

TAPFixer: Automatic Detection and Repair of Home Automation Vulnerabilities based on Negated-property Reasoning

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Home Automation (HA)

HA in daily life



Lighting Control



Security Monitoring



Smart Watering



Smart Cleaning

HA platforms



Samsung SmartThings











openHAB

Trigger-Action Programming (TAP)

Paradigm of a TAP rule

IF a [Trigger] occurs WHILE a [Condition] satisfies, THEN perform an [Action].

Example



Keep the room at the proper temperature:

IF [motion sensor detects user returning home], THEN [turn on heater].

[Trigger] [Action]



Power off before bedtime:

IF [11 p.m.] WHILE [motion sensor detects no one moving], THEN [turn off all power].

[Trigger] [Condition] [Action]

Interaction Vulnerabilities

TAP rule1 : IF [motion sensor detects user returning home] , THEN [turn on heater].

[Trigger] [Action]

TAP rule2 : IF [11 p.m.] WHILE [motion sensor detects no one moving] , THEN [turn off all power].

[Trigger] [Condition] [Action]

	Secure Case	Defective Case
Interaction Pattern	Rules run independently	Rules interact with each other
Execution Sequence	Rule2 triggers later than rule1 since the user is usually home before 11 p.m.	Rule2 triggers earlier than rule1 since the user is home after 11 p.m.
Vulnerability	Nonexistent	The heater may still be running while sleeping, which may cause a fire hazard .

Motivation: Previous Approaches

Vulnerability Detection:

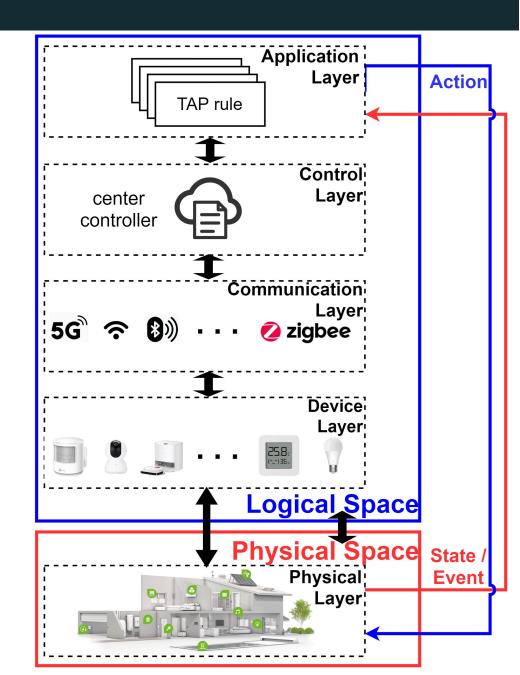
Model Checking-based / Symbolic Execution-based

Logical-physical Space:

- Determine correctness of rule interactions
- Latency l_1 , l_2 , l_3 :
- l_1 : delay defined in rules for specifying the device execution time
- l_2 : delay on a tardy attribute change to a certain value
- l_3 : platform delay
- Physical interaction phy_1, phy_2, phy_3 :
- phy_1 : implicit physical effect phy_2 : joint physical effect
- *phy*₃: nondeterminacy

Limitation:

- Existing works neglect above key logical-physical features
- Fail to detect vulnerabilities with such features



Motivation: Previous Approaches

Vulnerability Repair:

Dynamic Access Control-based / Static Semantic Modification-based

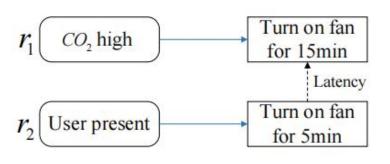
Limitation of Dynamic Approaches:

- Unable to fix flaws in rule semantics (root cause of vulnerabilities)
- Introduce additional runtime overhead

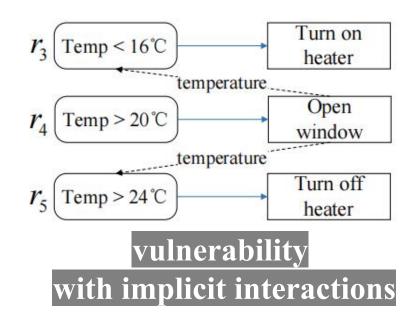
Limitation of Static Approaches:

• Not considering dynamic factors in physical space and fail to repair related expanded vulnerabilities.

Scenarios with Dynamic Factors

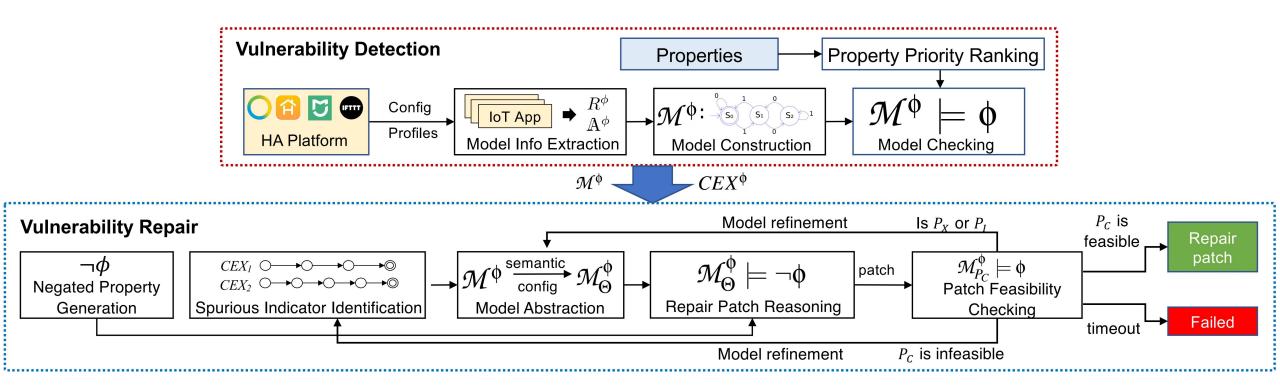


vulnerability with latency



TAPFixer

To our best knowledge, TAPFixer is the first work that can essentially detect and fix rule interaction vulnerabilities both in the logical and physical space.



Addressing Challenge 1: Comprehensive modeling logical-physical features

Physical model-based HA system modeling

1. Latency Modeling

 l_1 : modeled as a timer configured with a timeout value

 l_2 : set the physical changing threshold

 l_3 : set the updating interval threshold

3. Finite Automata Construction

$$\mathcal{M}_{RI}$$
: = { S , I , Σ }

Automata state universal set S, initial state set I:

modeled as device and environment attributes

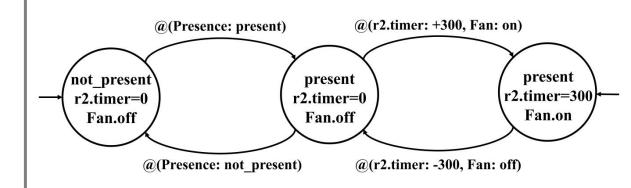
State transfer function $\phi \in \Sigma$:

transfer conditions: modeled as trigger and condition

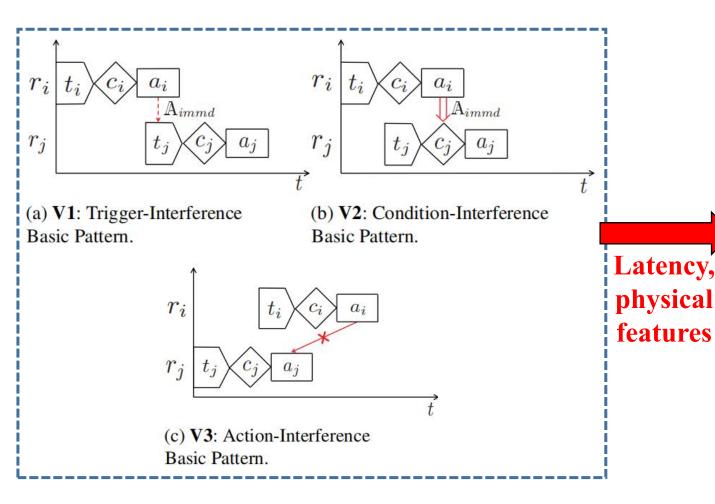
transfer label: modeled as action

2. Physical Interaction Modeling

phy₁: create a mapping of device actions to implicitly affected physical channels
phy₂: modeled as the sum of the physical effects of the device's individual operation
phy₃: modeled as a series of arbitrary values



Addressing Challenge 2: Detecting expanded vulnerabilities



 A_{tardy} (d) V4: Tardy-channel-based Trigger Interference. a_i a_j (e) V5: Disordered Action (f) V6: Action Overriding. physical Scheduling. **features** | A_{tardy} (g) V7: Action Breaking. (h) V8: Tardy-channel-based Condition Interference.

Basic Vulnerability Pattern (V1-V3)

Expanded Vulnerability Pattern (V4-V8)

Addressing Challenge 2: Detecting expanded vulnerabilities

Correctness property categorizing and ranking-based vulnerability detection

- Correctness Property: a criterion to describe what automation behavior is safe or not.
- Categorize 9 language templates of properties into 2 logical templates for property specification

Table 12: Logically equivalent correctness property types.

Summarised property types	Property types	Natural language templates		
Event-based	One-Event Unconditional	[event] should [never] hap pen		
	Event-State Conditional (always)	- [event] should [always] hap pen when [state ₁ ,, state _n]		
	Event-State Conditional (never)	[event] should [never] hap- pen when [state ₁ ,, state _n]		
	One-State Unconditional (always)	[state] should [always] be active		
State-based	One-State Unconditional (never)	[state] should [never] be active		

Addressing Challenge 2: Detecting expanded vulnerabilities

• Define pre- and post-proposition priority ranking to resolve property violations.

P.1: close the window if it rains

P.2: open the window if CO is detected

If it rains and CO is detected

Close or open window ??

Table 13: Sortin	g descriptions of	the pre-p	proposition	priority.
		T P I	I	L.

Scenarios in the pre-proposition of the correctness property	Pre-proposition priority		
General	user.not_present >user.present, smoke.detected = CO.detected >weather.raining > CO ₂ -related = humidity-related		
Temperature-related	user.not_present >heater.on = AC.on >the temperature is below / rises above a predefined value		

• Supply more properties for different scenarios (e.g., **properties with permitted latencies P.34**), finally conduct 53 properties for vulnerability detection.

P.31	WHEN CO is detected, the alarm should be activated.
P.32	IF humidity is greater than a predefined value, the ventilating fan should be turned on.
P.33	IF CO ₂ is greater than a predefined value, the ventilating fan should be turned on.
P.34	WHEN CO ₂ remains greater than a predefined value, the ventilating fan should be on for at least the permitted time.

• If the vulnerability exists, a system execution path that violates the correctness property (i.e., **counterexample**) will be returned by the model checker

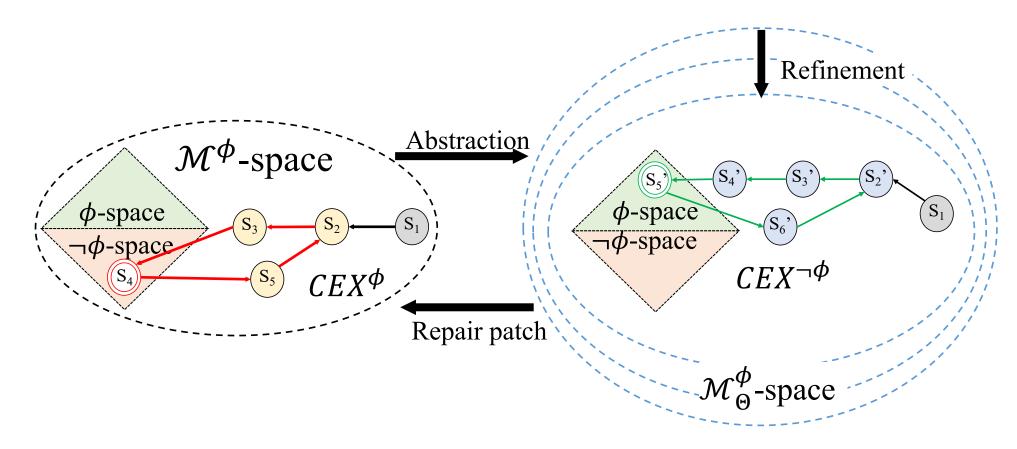
Addressing Challenge 3: Repairing expanded vulnerabilities

	Vulnerability Detection	How to repair vulnerability? Negated-property Reasoning (NPR) Algorithm
Applied Property	Correctness Property ϕ : IF the user is not at home, the heater should be turned off.	Negated-property $\neg \phi$ logically opposite to ϕ : IF the user is not at home, the heater should be turned on (negated).
Verification Process	model checking with ϕ on model \mathcal{M} $(\mathcal{M} \models \phi)$	model checking with $\neg \phi$ on model \mathcal{M} $(\mathcal{M} \vDash \neg \phi)$
Verification Result	Return a violation of ϕ in \mathcal{M} : scenarios where no one is home but the heater is running (fire hazard)	Return a violation of $\neg \phi$ in \mathcal{M} : scenarios where no one is home and the heater is off (potential fix information)

- Secure scenario reasoned by $\neg \phi$ can fix vulnerability violated ϕ
- Reasoned result of $\neg \phi$ does not violate ϕ , i.e., the reasoned result of $\neg \phi$ constitutes the repair space for the vulnerability violated ϕ

Addressing Challenge 3: Repairing expanded vulnerabilities

Negated-property reasoning (NPR) algorithm



- The core idea of our NPR algorithm for vulnerability repair.
- Model abstraction via interpolation is used to involve a larger state space for patch searching, so the negated counterexample $CEX \neg \phi$ can contain repair patches for the vulnerability $CEX \phi$.

Addressing Challenge 3: Repairing expanded vulnerabilities

Negated-property reasoning (NPR) algorithm

Step1: Negated Property Generation & Spurious Indicator Identification

- Negate the latter part of the LTL template divided by "⇒" to generate the negated property
- Spurious Indicator: assess whether a patch can fix the vulnerability
- Spurious indicator is the violating state in the counterexample

Step2: Patch Reasoning

- ullet Limited repair information in the state space of model ${\mathcal M}$
- Abstract model \mathcal{M} to $\mathcal{M}^{\phi}_{\Theta}$
- Reason patch P from abstract model $\mathcal{M}^{\phi}_{\Theta}$

Step3: Patch Feasibility Checking

- Patch Category
- Feasibility Checking
- Reasoning-guided Abstraction Refinement

Case Study of Vulnerability Detection and Repair Accuracy

Table 2: Accuracy comparison of the vulnerability detection. We use \checkmark , \circ , and \circ to denote true positive, false positive, and false negative, respectively.

Benchmark	SOATERIA*	SAFECHAIN	IOTCOM	TAPInspector*	TAPFixer
ID-1		€	€	✓	☑
ID-2	(8)	0	☑	♂	S
ID-3	€	0	€	✓	✓
ID-4	▼ 0	0	€	∀	☑
ID-5	0	∀	0	✓	☑
ID-6	⋖	0	€	⋖	⋖
ID-7		∀	✓	✓	S
ID-8		0	☑	✓	✓
ID-9	®	✓	✓	⋖	✓
N-1	0	0	0	∀	~
N-2	0	0	®	∀	✓
N-3	0	0	0	€	✓
N-6	0	0	0	0	✓
Gp-1	∀	0	€		S
Gp-2	✓	0	✓	⋖	✓
Gp-3		€		∀	⋖
Gp-4	0	0	✓	⋖	⋖
Gp-5	0	0	€	✓	⋖
Gp-6	0	0	€	✓	S
Gp-N4	0	0	0		⋖
Gp-N5	0	0	0	✓	✓

- Benchmark contains expanded vulnerabilities V4-V8
- TAPFixer is more effective at identifying and repairing expanded vulnerabilities

Table 4: Repair accuracy comparison of expanded vulnerabilities.

Benchmark	Liang et al. [35]	MenShen [18]	AutoTap [48]	TAPFixer
Group 1	3 †	®	•‡	✓
Group 2	®	®	®	⋖
Group 3	•	®	®	⋖
Group 4	0	0	®	⋖
Group 5	0	0	∀	⋖
N/A 1	0	0	€	⋖
N/A 2	0	0	€	✓

[†] Correctly fixed partial vulnerable rule interactions, but did not fix the rest.

[‡] Correctly fixed partial vulnerable rule interactions, but incorrectly fixed the rest.

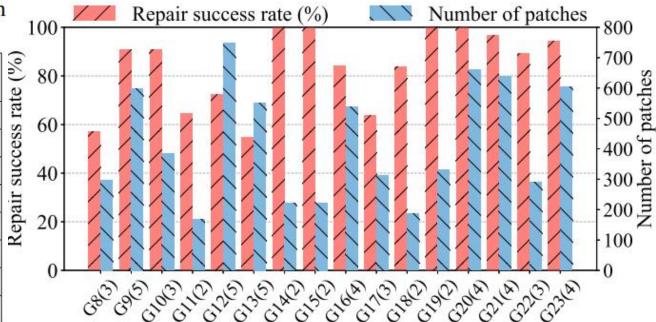
^{*} results obtained from [12, 20, 47]

Market App Study of Vulnerability Repair

- 1177 TAP rules from SmartThings SmartApp and IFTTT applets and 110 test groups
- Scenario-based Vulnerability Repair: repair 4544 of 5244 found vulnerabilities **Repair Success Rate (RSR): 86.65%**
- Priority-based Violation Repair: repair 4460 of 5335 found vulnerabilities Repair Success Rate (RSR): 83.60%

Table 5: Summary of detection and repair results for G1-G7 with 110 rule groups, each of which contains 15-30 TAP rules.

Application scenarios	Fixed violations	Unfixable violations	Safe cases	Generated patches	RSR↑
G1 (2 properties)	179	35	6	364	83.64%
G2 (21 properties)	1675	277	358	2228	85.81%
G3 (6 properties)	459	201	0	902	69.55%
G4 (8 properties)	687	59	134	1145	92.09%
G5 (9 properties)	870	68	52	1288	92.75%
G6 (3 properties)	272	58	0	491	82.42%
G7 (4 properties)	402	2	36	419	99.50%
Total	4544	700	586	6837	86.65%



Market App Study of Vulnerability Repair

- Comparison with the SOTA approach
- Modeling Success Rate (MSR): assess the integrity of rule modeling
- Repair Failure Rate (RFR)

RFR-MF: RFR caused by modeling failures

RFR-LIMIT: RFR caused by modeling failures repair algorithm limitations

Table 6: Comparison between AutoTap and TAPFixer.

	- ASS.	
Evaluation target	AutoTap	TAPFixer
MSR)	54.23%	100%
RSR of G1 (1 property)	20%	44%
RSR of G2 (14 properties)	51.43%	74.59%
RSR of G3 (1 property)	8%	92%
RSR of G4 (4 properties)	44%	94.32%
RSR of G7 (2 properties)	38%	91.49%
RFR-MF/RFR-LIMIT	23.99%/24.57%	0%/20.93%

User Survey on the Quality of Vulnerability Detection and Repair

- RSR: 94.5%
- Satisfaction Rate of the Detection and Repair Quality: 99%

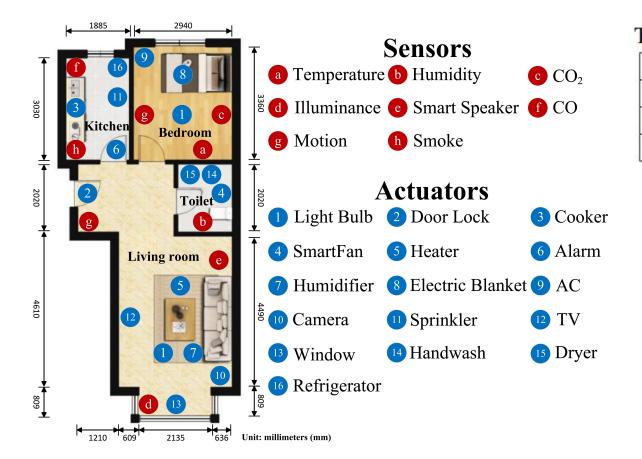
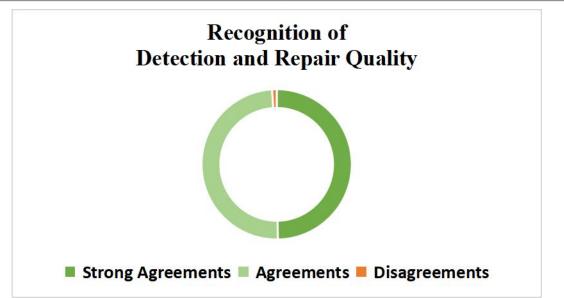


Table 7: Number of identified and fixed vulnerabilities in 129 rules.

	V1	V2	V3	V4	V5	V6	V7	V8
# found violations	5	6	9	23	0	4	5	3
# fixed violations	3	6	8	23	0	4	5	3
RSR	60%	100%	89.9%	100%	N/A	100%	100%	100%



Performance

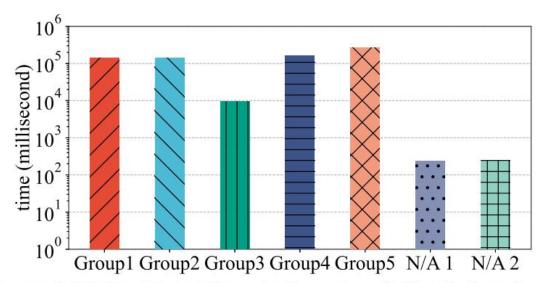


Figure 7: Verification and repair time of each 21-rule benchmark dataset and initialization scenarios.

Table 8: Average patch generation time for market apps.

Market apps	Total time (minute)	Number of generated patches	Avg. generation time per patch (second)
G1-G7 (110 rule groups)	364.070	6837	3.195
G8-G23 (110 rule groups)	431.261	6749	3.834

For case study:

- 21-rule benchmarks (Group 1-5) Average time 2.69 min
- Initialization Scenarios (N/A 1-2)
 Average time 228 ms

For 110 test groups in market apps:

- Average time 6.62 h
- Average time to reason a patch 3.51s

Conclusion

- We propse the physical model-based HA system modeling that can model rules with more practical latency and physical features.
- We propose the correctness property categorizing and ranking-based vulnerability detection that can identify more expanded interaction vulnerabilities.
- We propose a novel negated-property reasoning algorithm (NPR) that can accurately generate valid patches for eliminating vulnerabilities both in the logical and physical space.
- We implement **TAPFixer**, the first framework that can essentially detect and repair rule interaction vulnerabilities both in the logical and physical space.
- TAPFixer achieves very good results from aspects of accuracy analysis, repair capabilities of market apps, real user study, and execution performance.

Thanks

Q&A