msg-> cb context, NULL, -ETIME); cb context, NULL, -ETIME); kfree(hdev->cmd\_pending\_msg); kfree(hdev->cmd pending msg); hdev->cmd\_pending\_msg = NULL; hdev->cmd pending msg = NULL; mutex unlock(&hdev->msg tx mutex); mutex unlock(&hdev->msg tx mutex); } }



struct nfc hci dev { SHARED DATA struct hci msg { struct mutex msg\_tx\_mutex; void (\*cb)(...); void \*cb context; struct hci msg\* cmd pending msg; **THREAD** 2 **THREAD** 1 nfc hci msg\_tx\_work(hdev){ nfc hci msg tx work(hdev){ mutex\_lock(&hdev->msg\_tx\_mutex); mutex\_lock(&hdev->msg\_tx\_mutex); if(hdev->cmd pending msg->cb){ if(hdev->cmd pending msg->cb){ hdev->cmd pending msg->cb( hdev->cmd pending msg->cb( hdev->cmd pending msg-> hdev->cmd pending msg-> cb context, NULL, -ETIME); cb context, NULL, -ETIME); kfree(hdev->cmd pending msg); kfree(hdev->cmd pending msg); hdev->cmd pending msg = NULL; hdev->cmd pending msg = NULL; mutex unlock(&hdev->msg tx mutex); mutex unlock(&hdev->msg tx mutex); } }

**1. Bypass The** struct nfc hci dev { SHARED DATA struct hci msg { **Sync.** Primitive struct mutex msg\_tx\_mutex; void (\*cb)(...) void \*cb context struct hci msg\* cmd pending msg; **THREAD** 2 **THREAD** 1 nfc hci msg\_tx\_work(hdev){ nfc hci msg tx work(hdev){ mutex\_lock(&hdev->msg\_tx\_mutex); mutex\_lock(&hdev->msg\_tx\_mutex); if(hdev->cmd pending msg->cb){ if(hdev->cmd pending msg->cb){ hdev->cmd pending msg->cb( hdev->cmd pending msg->cb( hdev->cmd pending msg-> hdev->cmd pending msg-> cb context, NULL, -ETIME); cb context, NULL, -ETIME); kfree(hdev->cmd pending msg); kfree(hdev->cmd pending msg); hdev->cmd\_pending\_msg = NULL; hdev->cmd pending msg = NULL; mutex unlock(&hdev->msg tx mutex); mutex unlock(&hdev->msg tx mutex); } }

struct nfc\_hci\_dev { SHARED DATA struct hci\_msg {
 ...
 struct mutex msg\_tx\_mutex;
 ...
 struct hci\_msg\* cmd\_pending\_msg;
 ...
 }

# 2. Create an Exploitation Window

}

```
THREAD 2
           THREAD 1
                                    nfc hci msg_tx_work(hdev){
fc hci msg tx work(hdev){
mutex_lock(&hdev->msg_tx_mutex);
                                      mutex_lock(&hdev->msg_tx_mutex);
if(hdev->cmd pending msg->cb){
                                      if(hdev->cmd pending msg->cb){
   hdev->cmd pending msg->cb(
                                         hdev->cmd pending msg->cb(
     hdev->cmd pending msg->
                                            hdev->cmd pending msg->
     cb context, NULL, -ETIME);
                                            cb context, NULL, -ETIME);
kfree(hdev->cmd pending msg);
                                      kfree(hdev->cmd pending msg);
hdev->cmd_pending_msg = NULL;
                                      hdev->cmd pending msg = NULL;
mutex unlock(&hdev->msg tx mutex);
                                      mutex unlock(&hdev->msg tx mutex);
                                    }
```

1. Bypass The Sync. Primitive

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struct nfc\_hci
...
struct mute
...
freed Memory With
a Malicious Object

## 1. Bypass The Sync. Primitive

msg {

Þ)(...]

# 2. Create an Exploitation Window

<pre>struct hci_r } reed remory with _context, a Malicious Object</pre>	
<pre>fc_hci_msg_tx_work(hdev){</pre>	<pre>nfc_hci_msg_tx_work(hdev){</pre>
0 0 0	
<pre>mutex_lock(&amp;hdev-&gt;msg_tx_mutex);</pre>	<pre>mutex_lock(&amp;hdev-&gt;msg_tx_mutex);</pre>
if(hdev->cmd_pending_msg->cb){	if(hdev->cmd_pending_msg->cb){
hdev->cmd_pending_msg->cb(	hdev->cmd_pending_msg->cb(
hdev->cmd_pending_msg->	<pre>hdev-&gt;cmd_pending_msg-&gt;</pre>
<pre>cb_context, NULL, -ETIME);</pre>	<pre>cb_context, NULL, -ETIME);</pre>
}	}
<pre>kfree(hdev-&gt;cmd_pending_msg);</pre>	<pre>kfree(hdev-&gt;cmd_pending_msg);</pre>
<pre>hdev-&gt;cmd_pending_msg = NULL;</pre>	hdev->cmd_pending_msg = NULL;
• • •	• • •
<pre>mutex_unlock(&amp;hdev-&gt;msg_tx_mutex);</pre>	<pre>mutex_unlock(&amp;hdev-&gt;msg_tx_mutex);</pre>
}	}

struct nfc\_hci
...
struct mute
...
}
3. Reallocate The
Freed Memory With
a Malicious Object

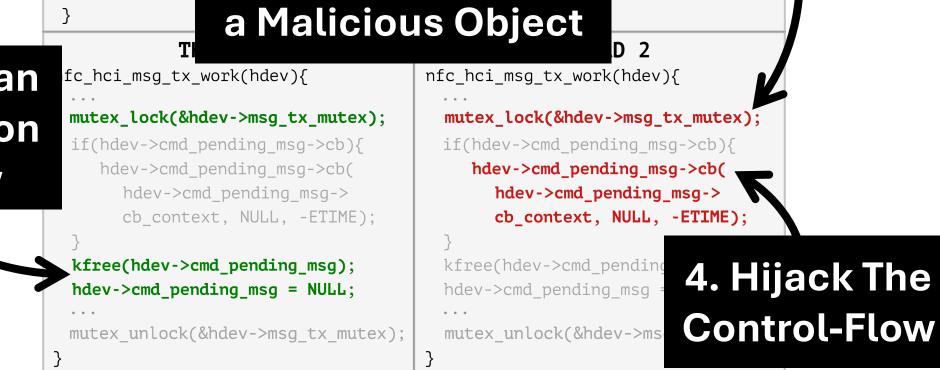
## 1. Bypass The Sync. Primitive

msg {

þ)(...)

contex

# 2. Create an Exploitation Window



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struct nfc\_hci\_dev { SHARED DATA struct hci\_msg {
 ...
 struct mutex msg\_tx\_mutex;
 ...
 struct hci\_msg\* cmd\_pending\_msg;
 ...
 }

# 2. Create an Exploitation Window

}

```
THREAD 2
           THREAD 1
                                    nfc hci msg_tx_work(hdev){
fc hci msg tx work(hdev){
mutex_lock(&hdev->msg_tx_mutex);
                                      mutex_lock(&hdev->msg_tx_mutex);
if(hdev->cmd pending msg->cb){
                                      if(hdev->cmd pending msg->cb){
   hdev->cmd pending msg->cb(
                                         hdev->cmd pending msg->cb(
     hdev->cmd pending msg->
                                            hdev->cmd pending msg->
     cb context, NULL, -ETIME);
                                            cb context, NULL, -ETIME);
kfree(hdev->cmd pending msg);
                                      kfree(hdev->cmd pending msg);
hdev->cmd_pending_msg = NULL;
                                      hdev->cmd pending msg = NULL;
mutex unlock(&hdev->msg tx mutex);
                                      mutex unlock(&hdev->msg tx mutex);
                                    }
```

#### 1. Bypass The Sync. Primitive

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**1. Bypass The** struct nfc hci dev { SHARED DATA struct hci msg { **Sync.** Primitive struct mutex msg\_tx\_mutex; void (\*cb)(...) void \*cb context struct hci msg\* cmd pending msg; **THREAD** 2 **THREAD** 1 nfc hci msg\_tx\_work(hdev){ nfc hci msg tx work(hdev){ mutex\_lock(&hdev->msg\_tx\_mutex); mutex\_lock(&hdev->msg\_tx\_mutex); if(hdev->cmd pending msg->cb){ if(hdev->cmd pending msg->cb){ hdev->cmd pending msg->cb( hdev->cmd pending msg->cb( hdev->cmd pending msg-> hdev->cmd pending msg-> cb context, NULL, -ETIME); cb context, NULL, -ETIME); kfree(hdev->cmd pending msg); kfree(hdev->cmd pending msg); hdev->cmd\_pending\_msg = NULL; hdev->cmd pending msg = NULL; mutex unlock(&hdev->msg tx mutex); mutex unlock(&hdev->msg tx mutex); } }

#### **Synchronization Primitives**

### **Synchronization Primitives**

Mutex lock on x86 arch:

• Conditional branch

```
void mutex_lock(struct mutex *lock){
  if (!__mutex_trylock_fast(lock))
    if (atomic_long_try_cmpxchg_acquire(&lock, ...))
```

### **Synchronization Primitives**

Mutex lock on x86 arch:

- Conditional branch
- A series of inlined function calls

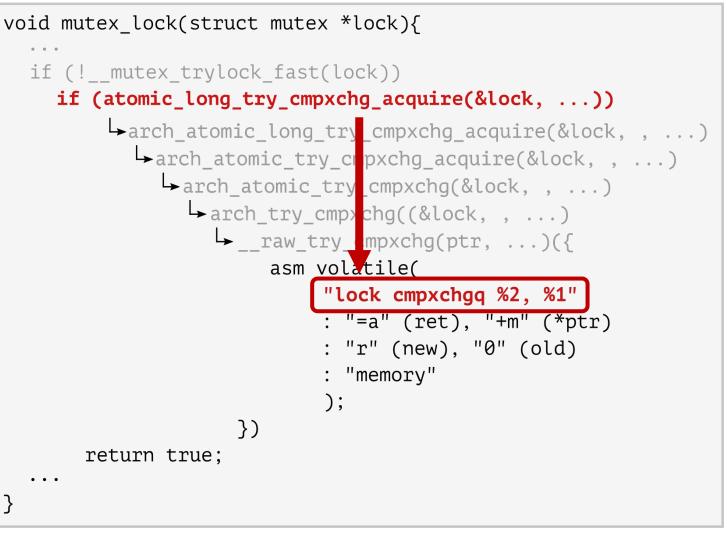
Mutex lock on x86 arch:

- Conditional branch
- A series of inlined function calls
- Atomic compare and exchange

```
void mutex lock(struct mutex *lock){
  if (! mutex trylock fast(lock))
    if (atomic_long_try_cmpxchg_acquire(&lock, ...))
         arch_atomic_long_try_cmpxchg_acquire(&lock, , ...)
           →arch_atomic_try_cmpxchg_acquire(&lock, , ...)
              \rightarrow arch atomic try cmpxchg(&lock, , ...)
                → arch_try_cmpxchg((&lock, , ...)
                  → raw try cmpxchg(ptr, ...)({
                       asm volatile(
                            "lock cmpxchgq %2, %1"
                            : "=a" (ret), "+m" (*ptr)
                            : "r" (new), "0" (old)
                            : "memory"
                            );
                    })
       return true:
```

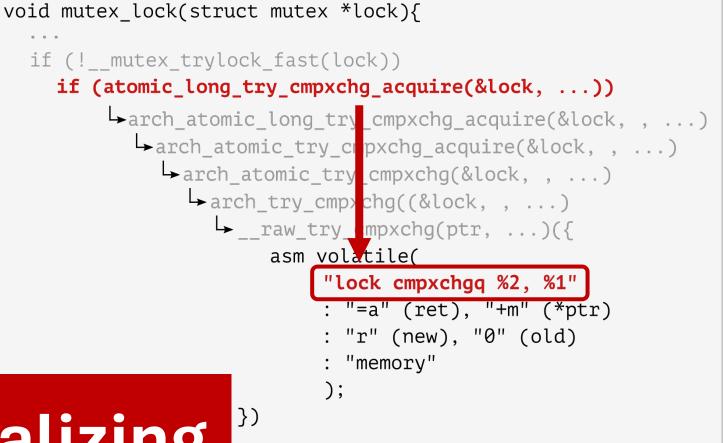
Mutex lock on x86 arch:

- Conditional branch
- A series of inlined function calls
- Atomic compare and exchange

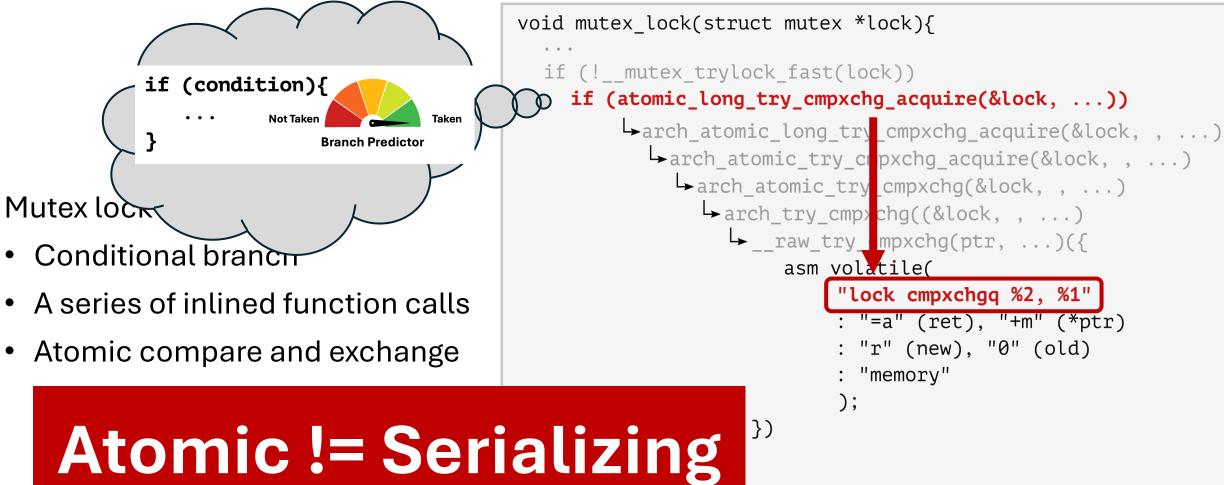


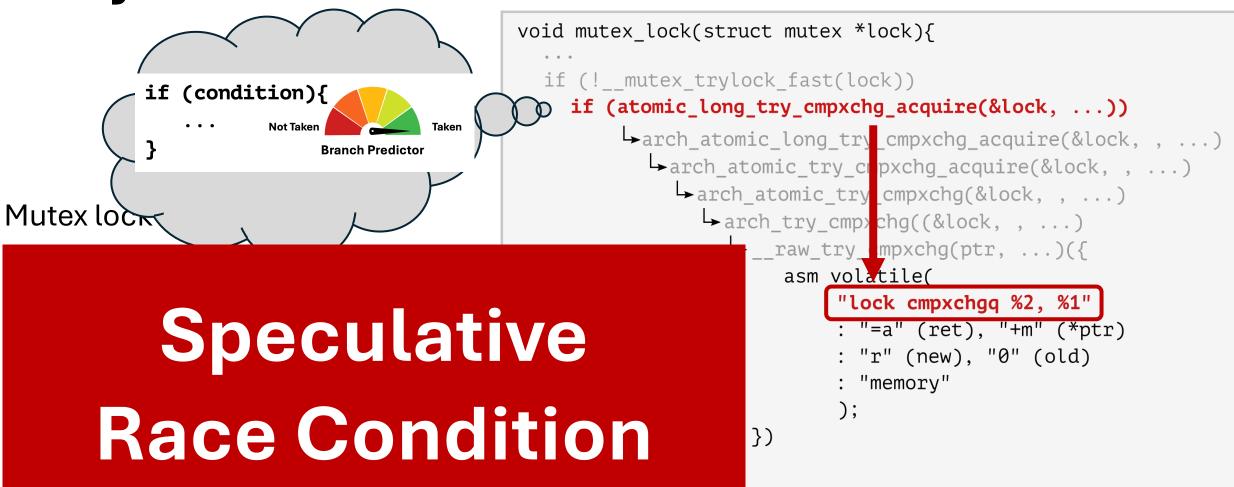
Mutex lock on x86 arch:

- Conditional branch
- A series of inlined function calls
- Atomic compare and exchange



## Atomic != Serializing





#### **UAF Attack Challenges**

**Speculative Race** struct nfc hci dev { SHARED DATA struct hci r Condition struct mutex msg\_tx\_mutex; void (\*cb void \*cb struct hci msg\* cmd pending msg; **THREAD** 2 **THREAD** 1 nfc hci msg tx work(hdev){ nfc hci msg\_tx\_work(hdev){ mutex\_lock(&hdev->msg\_tx\_mutex); mutex\_lock(&hdev->msg\_tx\_mutex); if(hdev->cmd pending msg->cb){ if(hdev->cmd pending msg->cb){ hdev->cmd pending msg->cb( hdev->cmd pending msg->cb( hdev->cmd pending msg-> hdev->cmd pending msg-> cb context, NULL, -ETIME); cb context, NULL, -ETIME); kfree(hdev->cmd pending msg); kfree(hdev->cmd pending msg); hdev->cmd\_pending\_msg = NULL; hdev->cmd pending msg = NULL; mutex unlock(&hdev->msg tx mutex); mutex unlock(&hdev->msg tx mutex); } }

**1. Bypass The Sync. Primitive:** 

#### **UAF Attack Challenges**

fc

mu

if

kf hd

mu

}

struct nfc\_hci\_dev { SHARED DATA struct hci\_r
...
struct mutex msg\_tx\_mutex; void (\*cb
...
struct hci\_msg\* cmd\_pending\_msg; void \*cb\_
...
}

# 2. Create an Exploitation Window

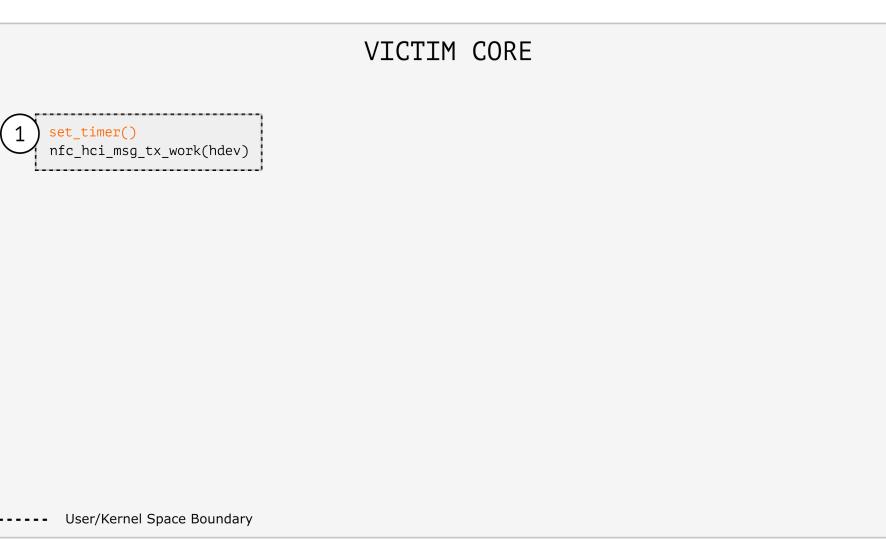
···	}	
THREAD 1	THREAD 2	
_hci_msg_tx_work(hdev){	<pre>nfc_hci_msg_tx_work(hdev){</pre>	
itex_lock(&hdev->msg_tx_mutex);	<pre>mutex_lock(&amp;hdev-&gt;msg_tx_mutex);</pre>	
f(hdev->cmd_pending_msg->cb){	if(hdev->cmd_pending_msg->cb){	
hdev->cmd_pending_msg->cb(	hdev->cmd_pending_msg->cb(	
hdev->cmd_pending_msg->	<pre>hdev-&gt;cmd_pending_msg-&gt;</pre>	
<pre>cb_context, NULL, -ETIME);</pre>	<pre>cb_context, NULL, -ETIME);</pre>	
	}	
<pre>free(hdev-&gt;cmd_pending_msg);</pre>	<pre>kfree(hdev-&gt;cmd_pending_msg);</pre>	
dev->cmd_pending_msg = NULL;	hdev->cmd_pending_msg = NULL;	
<pre>utex_unlock(&amp;hdev-&gt;msg_tx_mutex);</pre>	<pre>mutex_unlock(&amp;hdev-&gt;msg_tx_mutex)</pre>	
	}	

1. Bypass The Sync. Primitive: Speculative Race Condition

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VICTIM CORE

----- User/Kernel Space Boundary





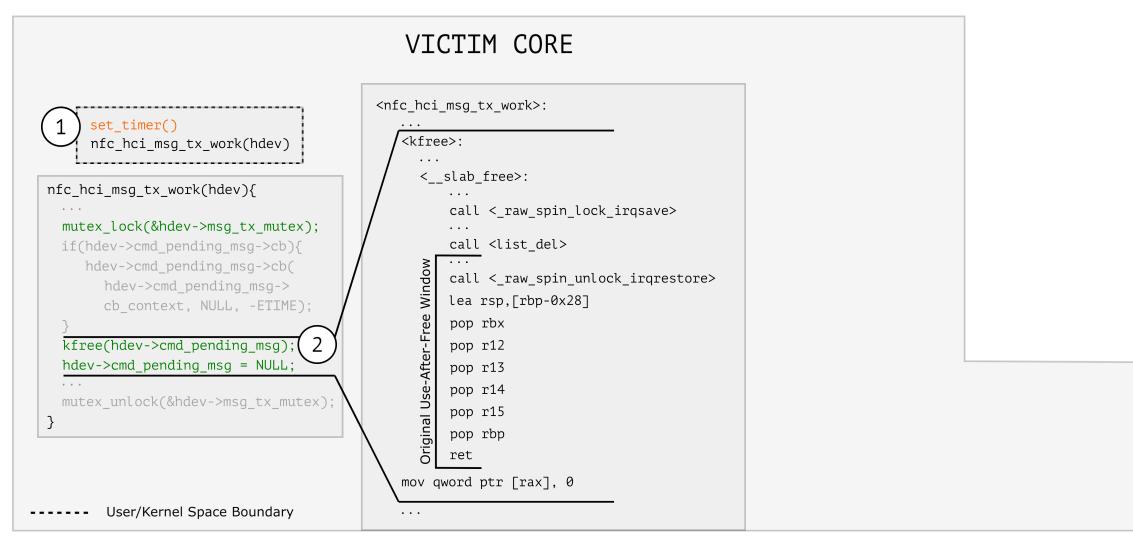
User/Kernel Space Boundary

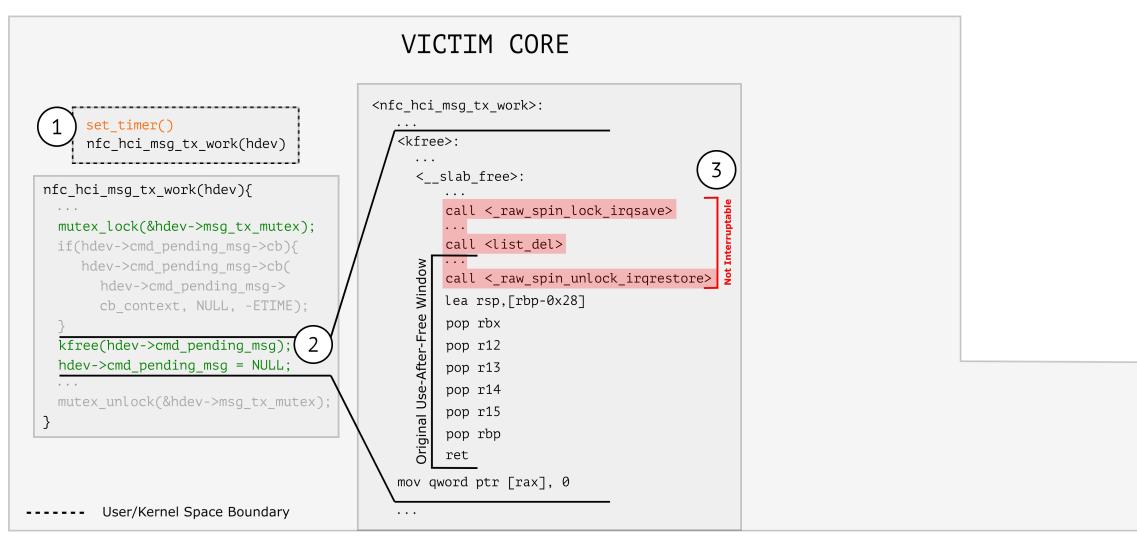
}

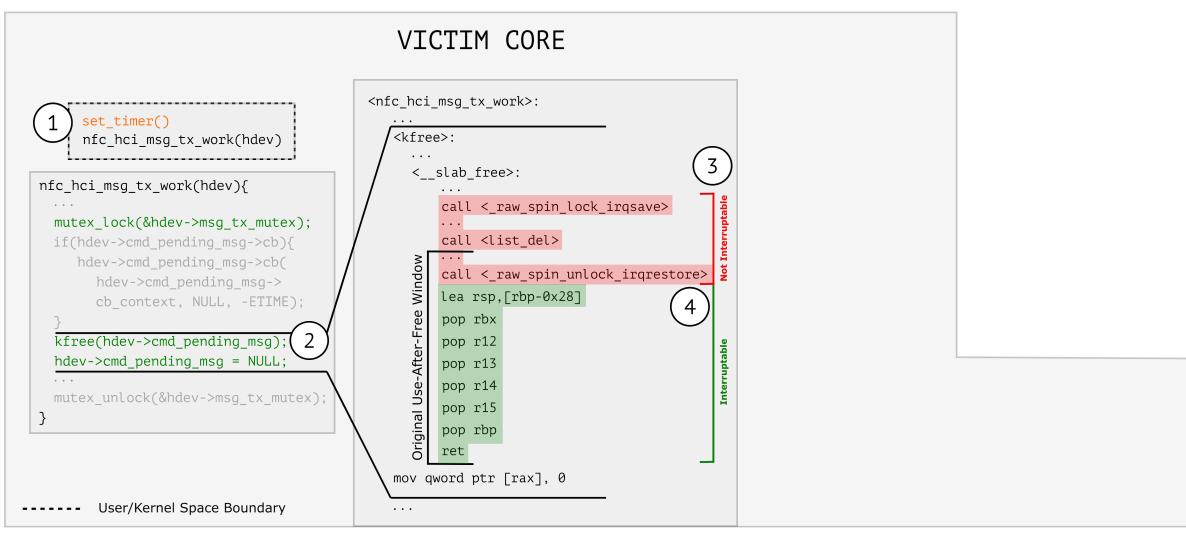


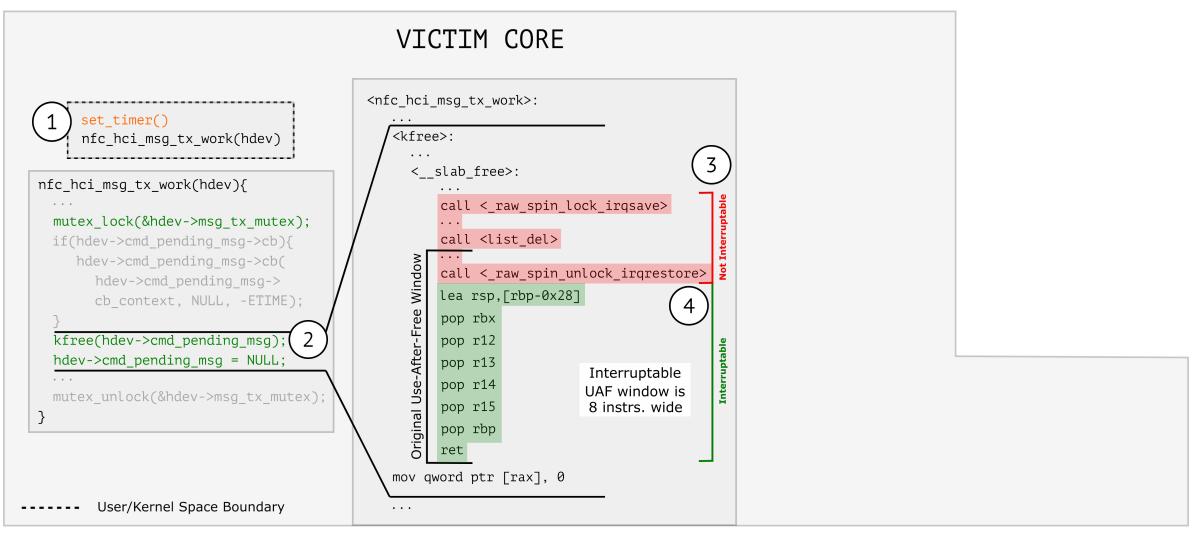
----- User/Kernel Space Boundary

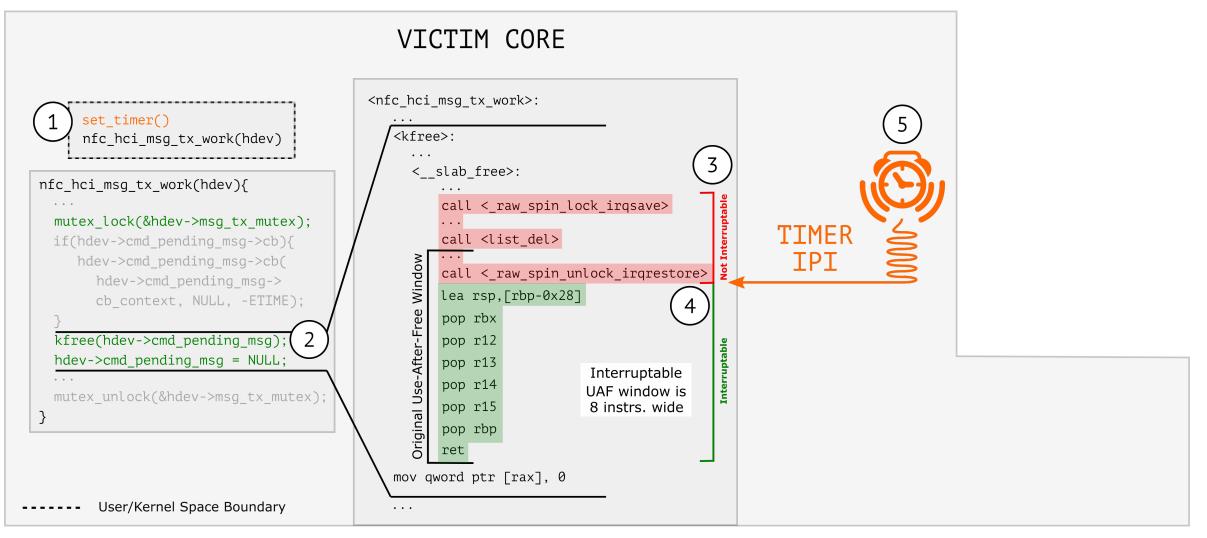
<pre> 1 set_timer()     nfc_hci_msg_tx_work(hdev)  nfc_hci_msg_tx_work(hdev)  nfc_hci_msg_tx_work(hdev){      mutex_lock(&amp;hdev-&gt;msg_tx_mutex);     if(hdev-&gt;cmd_pending_msg-&gt;cb(         hdev-&gt;cmd_pending_msg-&gt;cb(         hdev-&gt;cmd_pending_msg-&gt;(2)         hdev-&gt;cmd_pending_msg);     }     kfree(hdev-&gt;msg_tx_mutex);     }     mutex_unlock(&amp;hdev-&gt;msg_tx_mutex);     } </pre>

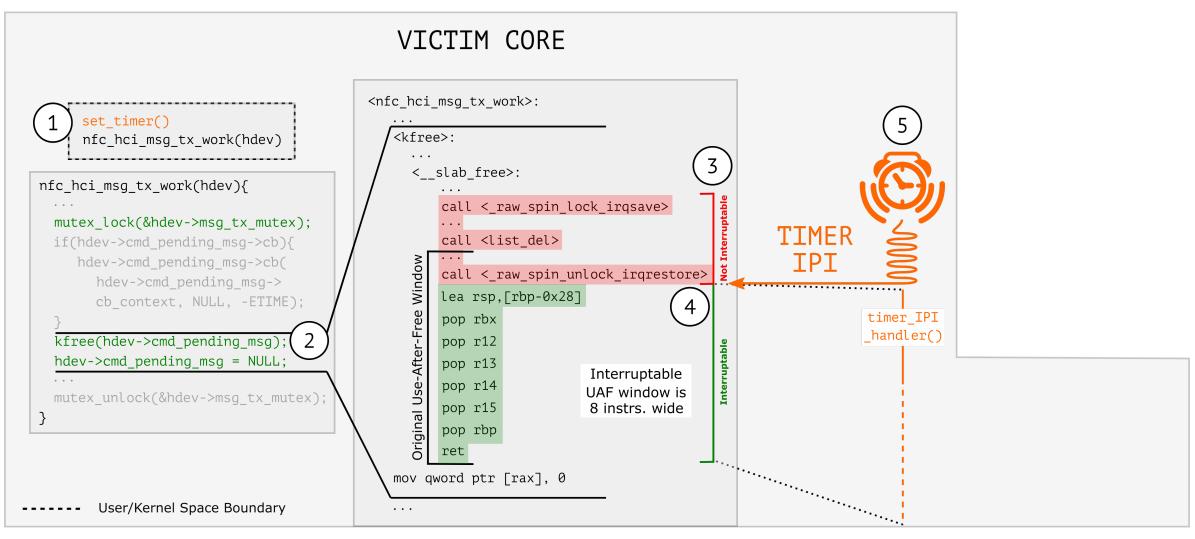


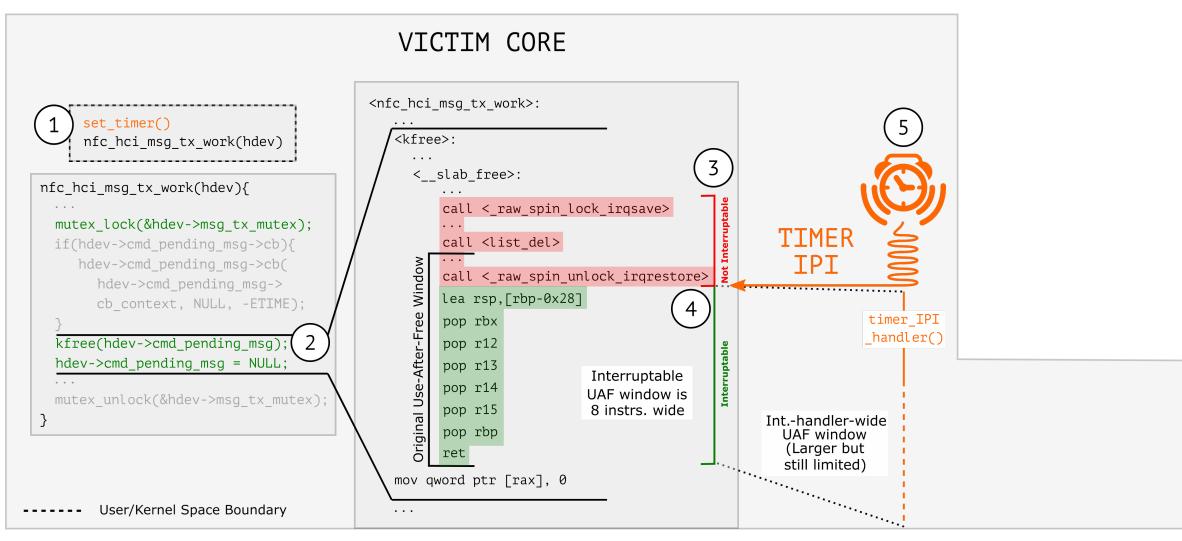


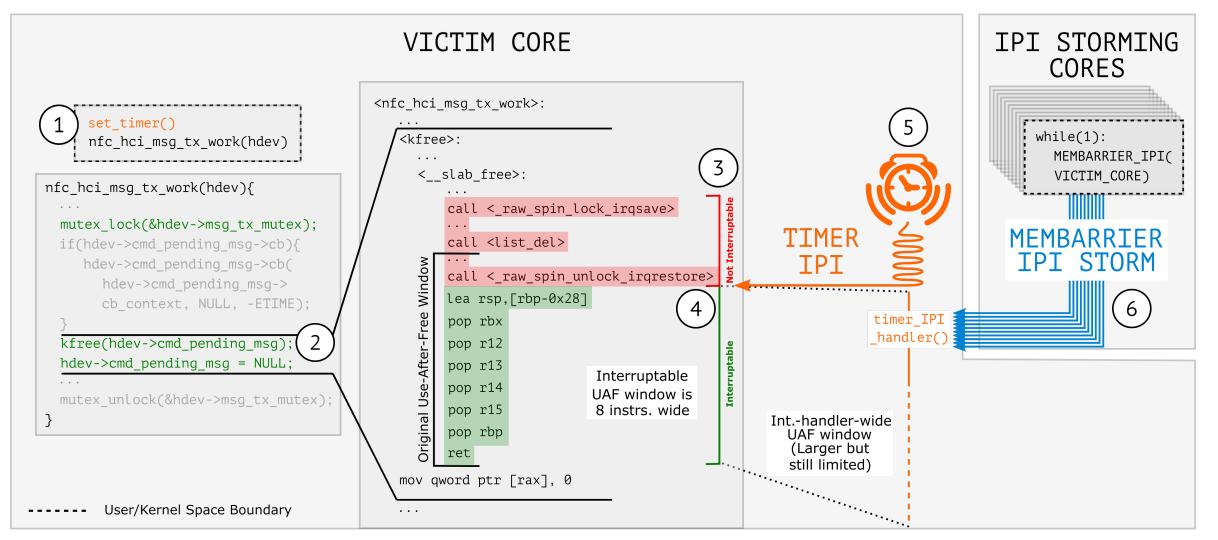


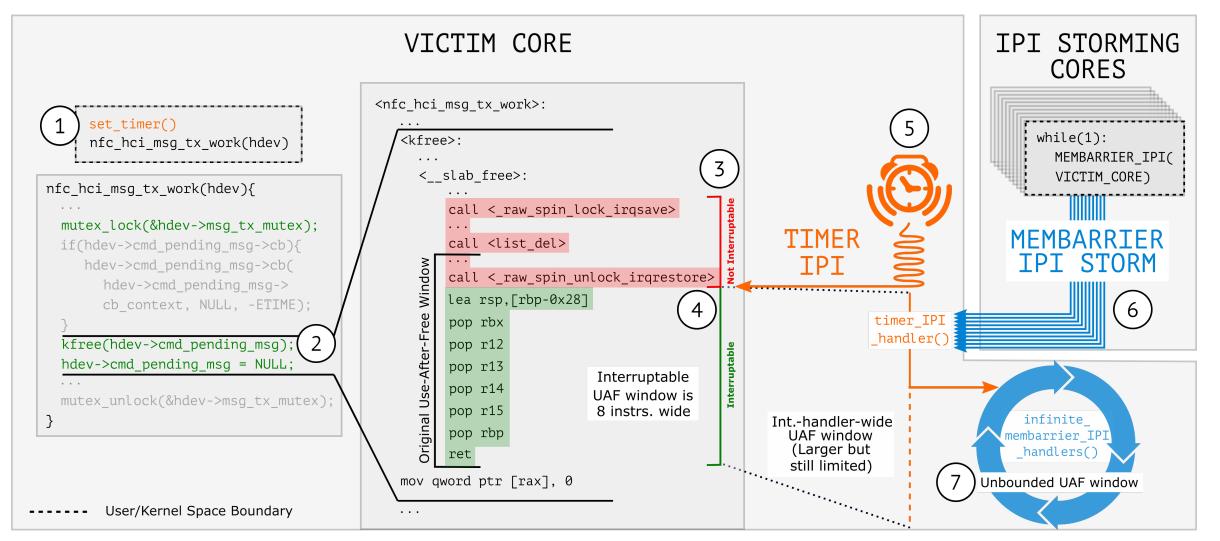


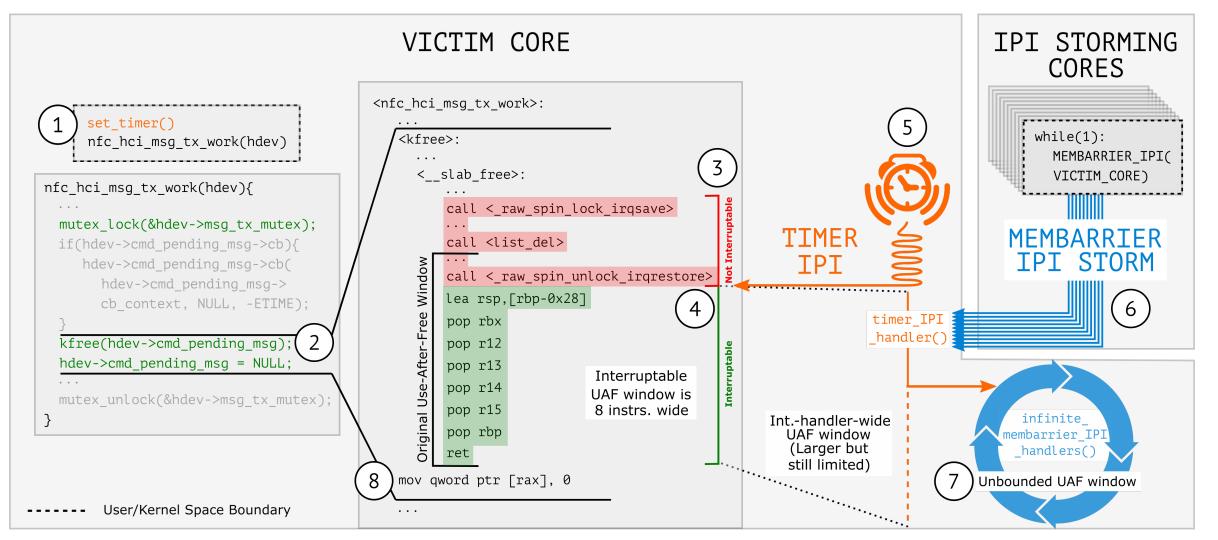


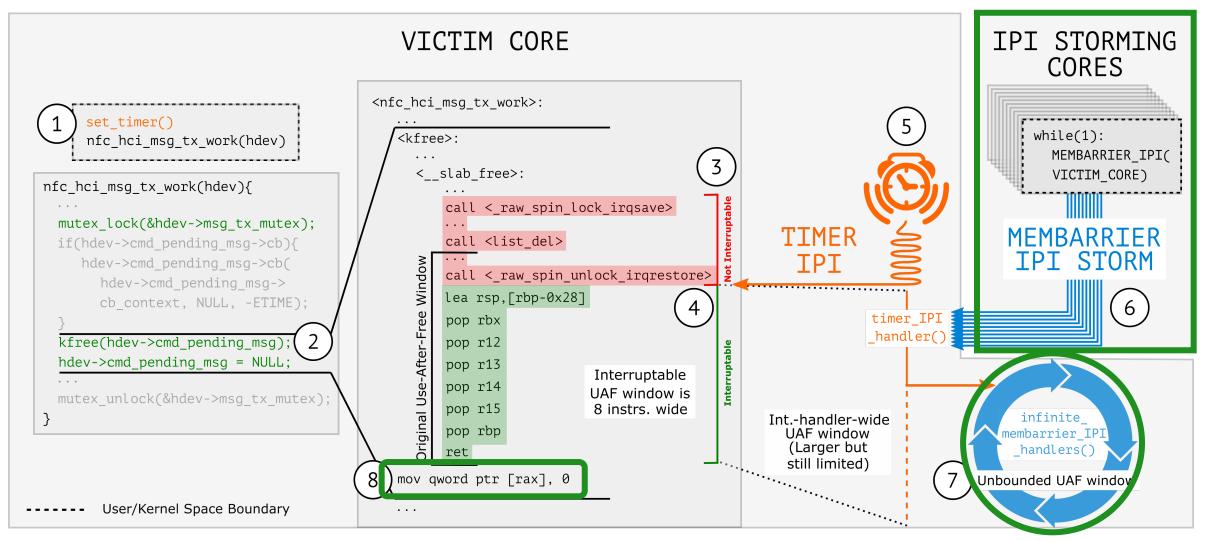












VICTIM CORE	IPI STORMING CORES
<pre> 1 se nfc_hci wutex if(hde hde } </pre> Cubounded Cubounde Cubound	while(1): MEMBARRIER_IPI( VICTIM_CORE) MEMBARRIER IPI STORM
hdev-i mutex } UAF Window (Larger but still limited) 7 User/Kernel Space Boundary 	<pre>infinite_ membarrier_IPI _handlers() Unbounded UAF window</pre>

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#### **UAF Attack Challenges**

struct nfc\_hci dev { SHARED DATA struct hci r struct mutex msg\_tx\_mutex; void (\*cb void \*cb struct hci msg\* cmd pending msg; **THREAD** 2 **THREAD** 1 i msg tx work(hdev){ nfc hci msg tx work(hdev){ k\_lock(&hdev->msg\_tx\_mutex); mutex\_lock(&hdev->msg\_tx\_mutex); dev->cmd pending msg->cb){ if(hdev->cmd pending msg->cb){ dev->cmd pending msg->cb( hdev->cmd pending msg->cb( hdev->cmd pending msg-> hdev->cmd pending msg-> cb context, NULL, -ETIME); cb context, NULL, -ETIME); kfree(hdev->cmd pending msg); kfree(hdev->cmd pending msg); hdev->cmd\_pending\_msg = NULL; hdev->cmd pending msg = NULL; mutex unlock(&hdev->msg tx mutex); mutex unlock(&hdev->msg tx mutex);

1. Bypass The Sync. Primitive: Speculative Race Condition

2. Create an Exploitation Window: IPI Storming

}

}

#### **UAF Attack Challenges**

. . .

struct nfc 3. Reallocate The Freed Memory: struct m msgsnd syscall struct h

#### **1. Bypass The Sync. Primitive: Speculative Race** Condition

2. Create an Exploitation Window: **IPI Storming** 

#### **THREAD** 1 i msg tx work(hdev){

#### k\_lock(&hdev->msg\_tx\_mutex);

dev->cmd pending msg->cb){ dev->cmd pending msg->cb( hdev->cmd pending msg-> cb context, NULL, -ETIME);

**THREAD** 2 nfc hci msg\_tx\_work(hdev){

t hci r

ld (\*cb

ld \*cb

mutex lock(&hdev->msq tx mutex); if(hdev->cmd pending msg->cb){ hdev->cmd pending msg->cb( hdev->cmd pending msg-> cb context, NULL, -ETIME);

kfree(hdev->cmd pending msc hdev->cmd\_pending\_msg = NUL mutex unlock(&hdev->msg tx }

4. Hijack The Control-Flow: **Speculatively Execute A** Controlled **Disclosure Gadget** 

GhostRace: Exploiting and Mitigating Spe

## Speculative Concurrent Use-After-Free Attack

struct nfc Struct m Struct h Struct h Struct h

#### 1. Bypass The Sync. Primitive: Speculative Race Condition

2. Create an Exploitation Window: IPI Storming **THREAD 1** i\_msg\_tx\_work(hdev){

#### k\_lock(&hdev->msg\_tx\_mutex);

lev->cmd\_pending\_msg->cb){
dev->cmd\_pending\_msg->cb(
 hdev->cmd\_pending\_msg->
 cb\_context, NULL, -ETIME);

THREAD 2
nfc\_hci\_msg\_tx\_work(hdev){

mutex\_lock(&hdev->msg\_tx\_mutex); if(hdev->cmd\_pending\_msg->cb){ hdev->cmd\_pending\_msg->cb( hdev->cmd\_pending\_msg-> cb\_context, NULL, -ETIME);

t hci r

ld (\*cb

ld \*cb

kfree(hdev->cmd\_pending\_msg
hdev->cmd\_pending\_msg = NUL
...
mutex\_unlock(&hdev->msg\_tx\_

4. Hijack The Control-Flow: Speculatively Execute A Controlled Disclosure Gadget

GhostRace: Exploiting and Mitigating Spe

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• Gadget scanner using the Coccinelle patching engine

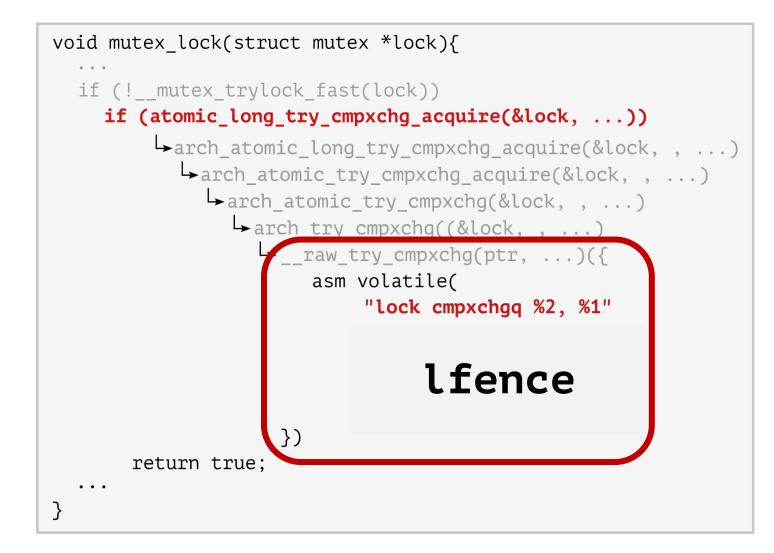
- Gadget scanner using the Coccinelle patching engine
- Statically analyzed Linux v5.15.83

- Gadget scanner using the Coccinelle patching engine
- Statically analyzed Linux v5.15.83

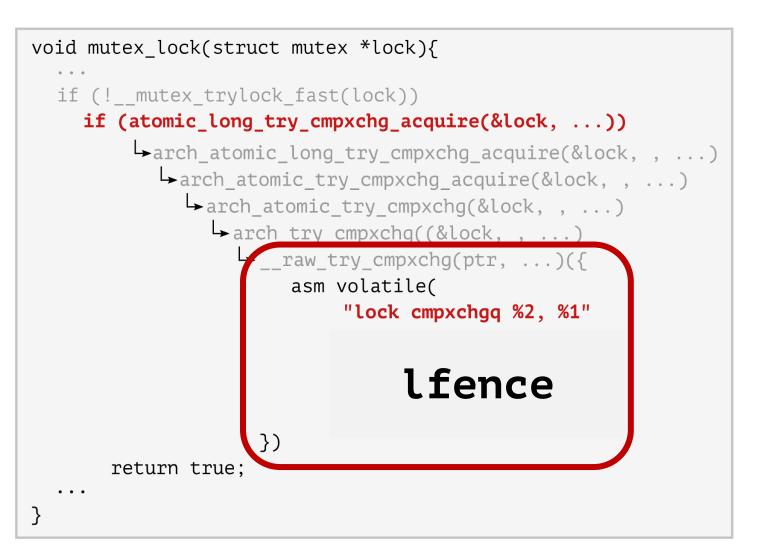
# Identified 1283 Speculative Concurrent Use-After-Free Gadgets

```
void mutex lock(struct mutex *lock){
  if (! mutex trylock fast(lock))
    if (atomic_long_try_cmpxchg_acquire(&lock, ...))
         arch_atomic_long_try_cmpxchg_acquire(&lock, , ...)
           →arch atomic try cmpxchg acquire(&lock, , ...)
             arch_atomic_try_cmpxchg(&lock, , ...)
                → arch try_cmpxchg((&lock, , ...)
                  → raw try cmpxchg(ptr, ...)({
                       asm volatile(
                           "lock cmpxchgq %2, %1"
                           : "=a" (ret), "+m" (*ptr)
                           : "r" (new), "0" (old)
                           : "memory"
                           );
                    })
       return true;
}
```





LMBench: Geomean **≈5%** 



• Speculative Race Condition



- Speculative Race Condition
- Inter-Process Interrupt Storming



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GhostRace: Exploiting and Mitigating Speculative Race Conditions | Hany Ragab @ USENIX Security 2024

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